# SDPNAL+: A MATLAB software package for large-scale SDPs with a user-friendly interface 

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SIAM ALA Conference 2018

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- SDP and SDP+ (variable is positive semidefinite and bounded)
- Some examples of SDP+
- User-friendly interface
- Phase I: An inexact symmetric Gauss-Seidel (sGS) ADMM for SDP+
- An sGS decomposition theorem for convex composite QP
- Phase II: An augmented Lagrangian method (ALM) for SDP+
- A semismooth Newton-CG (SNCG) method for solving ALM subproblems
- SDPNAL+: practical implementation of the 2 phase method
- Numerical experiments
$\mathbb{S}_{+}^{n}=$ cone of positive semidefinite matrices. Write $X \succeq 0$ if $X \in \mathbb{S}_{+}^{n}$.

$$
(\mathrm{SDP}) \min \left\{\langle C, X\rangle \mid \mathcal{A}(X)=b, X \in \mathbb{S}_{+}^{n}\right\}
$$

where $C \in \mathbb{S}^{n}, b \in \mathbb{R}^{m}$ are given data; $\mathcal{A}: \mathbb{S}^{n} \rightarrow \mathbb{R}^{m}$ is a linear map.

$$
(\mathrm{SDP}+) \min \left\{\langle C, X\rangle \mid \mathcal{A}(X)=b, X \in \mathbb{S}_{+}^{n}, X \in \mathcal{N}\right\}
$$

where $\mathcal{N}=\left\{X \in \mathbb{S}^{n} \mid L \leq X \leq U\right\}$ and $L, U$ are given bounds (entries allow to take $-\infty, \infty$ respectively).
Important case: $\mathcal{N}=\left\{X \in \mathbb{S}^{n} \mid X \geq 0\right\}$, i.e., DNN (doubly nonnegative) SDP.
(SDP) is solvable by powerful interior-point methods if $n$ and $m$ are not too large, say, $n \leq 2000, m \leq 10,000$.
$m$ large $\Rightarrow m \times m$ dense "Hessian" matrix cannot be stored explicitly. For $m=10^{5}$, needs 100GB RAM memory!

Current research interests focus on $n \leq 5000$ but $m \gg 10,000$.

SDPNAL was developed around 2008/09 for (SDP).
In 2012/13, it was extended to SDPNAL+ for (SDP+) directly without introducing extra equality constraints $X=Y$ to convert $X \in$ $\mathbb{S}_{+}^{n} \cap \mathcal{N}$ to $X \in \mathbb{S}_{+}^{n}$ and $Y \in \mathcal{N}$.

Now our solver SDPNAL+ can solve general SDP problems:
$($ genSDP $) \min \quad \sum_{i=1}^{N}\left\langle C_{i}, X_{i}\right\rangle$

$$
\begin{array}{ll}
\text { s.t. } & \sum_{i=1}^{N} \mathcal{A}_{i}\left(X_{i}\right)=b \quad \text { (equalities) } \\
& l \leq \sum_{i=1}^{N} \mathcal{B}_{i}\left(X_{i}\right) \leq u \quad \text { (inequalities) } \\
& X_{i} \in \mathbb{K}_{i} \text { (cone), } \quad X_{i} \in \mathcal{N}_{i} \text { (bounds), } i=1: N
\end{array}
$$

where $\mathbb{K}_{i}$ is either a PSD cone or nonnegative orthant. Currently extending $\mathbb{K}_{i}$ to other cones such as SOCP.

- Parallel IPM [Benson, Borchers, Fujisawa, ... 03-present]
- First-order gradient methods on NLP formulation (low accuracy) [Burer-Monteiro 03]
- Inexact IPM [Kojima, Toh 04]
- Gen. Lag. method on barrier-penalized dual [Kocvara-Stingl 03]
- ALM on primal SDP from relaxation of lift-and-project scheme [Burer-Vandenbussche 06]
- Boundary-point method: BCD-ALM on dual [Rendl et al. 06] Reg. methods for SDP $\equiv$ ADMM on dual [Malick-Povh-Rendl 09]
- SDPNAL: ADMM+SNCG-ALM on dual [Zhao-Sun-Toh 10]
- SDPAD: ADMM on dual [Wen et al. 10] (used SDPNAL template)
- 2EBD: hybrid proximal extra-gradient method on primal [Monteiro et al. 13] (used SDPNAL template)
- ADMM+: convergent sGS-ADMM on SDP + [Sun-Toh-Yang 15]
- SDPNAL+: SNCG-ALM on SDP+ [Yang-Sun-Toh 15]

In nearest correlation matrix problem, given data matrix $U \in \mathbb{S}^{n}$, we want to solve
(NCM) $\min _{X}\left\{\left.\frac{1}{2}\|H \circ(X-U)\|_{1} \right\rvert\, \operatorname{Diag}(X)=\mathbf{1}, X \succeq 0\right\}$
where $H \in \mathbb{S}^{n}$ has nonnegative entries and " o " is the Hardamard product.

