# **ESDP Relaxation of Sensor Network Localization**

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# **Talk Outline**

- Sensor network localization
- SDP and ESDP relaxations: formulations and examples
- Properties of ESDP relaxations
- A robust version of ESDP for the noisy case
- Conclusion & Ongoing work

## **Sensor Network Localization**

• n pts in  $\mathbb{R}^2$ .

• Know last n - m pts ('anchors')  $x_{m+1}, ..., x_n$  and Eucl. dist. estimate for some pairs of 'neighboring' pts (i.e. within 'radio range')

$$d_{ij} \ge 0 \quad \forall (i,j) \in \mathcal{A}$$

with  $\mathcal{A} \subseteq \{(i, j) : 1 \leq i < j \leq n\}.$ 

• Estimate first m pts ('sensors').

ESDP RELAXATION OF SENSOR NETWORK LOCALIZATION Easy cases:

- When there are lots of anchors.
- When sensors neighbor lots of anchors.

### Hard cases:

- When m is large, n m is small, and sensors have few neighbors.
- Dist. measurement can have noise.

# **Optimization Problem Formulation**

$$v_{\text{opt}} := \min_{x_1, \dots, x_m} \sum_{(i,j) \in \mathcal{A}} |\|x_i - x_j\|^2 - d_{ij}^2|$$

- Objective function is nonconvex. m can be large (m > 1000).
- Problem is NP-hard (reduction from PARTITION).
- Use a convex relaxation.

### **SDP Relaxation**

Let  $X := [x_1 \cdots x_m].$ 

SDP relaxation (Biswas, Ye '03):

$$\begin{split} \upsilon_{\mathrm{sdp}} &:= \min_{Z} \sum_{\substack{(i,j) \in \mathcal{A}, j > m \\ + \sum_{\substack{(i,j) \in \mathcal{A}, j \le m \\ (i,j) \in \mathcal{A}, j \le m \\ }} \left| Y_{ii} - 2Y_{ij} + Y_{jj} - d_{ij}^2 \right| \\ \text{s.t.} \quad Z &= \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \succeq 0 \end{split}$$

Adding the nonconvex constraint rank Z = 2 yields original problem. Expensive to solve for *m* large.

Relax  $Z \succeq 0$  further.

### **ESDP Relaxation**

ESDP relaxation (Wang, Zheng, Boyd, Ye '07):

$$\begin{split} \upsilon_{\text{esdp}} &\coloneqq \min_{Z} \sum_{\substack{(i,j) \in \mathcal{A}, j > m \\ (i,j) \in \mathcal{A}, j > m \\ + \sum_{\substack{(i,j) \in \mathcal{A}, j \leq m \\ (i,j) \in \mathcal{A}, j \leq m \\ (i,j) \in \mathcal{A}, j \leq m \\ \text{s.t.} \quad Z = \begin{bmatrix} Y & X^T \\ X & I \\ \begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix}} \succeq 0 \quad \forall (i,j) \in \mathcal{A}, j \leq m \\ \begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \succeq 0 \quad \forall i = 1, \cdots, m \end{split}$$

In simulation, ESDP is nearly as strong as SDP, and solvable much faster by IP method.

ESDP seems more amenable to solution by distributed method (compared to SDP), which is important for practical implementation.

### **An Example**

$$n = 3, m = 1, d_{12} = d_{13} = 2$$

Problem:

$$0 = \min_{x_1 \in \Re^2} |||x_1 - (1,0)||^2 - 4| + |||x_1 - (-1,0)||^2 - 4|$$



### ESDP Relaxation

$$0 = \min_{\substack{x_1 = (x_{11}, x_{12}) \in \Re^2 \\ Y_{11} \in \Re}} |Y_{11} - 2x_{11} - 3| + |Y_{11} + 2x_{11} - 3|}$$
  
s.t. 
$$\begin{bmatrix} Y_{11} & x_{11} & x_{12} \\ x_{11} & 1 & 0 \\ x_{12} & 0 & 1 \end{bmatrix} \succeq 0$$



If solve ESDP by IP method, then likely get analy. center.

### **SDP & ESDP Relaxation: a larger example**

- n = 220, m = 200. Anchors (" $\diamond$ ") and sensors (" $\circ$ ") uniformly distributed in  $[-.5, .5]^2$ .
- $(i,j) \in \mathcal{A}$  whenever  $\|x_i^{\text{true}} x_j^{\text{true}}\| \le 0.15$
- The soln found by SeDuMi 1.05 (implementation of a (primal-dual) IP method in C with a Matlab interface) is shown ("\*") joined to its true position ("o") by a line.
- Run on CPU Intel Celeron 1.6 GHz, 448Mb RAM.



