Merging Trust-Region and Limited Memory Technologies for Large-Scale Nonlinear Optimization

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Outline

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 - General format of large-scale nonlinear optimization problem.
 - Main difficulties.
- 2. Large-scale unconstrained problem.
 - Description of the algorithm.
 - Numerical Algebra.
 - Numerical Results.
- 3. Active-set Trust-region Algorithm(ASTRAL) for bounded constrained problems.
 - Description of the algorithm.
 - Numerical Algebra.
 - Numerical Results.

 $\begin{array}{ll} \min & f(x) \\ \text{s.t.} & l \leq x \leq u \end{array}$

Assumptions:

 f : ℝⁿ → ℝ is smooth, but difficult to evaluate. high dimensional integration simulation systems of PDE/ODE's multiple levels of optimization

2. $-\infty \leq I_i \leq u_i \leq \infty$

3. *n* is large. ($n \ge 1000$, e.g. $n = 2^{19}$ for imaging applications).

(1) Dominant Costs: function and gradient evaluations.

Jorge Moré and Nick Gould No test sets now available for *hard* to evaluate functions.

(2) Numerical linear algebra.

A very significant concern due to dimensionality. But, due to *expensive* functions, we want to exploit all data as fully as possible. (1) Dominant Costs: function and gradient evaluations.

line-search vs trust-region

(2) Numerical linear algebra.

matrix multiplies vs equation solves

The Unconstrained Problem

- Objective: $\min f(x)$
- Algorithm: $x^{k+1} = x^k + s^k$

Notation: $g^k = \nabla f(x^k), \quad B_k \approx \nabla^2 f(x^k)^{-1}, \quad H_k \approx \nabla^2 f(x^k)$

• Line-Search: $s^k = -\lambda_k B_k d^k$

 λ_k a stepsize (weak or strong Wolfe conditions). Requires repeated function and gradient evaluations.

• Trust-Region: s^k solves

$$\begin{array}{ll} \min & (g^k)^T s + \frac{1}{2} s^T H_k s \\ \text{s.t.} & \|s\| \le \Delta_k \end{array}$$

Hessian Approximations

Scalar Secant Equation:

$$H_k = \frac{\nabla f(x^k) - \nabla f(x^{k-1})}{x^k - x^{k-1}}$$

Matrix Secant Equation:

$$H_k(x^k - x^{k-1}) = \nabla f(x^k) - \nabla f(x^{k-1})$$

n linear equations in $\frac{n(n+1)}{2}$ unknowns

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BFGS Update:

 H_0 sym. and positive definite \Rightarrow H_k sym. and positive definite

$$H_{k+1} = H_{k} - \frac{H_{k}s^{k}s^{k^{T}}H_{k}}{s^{k^{T}}H_{k}s^{k}} + \frac{y^{k}y^{k^{T}}}{s^{k^{T}}y^{k}} \quad \text{whenever } s^{k^{T}}y^{k} > 0$$
$$= H_{k} - [y^{k} H_{k}s^{k}] \begin{pmatrix} -s^{k^{T}}y^{k} & 0\\ 0 & s^{k^{T}}H_{k}s^{k} \end{pmatrix}^{-1} [y^{k} H_{k}s^{k}]^{T}$$

where

$$s^k = x^{k+1} - x^k$$
 and $y^k = g^{k+1} - g^k$.

Compact *m*-Step Representation

Nocedal, Byrd and Schnabel (1994)

$$H_{k+m} = H_k - [Y \ H_k S] \begin{pmatrix} -D & L^T \\ L & s^{k^T} H_k s^k \end{pmatrix}^{-1} [Y \ H_k S]^T$$

where

$$S = [s^k, \cdots, s^{k+m-1}], \qquad Y = [y^k, \cdots, y^{k+m-1}],$$

and

$$S^T Y = L + D + R$$

with *L* strictly lower triangular, *D* diagonal, and *R* strictly upper triangular.