In clustering, given data vectors $\left\{p_{i}\right\}_{i=1}^{n}$, the goal is to cluster them into $k$ clusters. A possible model [Peng-Wei 07] is:

$$
\min \left\{\langle D, X\rangle \mid\langle I, X\rangle=k, X \mathbf{1}=\mathbf{1}, X \in \mathbb{S}_{+}^{n}, X \geq 0\right\}
$$

where $D_{i j}=\left\|p_{i}-p_{j}\right\|^{2}$.
Note: $D$ can also be other affinity matrix.

A stable set $S$ is subset of $V$ such that no vertices in $S$ are adjacent. Maximum stable set problem: find $S$ with maximum cardinality. Let

$$
x_{i}=\left\{\begin{array}{ll}
1 & \text { if } i \in S \\
0 & \text { otherwise }
\end{array} \quad \Rightarrow|S|=\sum_{i=1}^{n} x_{i}\right.
$$

A common formulation of the max-stable-set problem:

$$
\begin{aligned}
\alpha(G):= & \max \left\{\left.|S|=\frac{1}{|S|} \sum_{i j} x_{i} x_{j} \right\rvert\, x_{i} x_{j}=0 \forall(i, j) \in \mathcal{E}, x \in\{0,1\}^{n}\right\} \\
& \Downarrow X:=x x^{T} /|S| \\
& \max \left\{\langle E, X\rangle \mid X_{i j}=0 \forall(i, j) \in \mathcal{E}, \quad\langle I, X\rangle=1\right\}
\end{aligned}
$$

SDP relaxation: $X=x x^{T} /|S| \Rightarrow X \succeq 0$, get

$$
\theta(G):=\max \left\{\langle E, X\rangle: X_{i j}=0 \forall(i, j) \in \mathcal{E},\langle I, X\rangle=1, X \succeq 0\right\}
$$

$\theta_{+}(G):=n(n+1) / 2$ additional constraints $X \geq 0$

Assign $n$ facilities to $n$ locations [Koopmans and Beckmann (1957)]
$A=\left(a_{i j}\right)$ where $a_{i j}=$ flow from facility $i$ to facility $j$
$B=\left(b_{k l}\right) \quad$ where $b_{k l}=$ distance from location $k$ to location $l$
cost of assignment $\pi=\sum_{i=1}^{n} \sum_{j=1}^{n} a_{i j} b_{\pi(i) \pi(j)}$
$\min _{P}\left\{\left\langle B \otimes A, \operatorname{vec}(P) \operatorname{vec}(P)^{T}\right\rangle \mid P\right.$ is $n \times n$ permutation matrix $\}$
SDP + relaxation [Povh and Rendl, 09]: relax $\operatorname{vec}(P) \operatorname{vec}(P)^{T}$ to the $n^{2} \times n^{2}$ variable $X \in \mathbb{S}_{+}^{n^{2}}$ and $X \geq 0$

$$
(\mathrm{QAP}) \min \left\{\langle B \otimes A, X\rangle \mid \mathcal{A}(X)-b=0, X \in \mathbb{S}_{+}^{n^{2}}, X \geq 0\right\}
$$

where the linear constraints (with $m=3 n(n+1) / 2$ ) encode the condition $P^{T} P=I_{n}, P \geq 0$.

Consider the NCM problem.

```
n = 100;
G = randn (n,n);
G = 0.5*(G + G');
model = ccp_model('NCM');
    X = var_sdp (n,n);
    model.add_variable(X);
    model.minimize(l1_norm(X-G));
    model.add_affine_constraint(map_diag(X)==ones(n,1));
model.solve;
```

Consider the $\theta+$ problem of a graph with adjacency matrix $G$.

```
n = 200;
G = triu(sprand(n,n,0.5),1);
[IE,JE] = find(G);
n = length(G);
```

model = ccp_model('theta');
$X=\operatorname{var}_{-} \operatorname{sdp}(\mathrm{n}, \mathrm{n})$;
model. add_variable(X);
model.maximize(sum (X));
model. add_affine_constraint (trace (X) == 1);
model. add_affine_constraint (X(IE, JE) == 0);
model. add_affine_constraint ( $\mathrm{X}>=0$ ) ;
model.solve;