### **ESDP: Testing Solution Accuracy**

Aim: Determine which sensors are accurately positioned by ESDP.

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### Two cases:

•  $\exists x_1, \dots, x_m$  such that  $\|x_i - x_j\| = d_{ij}$  for  $(i, j) \in \mathcal{A}$  (Noiseless);

• 
$$d_{ij}^2 = ||x_i^{\text{true}} - x_j^{\text{true}}||^2 + \delta_{ij}$$
 (Noisy).

**Defn**:  $x_i$  is *invariant* over Sol (ESDP) if  $\forall Z \in Sol (ESDP)$ , its  $x_i$  entry is the same.

- $v_{\text{esdp}} = 0$ .
- Sol (P)  $\subseteq \left\{ X : \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix} \in$ Sol (ESDP)  $\right\}$
- Observation: If  $x_i$  is invariant over Sol (ESDP), then  $x_i$  is invariant over Sol (P) and hence  $x_i = x_i^{\text{true}}$ .

• Consider the individual trace for each sensor location i:  $\operatorname{tr}_i[Z] := Y_{ii} - \|x_i\|^2, \quad i = 1, ..., m.$  (Biswas, Ye '03)

Note:

- $\operatorname{tr}_i[Z] \ge 0$  for all Z feasible for ESDP;
- $\operatorname{tr}_i[Z] = 0$  for some  $Z \in \operatorname{ri}(\operatorname{Sol}(\operatorname{ESDP})) \Leftrightarrow \operatorname{tr}_i[Z] = 0$  for all  $Z \in \operatorname{Sol}(\operatorname{ESDP})$ .

- Consider the individual trace for each sensor location *i*:  $\operatorname{tr}_i[Z] := Y_{ii} - ||x_i||^2, \quad i = 1, ..., m.$
- (Wang et al '07) If  $tr_i(Z) = 0$  for some  $Z \in ri(Sol(ESDP))$ , then  $x_i$  is invariant.

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- (Wang et al '07) If  $tr_i(Z) = 0$  for some  $Z \in ri(Sol(ESDP))$ , then  $x_i$  is invariant.
- (P, Tseng '08) The converse holds: if  $x_i$  is invariant, then  $tr_i(Z) = 0$  for all  $Z \in Sol(ESDP)$ .

### Proof sketch for the converse implication:

**1.** For  $(i, j) \in A$ , j > m, if  $x_i$  is invariant over Sol(ESDP), then  $tr_i(Z) = 0$  for all  $Z \in Sol(ESDP)$ .

**Why:**  $v_{opt} = 0$  and  $x_i$  invariant over Sol(ESDP) imply, for any  $Z \in Sol(ESDP)$ ,

$$Y_{ii} - 2x_j^T x_i + ||x_j||^2 = d_{ij}^2, \qquad ||x_i - x_j||^2 = d_{ij}^2$$

So 
$$\operatorname{tr}_i(Z) = Y_{ii} - ||x_i||^2 = d_{ij}^2 - ||x_i - x_j||^2 = 0.$$

**2.** For  $(i, j) \in A$ ,  $j \leq m$ , if  $x_i, x_j$  are invariant over Sol(ESDP), then  $\operatorname{tr}_i(Z) = \operatorname{tr}_j(Z)$  for all  $Z \in \operatorname{Sol}(ESDP)$ .

Why?  $v_{opt} = 0$  and  $x_i, x_j$  invariant over Sol(ESDP) imply,

for any  $Z \in Sol(ESDP)$ ,

$$Y_{ii} - 2Y_{ij} + Y_{jj} = d_{ij}^2, \qquad ||x_i - x_j||^2 = d_{ij}^2$$
  
So  $Y_{ij} - x_i^T x_j = \frac{1}{2} (\operatorname{tr}_i(Z) + \operatorname{tr}_j(Z)).$ 

Then

This is psd, which implies ...that  $tr_i(Z) = tr_j(Z)$ .

### **Noisy Case**

Suppose dist. measurement has errors:

$$d_{ij}^2 = d_{ij}^{\text{true}^2} + \delta_{ij},$$

where  $\delta_{ij} \in \mathbb{R}$  and  $d_{ij}^{\text{true}} := \|x_i^{\text{true}} - x_j^{\text{true}}\|$ .