Limited Memory BFGS Updating

$$H = \lambda I - \Psi \Gamma^{-1} \Psi^{T}$$

where

$$\lambda = \frac{y^{k^T} y^k}{s^{k^T} y}, \qquad \Psi = [\Upsilon, \ \lambda S], \qquad \Gamma = \begin{bmatrix} -D & L^T \\ L & \lambda S^T S \end{bmatrix}_{2m \times 2m},$$

$$S^T Y = L + D + R$$

$$S = [s^{k_1}, \cdots, s^{k_m}], \qquad Y = [y^{k_1}, \cdots, y^{k_m}]$$

Typically, m = 5. λ -scaling – an approximate Rayleigh quotient Oren-Spedicato (1976), Phua-Shanno (1980), Barzilai-Borwein (1988).

The Trust-Region Subproblem

$$\begin{array}{ll} \min & (g^k)^T s + \frac{1}{2} s^T H_k s \\ \text{s.t.} & \|s\| \le \Delta_k \end{array}$$

Apply Newton's method to find $\mu > 0$ so that

$$\phi(\mu) = 0$$

where

$$\phi(\mu) = rac{1}{\Delta_k} - rac{1}{\|s(\mu)\|}$$
 and $s(\mu) = -(\mu I + H_k)^{-1} g^k$.
 $\mu_+ = \mu - rac{\phi(\mu)}{\phi'(\mu)},$ where $\phi'(\mu) = -rac{g^T (\mu I + H)^{-3} g}{\|s(\mu)\|^3}.$

(Our implementation avoids the hard case.)

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Powers of $(\mu I + H)^{-1}$

$$(\mu I + H)^{-1} = \frac{1}{\tau} [I + \Psi (\tau \Gamma - \Psi^T \Psi)^{-1} \Psi^T]$$

$$(\mu I + H)^{-2} =$$

$$\frac{1}{\tau^2}[I + \Psi(\tau \Gamma - \Psi^T \Psi)^{-1}\Psi^T + \tau \Psi(\tau \Gamma - \Psi^T \Psi)^{-1}\Gamma(\tau \Gamma - \Psi^T \Psi)^{-1}\Psi^T],$$

where $\tau = \mu + \lambda$.

Triangular Factorization of $\tau \Gamma - \Psi^T \Psi$

Set

$$\hat{D} = (\mu + \lambda)D + Y^T Y, \quad \hat{W} = \lambda \mu S^T S, \text{ and } \hat{L} = \mu L - \lambda (R + D).$$

Compute Cholesky factorizations

$$\hat{D} = M M^{T}$$
$$\hat{W} + \hat{L}\hat{D}^{-1}\hat{L}^{T} = J J^{T}$$

Then

$$\tau \Gamma - \Psi^{T} \Psi = \begin{bmatrix} M & 0 \\ -\hat{L}M^{-T} & J \end{bmatrix} \begin{bmatrix} -M^{T} & M^{-1}\hat{L}^{T} \\ 0 & J^{T} \end{bmatrix}$$

The trust-region Newton iteration occurs entirely in dimension 2*m*. Cost $\approx O(mn)$ assuming $m^3 \ll n$.

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TEST SET: 24 problems from MINPACK-2 set, $2500 \le n \le 160,000$.

Termination Criteria:

 $f_{\rm best} \sim$ best known function value.

1.
$$|f^k - f_{\text{best}}| / \max(1, |f_{\text{best}}|) \le \epsilon$$
.

 $\text{2. nf} \geq 1000.$

Performance Profile: (Dolen and Moré 2001)

Given a problem set \mathcal{P} ,

$$t_{i,s_j}$$
 = nfg(CPU time) of solving prob i by alg s_j

$$r_{i,\mathbf{s}_j} = rac{t_{i,\mathbf{s}_j}}{\min_j t_{i,\mathbf{s}_j}}.$$

Set

$$\rho_{s_j}(\tau) = \frac{1}{n_p} |\{i \in \mathcal{P} : r_{i,s_j} \leq \tau\}|.$$

Comparison of number of the function and gradient evaluations



(a) relative accuracy 10^{-5}

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Comparison of number of the function and gradient evaluations