$$
\begin{array}{ll}
\min & \operatorname{trace}\left(X^{(1)}\right)+\operatorname{trace}\left(X^{(2)}\right)+\operatorname{sum}\left(X^{(3)}\right) \\
\text { s.t. } & -X_{12}^{(1)}+2 X_{33}^{(2)}+2 X_{2}^{(3)}=4 \quad \text { (equalities) } \\
& 2 X_{23}^{(1)}+X_{42}^{(2)}-X_{4}^{(3)}=3 \\
& 2 \leq-X_{12}^{(1)}-2 X_{33}^{(2)}+2 X_{2}^{(3)} \leq 7 \quad \text { (inequalities) } \\
& X^{(1)} \in \mathbb{S}_{+}^{6}, X^{(2)} \in \mathbb{R}^{5 \times 5}, X^{(3)} \in \mathbb{R}_{+}^{7} \quad \text { (cones) } \\
& 0 \leq X^{(1)} \leq 10 E_{6}, \quad 0 \leq X^{(2)} \leq 8 E_{5} \quad \text { (bounds) }
\end{array}
$$

```
n1 = 6; n2 = 5; n3 = 7;
M = ccp_model('Example_simple');
X1=var_sdp(n1,n1); X2=var_nn(n2,n2); X3=var_nn(n3);
M.add_variable(X1,X2,X3);
M.minimize(trace(X1) + trace(X2) + sum(X3));
M.add_affine_constraint (-X1 (1, 2) +2* X2 (3,3) +2* X 3 (2)==4);
M.add_affine_constraint (2*X1 (2,3) + X2(4,2)-X3(4) == 3);
M.add_affine_constraint (2<=-X1 (1, 2) - 2* X2 ( 3, 3) +2* X 3 (2) <= 7);
M.add_affine_constraint(0 <= X1 <= 10);
M.add_affine_constraint(X2 <= 8);
M.solve;
```

For simplicity, consider only $\mathcal{N}=\left\{X \in \mathbb{S}^{n} \mid X \geq 0\right\}$.
Dual of SDP+ and its augmented Lagrangian function are given by:
(D) $\quad \min \left\{-\langle b, y\rangle+\delta_{\mathbb{S}_{+}^{n}}(S)+\delta_{\mathcal{N}}(Z) \mid \mathcal{A}^{*} y+S+Z=C\right\}$
(a linearly constrained convex problem with 3 blocks of variables);

$$
\begin{aligned}
& \mathcal{L}_{\sigma}(y, S, Z ; X)=-\langle b, y\rangle+\delta_{\mathbb{S}_{+}^{n}}(S)+\delta_{\mathcal{N}}(Z) \\
& \quad+\left\langle\mathcal{A}^{*} y+S+Z-C, X\right\rangle+\frac{\sigma}{2}\left\|\mathcal{A}^{*} y+S+Z-C\right\|^{2}
\end{aligned}
$$

(quadratic in $(y, S, Z)+$ nonsmooth terms in $S, Z)$
KKT conditions:

$$
\mathcal{R}_{\mathrm{KKT}}(y, S, Z ; X):=\left(\begin{array}{c}
A X-b \\
S-\Pi_{\mathbb{S}_{+}^{n}}(S-X) \\
Z-\Pi_{\mathcal{N}}(Z-X) \\
\mathcal{A}^{*} y+S+Z-C
\end{array}\right)=\left(\begin{array}{l}
0 \\
0 \\
0 \\
0
\end{array}\right) .
$$

## A directly extended ADMM for dual SDP+

Input $\left(y_{0}, S_{0}, Z_{0} ; X_{0}\right)$. For $k=0,1, \ldots$, let $\widehat{C}^{k}=C-\sigma^{-1} X^{k}$
(1a) $y^{k+1}=\operatorname{argmin}_{y \in \mathbb{R}^{m}} \mathcal{L}_{\sigma}\left(y, S^{k}, Z^{k} ; X^{k}\right)$
(1b) $S^{k+1}=\operatorname{argmin}_{S \in \mathbb{S}_{+}^{n}} \mathcal{L}_{\sigma}\left(y^{k+1}, S, Z^{k} ; X^{k}\right)=\Pi_{\mathbb{S}_{+}^{n}}\left(\widehat{C}^{k}-\mathcal{A}^{*} y^{k+1}-Z^{k}\right)$
(2) $Z^{k+1}=\operatorname{argmin}_{Z \in \mathcal{N}} \mathcal{L}_{\sigma}\left(y^{k+1}, S^{k+1}, Z ; X^{k}\right)=\Pi_{\mathcal{N}}\left(\widehat{C}^{k}-\mathcal{A}^{*} y^{k+1}-S^{k+1}\right)$
(3) $X^{k+1}=X^{k}+\tau \sigma\left(\mathcal{A}^{*} y^{k+1}+S^{k+1}+Z^{k+1}-C\right)$, where $\tau \in\left(0, \frac{1+\sqrt{5}}{2}\right)$ is the step-length.

Direct extension of 2-block ADMM is not guaranteed to converge [Chen-He-Ye-Yuan, v155, MP 2016]

## A convergent symmetric Gauss-Seidel (sGS) ADMM for dual SDP+

But sGS-ADMM is guaranteed to converge!