Does  $tr_i[Z] = 0$  (with  $Z \in ri(Sol(ESDP))$ ) imply  $x_i$  is near the true position of sensor *i*?

### **Noisy Case**

Unfortunately, no.

• (P, Tseng '08): For  $|\delta_{ij}| \approx 0$ ,

$$\operatorname{tr}_i[Z] = 0$$
 for some  $Z \in \operatorname{ri}(\operatorname{Sol}(\operatorname{ESDP})) \not\Longrightarrow ||x_i - x_i^{\operatorname{true}}|| \approx 0.$ 

Proof is by counter-example.



An example of sensitivity of ESDP solns to measurement noise:



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Thus, even when  $Z \in Sol(ESDP)$  is unique,  $tr_i[Z] = 0$  certifies accuracy of  $x_i$  only in the noiseless case! (Still true when changed to SDP)

### **Robust ESDP**

Fix  $\rho > \delta_{\max} := \max_{(i,j) \in \mathcal{A}} |\delta_{ij}|$ .

Sol(
$$\rho$$
ESDP) denotes the set of  $Z = \begin{bmatrix} Y & X^T \\ X & I \end{bmatrix}$  satisfying  
$$\begin{bmatrix} Y_{ii} & Y_{ij} & x_{ij}^T \end{bmatrix}$$

$$\begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \succeq 0 \quad \forall (i,j) \in \mathcal{A}, j \leq m$$
$$\begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \succeq 0 \quad \forall i = 1, \cdots, m$$
$$|Y_{ii} - 2x_j^T x_i + ||x_j||^2 - d_{ij}^2| \leq \rho \quad \forall (i,j) \in \mathcal{A}, j > m$$
$$|Y_{ii} - 2Y_{ij} + Y_{jj} - d_{ij}^2| \leq \rho \quad \forall (i,j) \in \mathcal{A}, j \leq m$$

Note: 
$$\begin{bmatrix} X^{\text{true}} & I \end{bmatrix}^T \begin{bmatrix} X^{\text{true}} & I \end{bmatrix} \in \text{Sol}(\rho \text{ESDP}).$$
  
Let

$$Z^{\rho} := \underset{Z \in \text{Sol}(\rho \text{ESDP})}{\arg\min} \quad - \quad \sum_{(i,j) \in \mathcal{A}, j \le m} \ln \det \left( \begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \right)$$
$$- \quad \sum_{i \le m} \ln \det \left( \begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \right)$$

Note: 
$$\begin{bmatrix} X^{\text{true}} & I \end{bmatrix}^T \begin{bmatrix} X^{\text{true}} & I \end{bmatrix} \in \text{Sol}(\rho \text{ESDP}).$$
  
Let

$$Z^{\rho} := \underset{Z \in \text{Sol}(\rho \text{ESDP})}{\operatorname{arg\,min}} - \underset{(i,j) \in \mathcal{A}, j \leq m}{\sum} \ln \det \left( \begin{bmatrix} Y_{ii} & Y_{ij} & x_i^T \\ Y_{ij} & Y_{jj} & x_j^T \\ x_i & x_j & I \end{bmatrix} \right)$$
$$- \underset{i \leq m}{\sum} \ln \det \left( \begin{bmatrix} Y_{ii} & x_i^T \\ x_i & I \end{bmatrix} \right)$$

• (P, Tseng '08)  $\exists \tau, \overline{\rho} > 0$  such that for each i,

$$\operatorname{tr}_{i}[Z^{\rho}] < \tau \quad \forall 0 \leq \delta_{\max} < \rho \leq \overline{\rho}$$
$$\implies ||x_{i}^{\rho} - x_{i}^{\operatorname{true}}|| \leq \sqrt{2|\mathcal{A}| + n} \left(\operatorname{tr}_{i}(Z^{\rho})\right)^{\frac{1}{2}}$$

## **Conclusion & Ongoing work**

SDP and ESDP are stronger relaxations, but the soln is not stable relative to measurement noise. Lack soln accuracy certificate.

 $\rho$ ESDP is a weaker relaxation, but has more stable solns. Has soln accuracy certificate.

- Distributed method to compute  $Z^{\rho}$ ?
- Simulation and numerical testing?

# Thanks for coming!