(c) relative accuracy 10^{-5} (d) relative accuracy 10^{-3}

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TR-L-BFGS

Comparison of CPU time



(e) relative accuracy 10^{-5} (f) relative accuracy 10^{-3}

TR-L-BFGS

More Implementation Details

Initialization: x^0 , $g^0 = \nabla f(x^0)$, $H_0 \succeq 0$, $0 < \kappa < 1$, $0 < \sigma < 1$. Iteration:

1. Set
$$\bar{s} = -H_k^{-1}g^k$$
 and $r(\bar{s}) = \frac{f(x^k+\bar{s})-f(x^k)}{q(\bar{s})}$. (98% acceptance)

2. WHILE
$$r(\bar{s}) < \kappa$$

a. Let $\delta = \sigma ||\bar{s}||$.
b. Solve

$$\min_{\substack{s.t.\\ ||s|| \le \delta}} q(s) = s^T g^k + \frac{1}{2} s^T H_k s$$
s.t.

$$||s|| \le \delta.$$
c. Compute $r(\bar{s}) = \frac{f(x^k + \bar{s}) - f(x^k)}{q(\bar{s})}$.
END WHILE

3. Set
$$x^{k+1} = x^k + \bar{s}$$
. Update H_{k+1} .

Minimization with bounded constraints

$$\begin{array}{ll} (\mathcal{P}) & \min & f(x) \\ & I \leq x \leq u. \end{array}$$

Active-set Trust-region Algorithm (ASTRAL).

We use an ℓ^{∞} trust-region to conform with the constraint geometry.

$$\Omega = \{ \boldsymbol{x} \mid \boldsymbol{I} \leq \boldsymbol{x} \leq \boldsymbol{u} \}, \quad \mathbb{B}_{\infty} = \{ \boldsymbol{x} \mid \| \boldsymbol{x} \|_{\infty} \leq \boldsymbol{1} \}$$
$$\Omega_{k} = \Omega \cap (\boldsymbol{x}_{k} + \Delta_{k} \mathbb{B}_{\infty}) = \left\{ \boldsymbol{x} \mid \boldsymbol{I}^{k} \leq \boldsymbol{x} \leq \boldsymbol{u}^{k} \right\}.$$

and

Binding Constraints

Active Constraints

$$A(\mathbf{x}) = \{i \mid \mathbf{x}_i = \mathbf{I}_i \text{ or } \mathbf{x}_i = \mathbf{u}_i\}$$

Binding Constraints

$$\mathcal{B}(\boldsymbol{x}) = \left\{ i \left| \begin{array}{c} \boldsymbol{x}_i = \boldsymbol{l}_i \text{ and } (\nabla f(\boldsymbol{x}))_i > \boldsymbol{0}, \\ \text{or } \boldsymbol{x}_i = \boldsymbol{u}_i \text{ and } (\nabla f(\boldsymbol{x}))_i < \boldsymbol{0} \end{array} \right\} \right.$$

Non-Binding Constraints

$$\mathcal{B}^{c}(\mathbf{x}) = \{1, 2, \dots, n\} \setminus \mathcal{B}(\mathbf{x}), \qquad \nu(\mathbf{x}) = |\mathcal{B}^{c}(\mathbf{x})|.$$

 $\Phi(x)$ is an $n \times \nu(x)$ matrix whose columns are those of the identity matrix corresponding to the non-binding constraints $\mathcal{B}^{c}(x)$.

Active-set Trust-region Algorithm

1. Identify $\mathcal{B}(x^k)$ and set $\Phi_k = \Phi(x^k)$.

2. Set $\tilde{H}_k = \Phi_k^T H_k \Phi_k$, $\tilde{g}^k = \Phi_k^T g^k$, $\tilde{l}^k = \Phi_k^T l^k$, and $\tilde{u}^k = \Phi_k^T u^k$.