$$
\begin{aligned}
& \text { Input }\left(y_{0}, S_{0}, Z_{0} ; X_{0}\right) . \text { For } k=0,1, \ldots, \text { let } \widehat{C}^{k}=C-\sigma^{-1} X^{k} \\
& \text { (1a) } \widehat{\mathbf{y}}^{k+1} \approx \operatorname{argmin}_{y \in \mathbb{R}^{m}} \mathcal{L}_{\sigma}\left(y, S^{k}, Z^{k} ; X^{k}\right) \\
& \text { (1b) } S^{k+1}=\operatorname{argmin}_{S \in \mathbb{S}_{+}^{n^{\mathcal{L}}}} \mathcal{L}_{\sigma}\left(\widehat{\mathbf{y}}^{k+1}, S, Z^{k} ; X^{k}\right)=\Pi_{\mathbb{S}_{+}^{n}}\left(\widehat{C}^{k}-\mathcal{A}^{*} \widehat{\mathbf{y}}^{k+1}-Z^{k}\right) \\
& \text { (1c) } y^{k+1} \approx \operatorname{argmin}_{y \in \mathbb{R}^{m}} \mathcal{L}_{\sigma}\left(y, S^{k+1}, Z^{k} ; X^{k}\right) \\
& \text { (2) } Z^{k+1}=\operatorname{argmin}_{Z \in \mathcal{N}} \mathcal{L}_{\sigma}\left(y^{k+1}, S^{k+1}, Z ; X^{k}\right)=\Pi_{\mathcal{N}}\left(\widehat{C}^{k}-\mathcal{A}^{*} y^{k+1}-S^{k+1}\right) \\
& \text { (3) } X^{k+1}=X^{k}+\tau \sigma\left(\mathcal{A}^{*} y^{k+1}+S^{k+1}+Z^{k+1}-C\right)
\end{aligned}
$$

In Step 1, the AL function $\mathcal{L}_{\sigma}$ for the block $(y, S)$ has the form:

$$
\begin{gathered}
\mathcal{L}_{\sigma}(y, S) \equiv \delta_{\mathbb{S}_{+}^{n}}(S)+\frac{\sigma}{2}\left\|\mathcal{A}^{*} y+S+Z^{k}+\widehat{C}^{k}\right\|^{2}-\langle b, y\rangle \\
(\mathrm{QP} \text { in }(y, S)+\text { nonsmooth term in } S)
\end{gathered}
$$

(1a)-(1c) is equivalent to minimizing $\mathcal{L}_{\sigma}(y, S)+\mathrm{sGS}$ proximal term.
The steps are based on an sGS decomposition theorem.

Theorem Suppose that the KKT conditions of (SDP+) has a solution. Let $\left\{\left(y^{k}, S^{k}, Z^{k}, X^{k}\right)\right\}$ be the sequence generated by the inexact sGS-ADMM. Then $\left\{X^{k}\right\}$ converges to an optimal solution of (SDP+) and $\left\{\left(y^{k}, S^{k}, Z^{k}\right)\right\}$ converges to an optimal solution of its dual.
[1] D.F. Sun, K.C. Toh and L.Q. Yang, A convergent 3-block semi-proximal ADMM for conic programming with 4-type constraints, v25, SIOPT 2015.
[2] X.D. Li, D.F. Sun, K.C. Toh, A Schur complement based semiproximal ADMM for convex ..., v155, MP 2016. Schur-complement-ADMM
[3] X.D. Li, D.F. Sun, K.C. Toh, QSDPNAL: A two-phase augmented Lagrangian method for convex quadratic SDP, MPC 2018. Section 2: sGS decomposition theorem, Schur-complement-ADMM $=$ sGS-ADMM
[4] L. Chen, D.F. Sun, K.C. Toh, An efficient inexact symmetric Gauss-Seidel based majorized ADMM for ..., v161, MP 2017. inexact sGS-ADMM
[5] X.D. Li, D.F. Sun, K.C. Toh, A block sGS decomposition theorem for convex composite quadratic programming and its applications, MP 2018. sGS-ADMM $=$ Schur-complement-ADMM, sSOR-extension

Theorem [Han-Sun-Zhang, MOR 2018: exact version]
Let $\Omega_{\mathrm{KKT}} \neq \emptyset$ be the KKT solution set. Suppose that an error bound condition holds for $\mathcal{R}_{\text {KKT }}$ at an optimal solution $u^{*}=$ $\left(y^{*}, S^{*}, Z^{*}, X^{*}\right)$ that $u^{k}=\left(y^{k}, S^{k}, Z^{k}, X^{k}\right)$ converges to, i.e, $\exists$ $\eta, r>0$ s.t.

$$
\operatorname{dist}\left(u, \Omega_{\mathrm{KKT}}\right) \leq \eta\left\|\mathcal{R}_{\mathrm{KKT}}(u)\right\| \quad \forall u \in B_{r}\left(u^{*}\right)
$$

Then $\exists \mu \in(0,1)$ depending on $\eta$ s.t.

$$
\operatorname{dist}\left(u^{k+1}, \Omega_{\mathrm{KKT}}\right) \leq \mu \operatorname{dist}\left(u^{k}, \Omega_{\mathrm{KKT}}\right) \quad \forall k \text { sufficiently large. }
$$

Inexact version can be established via the analysis in [Chen-Sun-Toh, MP 2017] and [Han-Sun-Zhang, MOR 2018].

Consider a convex composite QP with 3 blocks:

$$
\min \left\{p\left(x_{1}\right)+h(x) \mid x=\left(x_{1} ; x_{2}, x_{3}\right) \in \mathbb{R}^{n_{1}} \times \mathbb{R}^{n_{2}} \times \mathbb{R}^{n_{3}}\right\}
$$

Convex quadratic function $h(x):=\frac{1}{2}\langle x, \mathcal{H} x\rangle-\langle b, x\rangle$
Closed proper convex fun. $p: \mathbb{R}^{n_{1}} \rightarrow(-\infty,+\infty]$, e.g. $p\left(x_{1}\right)=$ $\left\|x_{1}\right\|_{\infty}$
Write $\mathcal{H}=\mathcal{U}^{*}+\mathcal{D}+\mathcal{U}, \mathcal{D}$ diagonal blocks, $\mathcal{U}$ strict upper triangular part. Assume $\mathcal{D}$ invertible.

Define $\operatorname{sGS}(\mathcal{H}):=\mathcal{U D}^{-1} \mathcal{U}^{*} \quad$ (symmetric Gauss-Seidel decomp)
Given $\bar{x}$, define

$$
x^{+}:=\operatorname{argmin}_{x}\left\{p\left(x_{1}\right)+h(x)+\frac{1}{2}\|x-\bar{x}\|_{\mathrm{sGS}(\mathcal{H})}^{2}\right\}
$$

Next theorem: can compute $x^{+}$using one sGS cycle! If $p\left(x_{1}\right)$ is absent, we get the classical block $\mathbf{s G S}$ iteration.

## Detour: block symmetric Gauss-Seidel (sGS) decomposition

## Theorem [Li-Sun-Toh 2015]

It holds that $\mathcal{H}+\operatorname{sGS}(\mathcal{H})=(\mathcal{D}+\mathcal{U}) \mathcal{D}^{-1}\left(\mathcal{D}+\mathcal{U}^{*}\right) \succ 0$.
Backward GS: $3 \rightarrow 2$. Compute

$$
\begin{aligned}
& x_{3}^{\prime}=\operatorname{argmin} p\left(\bar{x}_{1}\right)+h\left(\bar{x}_{1}, \bar{x}_{2}, x_{3}\right)=\mathcal{H}_{33}^{-1}\left(b_{3}-\mathcal{H}_{13}^{*} \bar{x}_{1}-\mathcal{H}_{23}^{*} \bar{x}_{2}\right) \\
& x_{2}^{\prime}=\operatorname{argmin} p\left(\bar{x}_{1}\right)+h\left(\bar{x}_{1}, x_{2}, x_{3}^{\prime}\right)=\mathcal{H}_{22}^{-1}\left(b_{2}-\mathcal{H}_{12}^{*} \bar{x}_{1}-\mathcal{H}_{23} \bar{x}_{3}^{\prime}\right)
\end{aligned}
$$

Forward GS: $1 \rightarrow 2 \rightarrow 3$. Compute

$$
\begin{aligned}
& x_{1}^{+}=\operatorname{argmin} p\left(x_{1}\right)+h\left(x_{1}, x_{2}^{\prime}, x_{3}^{\prime}\right) \quad \text { (non-smooth/non-quadratic) } \\
& x_{2}^{+}=\operatorname{argmin} p\left(x_{1}^{+}\right)+h\left(x_{1}^{+}, x_{2}, x_{3}^{\prime}\right)=\mathcal{H}_{22}^{-1}\left(b_{2}-\mathcal{H}_{12}^{*} x_{1}^{+}-\mathcal{H}_{23} x_{3}^{\prime}\right) \\
& x_{3}^{+}=\operatorname{argmin} p\left(x_{1}^{+}\right)+h\left(x_{1}^{+}, x_{2}^{+}, x_{3}\right)=\mathcal{H}_{33}^{-1}\left(b_{3}-\mathcal{H}_{13}^{*} x_{1}^{+}-\mathcal{H}_{23}^{*} x_{2}^{+}\right)
\end{aligned}
$$

Inexact computation is also allowed! So can use PCG to solve large linear systems.

Theorem [Li-Sun-Toh 2015]
Backward GS: For $i=s, \ldots, 2$, compute

$$
x_{i}^{\prime}=\mathcal{H}_{i i}^{-1}\left(b_{i}+e_{i}^{\prime}-\sum_{j=1}^{i-1} \mathcal{H}_{j i}^{*} \bar{x}_{j}-\sum_{j=i+1}^{s} \mathcal{H}_{i j} x_{j}^{\prime}\right) .
$$

Forward GS: For $i=2, \ldots, s$,

$$
\begin{aligned}
x_{1}^{+} & =\operatorname{argmin} p\left(x_{1}\right)+h\left(x_{1}, x_{\geq 2}^{\prime}\right)-\left\langle e_{1}^{+}, x_{1}\right\rangle, \\
x_{i}^{+} & =\mathcal{H}_{i i}^{-1}\left(b_{i}+e_{i}^{+}-\sum_{j=1}^{i-1} \mathcal{H}_{j i}^{*} x_{j}^{+}-\sum_{j=i+1}^{s} \mathcal{H}_{i j} x_{j}^{\prime}\right)
\end{aligned}
$$

$e^{+}, e^{\prime}$ are error vectors. In this case, $x^{+}$is the exact solution to a slightly perturbed proximal problem:
$x^{+}:=\operatorname{argmin}_{x}\left\{p\left(x_{1}\right)+h(x)+\frac{1}{2}\|x-\bar{x}\|_{\mathrm{sGS}(\mathcal{H})}^{2}-\left\langle x, \Delta\left(e^{\prime}, e^{+}\right)\right\rangle\right\}$
$\Delta\left(e^{\prime}, e^{+}\right)=e^{+}+\mathcal{U} \mathcal{D}^{-1}\left(e^{+}-e^{\prime}\right)$.