3. Solve trust-region subproblem

min
$$\frac{1}{2}s^T \tilde{H}_k s + \tilde{g}^k s$$

subject to $\tilde{l}^k \le s \le \tilde{u}^k$

4. Form the ratio
$$r = rac{f(x^k) - f(x^k + \bar{s})}{q_k(\bar{s})}$$
 and update.

Trust Region Sub-Probroblem Reduction

Let
$$w = s - \tilde{l}^k$$
, $h = \tilde{u}^k - \tilde{l}^k$.

$$\begin{array}{ll} \min & \frac{1}{2}s^{T}\tilde{H}_{k}s + \tilde{g}^{k}s & \min & \frac{1}{2}w^{T}\tilde{H}w + k^{T}w \\ \\ \text{s.t.} & \tilde{l}^{k} \leq s \leq \tilde{u}^{k} & \text{s.t.} & 0 \leq w \leq h. \end{array}$$

Since $\tilde{H} = \Phi^T H \Phi$ is positive definite, this QP is convex.

We solve using an interior point algorithm.

 $W = \operatorname{diag}(w)$

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Interior Point Newton Equations

$$p = z - v - k - t U^{-1} e + U^{-1} V h + t W^{-1} e$$

$$r = (\tilde{H} + U^{-1}V + W^{-1}Z)^{-1}p,$$

 $\Delta w = -w + r$

$$\Delta u = -w + h - u - \Delta w,$$

$$\Delta v = t U^{-1} e - v - U^{-1} V \Delta s,$$

$$\Delta z = -W^{-1}Z\Delta w + tW^{-1}e - z.$$

t = homotopy parameter.

$$(\tilde{H} + U^{-1}V + W^{-1}Z)^{-1}$$

$$\tilde{H} = \Phi^T H \Phi = \Phi^T (\lambda I - \Psi \Gamma^{-1} \Psi^T) \Phi$$

$$(\tilde{H} + U^{-1}V + W^{-1}Z)^{-1} = G^{-1} + G^{-1}(\Phi^{T}\Psi) \left[\Gamma - (\Phi^{T}\Psi)^{T}G^{-1}(\Phi^{T}\Psi) \right]^{-1} (\Phi^{T}\Psi)^{T}G^{-1}$$

where

$$\mathbf{G} = \lambda \mathbf{I} + \mathbf{U}^{-1}\mathbf{V} + \mathbf{W}^{-1}\mathbf{Z}.$$

TR-L-BFGS

Triangular Factorization

Write

$$\Gamma - (\Phi_k^T \Psi)^T G^{-1} (\Phi_k^T \Psi) = \begin{bmatrix} -\hat{D} & \hat{L}^T \\ \hat{L} & \hat{W} \end{bmatrix}.$$

Compute Cholesky factors

$$\hat{D} = MM^T$$
 and $\hat{W} + \hat{L}\hat{D}^{-1}\hat{L}^T = JJ^T$.

Then

$$\Gamma - (\Phi_k^T \Psi)^T G^{-1} (\Phi_k^T \Psi) = \begin{bmatrix} M & 0 \\ -\hat{L}M^{-T} & J \end{bmatrix} \begin{bmatrix} -M^T & M^{-1}\hat{L}^T \\ 0 & J^T \end{bmatrix}.$$

TEST SET: 23 problems from CUTEr set. $n \ge 1000$.21 problems have dimension ≥ 10000 .

Termination Criteria: $f_{best} = best$ known function value.

1.
$$|f^k - f_{\text{best}}| / \max(1, |f_{\text{best}}|) \le \epsilon$$
.

 $\text{2. nf} \geq 1000.$

Comparison of nfg between L-BFGS-B and ASTRAL

L-BFGS-B: Nocedal-Zhu-Byrd-Liu (1997)



Figure 1aFigure 1bFigure 1a. Performance profiles, sum of the function and gradient evaluations, relative
accuracy 10^{-5} . Figure 1b. Performance profiles, sum of the function and gradient
evaluations, relative accuracy 10^{-3} .

Comparison of CPU time

Note that the average cost of each function evaluation in our test set is 0.002s, and the average cost of each gradient evaluation is 0.008s.

Comparison of CPU time

Note that the average cost of each function evaluation in our test set is 0.002s, and the average cost of each gradient evaluation is 0.008s.





cpu(f) = 0.02s, cpu(g) = 0.08s

cpu(f) = 0.04s, cpu(g) = 0.16s

Thank You

Reference.

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