Adding a large proximal term slows the convergence of sGS-ADMM!
With no proximal term added, we consider the ALM for solving dual SDP+.
(1) Compute

$$
\begin{aligned}
& \left(y^{k+1}, S^{k+1}, Z^{k+1}\right) \approx \operatorname{argmin}\left\{\mathcal{L}_{k}(y, S, Z):=\mathcal{L}_{\sigma_{k}}\left(y, S, Z ; X^{k}\right)\right\} \\
= & \operatorname{argmin}\left\{-\langle b, y\rangle+\frac{\sigma}{2}\left\|\mathcal{A}^{*} y+S+Z+\widehat{C}^{k}\right\|^{2}+\delta_{\mathbb{S}_{+}^{n}}(S)+\delta_{\mathcal{N}}(Z)\right\}
\end{aligned}
$$

(2) Update $X^{k+1}=X^{k}+\sigma_{k}\left(\mathcal{A}^{*} y^{k+1}+S^{k+1}+Z^{k+1}-C\right)$; update $\sigma_{k+1} \uparrow \sigma_{\infty} \leq \infty$.

Define $X^{k+1}=X^{k}+\sigma_{k} R_{D}\left(y^{k+1}, S^{k+1}, Z^{k+1}\right)$,

$$
e^{k+1}=\left[\begin{array}{c}
\mathcal{A} X^{k+1}-b \\
X^{k+1}-\Pi_{\mathbb{S}_{+}^{n}}\left(X^{k+1}-S^{k+1}\right) \\
X^{k+1}-\Pi_{\mathcal{N}^{n}}\left(X^{k+1}-Z^{k+1}\right)
\end{array}\right] .
$$

In Step 1, we use the following easy-to-check stopping conditions:

$$
\begin{aligned}
& \text { (A) }\left\|e^{k+1}\right\| \leq \frac{\epsilon_{k}^{2}}{1+\left\|(X, y, S, Z)^{k+1}\right\|} \min \left\{\frac{1}{\sigma_{k}}, \frac{1}{1+\left\|X^{k+1}-X^{k}\right\|}\right\} \\
& \text { (B) }\left\|e^{k+1}\right\| \leq \frac{\eta_{k}^{2}\left\|X^{k+1}-X^{k}\right\|^{2}}{1+\left\|(X, y, S, Z)^{k+1}\right\|} \min \left\{\frac{1}{\sigma_{k}}, \frac{1}{1+\left\|X^{k+1}-X^{k}\right\|}\right\}
\end{aligned}
$$

where $\left\{\epsilon_{k}\right\}$ and $\left\{\delta_{k}\right\}$ are nonnegative summable sequences.
Theorem [Rockafellar 76] Let $\Omega_{P} \neq \emptyset$ be the primal optimal solution set and Slater's condition holds for primal problem (P). Under stopping condition (A), we have $X^{k} \rightarrow X^{*}$ and $\left(y^{k+1}, S^{k+1}, Z^{k+1}\right)$ converges to a dual optimal solution.

Theorem [Cui-Sun-Toh] If in addition, the blue stopping conditions are added, and the essential primal objective function $P^{\text {obj }}$ satisfies a quadratic growth condition at $X^{*}$, i.e., $\exists$ a neighborhood $\mathcal{U}$ of $X^{*}$ and $\kappa>0$ s.t.

$$
P^{\mathrm{obj}}(X) \geq P^{\mathrm{obj}}\left(X^{*}\right)+\kappa^{-1} \operatorname{dist}^{2}\left(X, \Omega_{P}\right) \quad \forall X \in \mathcal{U}
$$

Then for $k$ large, we have

$$
\operatorname{dist}\left(X^{k+1}, \Omega_{P}\right) \leq \theta_{k} \operatorname{dist}\left(X^{k}, \Omega_{P}\right)
$$

dual feasibility at $\left(y^{k+1}, S^{k+1}, Z^{k+1}\right) \leq \tau_{k} \operatorname{dist}\left(X^{k}, \Omega_{P}\right)$
dual objective gap at $\left(y^{k+1}, S^{k+1}, Z^{k+1}\right) \leq \tau_{k}^{\prime} \operatorname{dist}\left(X^{k}, \Omega_{P}\right)$
where $\theta_{k} \approx \frac{\kappa}{\sqrt{\kappa^{2}+\sigma_{k}^{2}}}, \quad \tau_{k} \approx \frac{1}{\sigma_{k}}, \quad \tau_{k}^{\prime} \approx \frac{\left\|X^{k}\right\|+\left\|X^{k+1}\right\|}{2 \sigma_{k}}$
Larger $\sigma_{k}$ gives faster convergence, but the inner problem is harder to solve.

For simplicity, assume $\mathcal{N}=\mathbb{S}^{n}$ and hence the variable $Z$ is absent.
$\operatorname{argmin}_{y, S}\left\{\mathcal{L}_{\sigma}(y, S) \equiv \delta_{\mathbb{S}_{+}^{n}}(S)+\frac{\sigma}{2}\left\|\mathcal{A}^{*} y+S-\widehat{C}^{k}\right\|^{2}-\langle b, y\rangle\right\}$
$\equiv \operatorname{argmin}_{y}\left\{\Phi^{k}(y):=-\langle b, y\rangle+\frac{\sigma}{2}\left\|\Pi_{\mathbb{S}_{+}^{n}}\left(\mathcal{A}^{*} y-\widehat{C}^{k}\right)\right\|^{2}\right\}($ project out $S)$
Optimality condition of unconstrained subproblem in $y$ is:

$$
\nabla \Phi^{k}(y)=-b+\sigma \mathcal{A} \Pi_{\mathbb{S}_{+}^{n}}\left(\mathcal{A}^{*} y-\widehat{C}^{k}\right)=0
$$

Solve for solution $y^{k+1}$ by the semismooth Newton-CG (SNCG) method. Then compute $S^{k+1}=\Pi_{\mathbb{S}_{+}^{n}}\left(\widehat{C}^{k}-\mathcal{A}^{*} y^{k+1}\right)$.
$\nabla \Phi^{k}(y)$ is not differentiable, but is strongly semismooth [Sun-Sun, 2002]. Thus SNCG is expected to have at least superlinear convergence.

## A semismooth Newton-CG method (SNCG) for ALM-subproblem

Solve $\nabla \Phi^{k}(y)=-b+\sigma \mathcal{A} \Pi_{\text {sin }_{+}^{n}}(U)=0, U=\mathcal{A}^{*} y-\widehat{C}^{k}$.
At the current iteration, $y_{l}$, we solve a generalized Newton equation:

$$
\begin{equation*}
\mathcal{H} \Delta y \approx \nabla \Phi^{k}\left(y_{l}\right), \quad \text { where } \mathcal{H} \Delta y=\sigma \mathcal{A} \Pi_{\mathbb{S}_{+}^{n}}^{\prime}(U)\left[\mathcal{A}^{*} \Delta y\right] \tag{1}
\end{equation*}
$$

From eigenvalue decomp: $U=Q D Q^{T}$ with $d_{1} \geq \cdots \geq d_{r} \geq 0>$ $d_{r+1} \geq \cdots \geq d_{n}$, we choose

$$
\begin{equation*}
\Pi_{\mathbb{S}_{+}^{n}}^{\prime}(U)[M]=Q\left(\Omega \circ\left(Q^{T} M Q\right)\right) Q^{T}, \tag{2}
\end{equation*}
$$

$\Omega_{i j}=\left(d_{i}^{+}-d_{j}^{+}\right) /\left(d_{i}-d_{j}\right)$. Let $\gamma=\{1, \ldots, r\}, \bar{\gamma}=\{r+1, \ldots, n\}$,

$$
\Omega=\left[\begin{array}{cc}
E_{\gamma \gamma} & \Omega_{\gamma \bar{\gamma}} \\
\Omega_{\bar{\gamma} \gamma} & 0
\end{array}\right] .
$$

When problem is primal nondegenerate, $\operatorname{cond}(\mathcal{H})$ is bounded:

$$
\operatorname{cond}(\mathcal{H}) \leq \sigma \Theta(1) \operatorname{cond}\left(\left[\mathcal{A} Q_{\gamma} \otimes Q_{\gamma}, \mathcal{A} Q_{\gamma} \otimes Q_{\bar{\gamma}}\right]\right)^{2}
$$

The structure in $\Omega$ allows for efficient computation of matrix-vector multiply for CG in solving (1). Direct evaluation of

$$
Y:=\Pi_{\mathbb{S}_{+}^{n}}^{\prime}(U)[M]=Q\left(\Omega \circ\left(Q^{T} M Q\right)\right) Q^{T}
$$

needs 4 matrix-matrix multiplications $=8 n^{3}$ operations. But with the structure of $\Omega$, can compute $Y$ as follows:

$$
Y=H+H^{T}, \quad H=Q_{\gamma}\left[\frac{1}{2}\left(U Q_{\gamma}\right) Q_{\gamma}^{T}+\left(\Omega_{\gamma \bar{\gamma}} \circ\left(U Q_{\bar{\gamma}}\right)\right) Q_{\bar{\gamma}}^{T}\right]
$$

where $U=Q_{\gamma} M$. The cost is at most $6 r n^{2}$.
If $r \approx n$, then use

$$
\begin{aligned}
Y & =Q\left(E \circ\left(Q^{T} M Q\right)\right) Q^{T}-Q\left(\bar{\Omega} \circ\left(Q^{T} M Q\right)\right) Q^{T} \\
& =M-Q\left(\bar{\Omega} \circ\left(Q^{T} M Q\right)\right) Q^{T}
\end{aligned}
$$

where $\bar{\Omega}=E-\Omega$ has a similar structure as $\Omega$ but with a large block of 0 . The cost is $6(n-r) n^{2}$.

Let ADMM + denote the sGS-ADMM.

1. Generate a good starting point to warm-start SNCG-ALM:

$$
\left(y^{0}, S^{0}, Z^{0}, X^{0}, \sigma_{0}\right) \leftarrow \mathrm{ADMM}+\left(\bar{y}^{0}, \bar{S}^{0}, \bar{Z}^{0}, \bar{X}^{0}, \bar{\sigma}_{0}\right)
$$

2. For $k=0,1, \ldots$

Generate $\left(y^{k+1}, S^{k+1}, Z^{k+1}\right)$ in ALM-subproblem via SNCG
Compute $X^{k+1}$ based on $\left(y^{k+1}, S^{k+1}, Z^{k+1}\right)$, update $\sigma_{k+1}$
If progress of SNCG-ALM is slow,
Rescale data
Let $\left(\bar{y}^{k}, \bar{S}^{k}, \bar{Z}^{k}, \bar{X}^{k}, \bar{\sigma}_{k}\right)$ denote rescaled $\left(y^{k}, S^{k}, Z^{k}, X^{k}, \sigma_{k}\right)$ Rescaling is chosen such that $\left\|\bar{X}^{k}\right\| \approx \max \left\{\left\|\bar{S}^{k}\right\|,\left\|\bar{Z}^{k}\right\|\right\}$
Goto Step 1: Restart with ADMM+( $\left.\bar{y}^{k}, \bar{S}^{k}, \bar{Z}^{k}, \bar{X}^{k}, \bar{\sigma}_{k}\right)$

$$
\eta \equiv \frac{\left\|\mathcal{R}_{\mathrm{KKT}}\left(y^{k+1}, S^{k+1}, Z^{k+1}, X^{k+1}\right)\right\|}{1+\left\|\left(y^{k+1}, S^{k+1}, Z^{k+1}, X^{k+1}\right)\right\|} \leq 10^{-6} .
$$

Performance of our SDPNAL+ and ADMM+ versus SDPAD: the directly extended ADMM implemented in [Wen et al.] 2EBD-HPE [Monteiro et al.]

Numbers of problems which are solved to the accuracy $\eta \leq 10^{-6}$

| problem set (No.) | SDPNAL+ | ADMM + | SDPAD | 2EBD |
| :---: | :---: | :---: | :---: | :---: |
| $\theta(58)$ | 58 | 56 | 53 | 53 |
| $\theta_{+}(58)$ | 58 | 58 | 58 | 56 |
| FAP (7) | 7 | 7 | 7 | 7 |
| QAP (95) | 95 | 39 | 30 | 16 |
| BIQ (134) | 134 | 134 | 134 | 134 |
| RCP (120) | 120 | 120 | 114 | 109 |
| R1TA (55) | 55 | 42 | 47 | 18 |
| Total (527) | 527 | 456 | 443 | 393 |

Performance profiles of SDPNAL+, ADMM + , SDPAD and 2EBD

Performance Profile ( 58 日, $58 \theta_{+}$, 7 FAP, 95 QAP, $134 \mathrm{BIQ}, 120$ RCP, 55 R1TA problems) tol $=1 \mathrm{e}-06$


## Numerical results for SDPNAL+

Implemented the algorithms in Matlab.
Runs perform on PC with (12 cores) Intel Xeon CPU E5-2680 @ 2.50 GHz and 128 GB RAM.

Stop SDPAD and 2EBD after 25000 iterations or 20 hours.

| Prob | $m ; n$ | $\begin{array}{c}\eta \\ \text { sDPAD\|2EBD }\end{array}$ SDPNAL+ |
| ---: | ---: | ---: | ---: |$)$

Results show that it is essential to use second-order information and second-order structured sparsity to solve hard problems!

- We have tested SDPNAL+ on about 520 SDPs from $\theta, \theta_{+}$, QAP, binary QP, rank-1 tensor approximation, etc
- When the problems are primal-dual nondegenerate, SDPNAL+ can efficiently solve large SDPs to high accuracy. SDPAD and 2EDB also performed well, though SDPNAL+ is often much more efficient.
- Many of the tested SDPs are degenerate, but SDPNAL+ can still solve them accurately with $\eta<10^{-6}$. On the other hand, SDPAD and 2EDB were not able to solve many such problems.

Currently under development:
(1) sparse SDPNAL+ so as to handle larger matrix variable when the data has conducive sparsity structure
(2) a more advanced user-friendly interface

Thank you for your attention!

