# The $Z$-eigenvalues of a symmetric tensor and its application to spectral hypergraph theory 

Guoyin $\mathrm{Li}^{1, *,}{ }^{1,}$, Liqun $\mathrm{Qi}^{2}$ and Gaohang $\mathrm{Yu}^{3}$<br>${ }^{1}$ Department of Applied Mathematics, University of New South Wales, Sydney 2052, Australia<br>${ }^{2}$ Department of Applied Mathematics, The Hong Kong Polytechnic University, Hung Hom, Kowloon, Hong Kong<br>${ }^{3}$ School of Mathematics and Computer Sciences, Gannan Normal University, Ganzhou, 341000, China

## SUMMARY

In this paper, using variational analysis and optimization techniques, we examine some fundamental analytic properties of $Z$-eigenvalues of a real symmetric tensor with even order. We first establish that the maximum $Z$-eigenvalue function is a continuous and convex function on the symmetric tensor space and so provide formulas of the convex conjugate function and $\epsilon$-subdifferential of the maximum $Z$-eigenvalue function. Consequently, for an $m$ th-order $n$-dimensional tensor $\mathcal{A}$, we show that the normalized eigenspace associated with maximum $Z$-eigenvalue function is $\rho$ th-order Hölder stable at $\mathcal{A}$ with $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. As a byproduct, we also establish that the maximum $Z$-eigenvalue function is always at least $\rho$ th-order semismooth at $\mathcal{A}$. As an application, we introduce the characteristic tensor of a hypergraph and show that the maximum $Z$-eigenvalue function of the associated characteristic tensor provides a natural link for the combinatorial structure and the analytic structure of the underlying hypergraph. Finally, we establish a variational formula for the second largest $Z$-eigenvalue for the characteristic tensor of a hypergraph and use it to provide lower bounds for the bipartition width of a hypergraph. Some numerical examples are also provided to show how one can compute the largest/second-largest $Z$-eigenvalue of a medium size tensor, using polynomial optimization techniques and our variational formula. Copyright © 2013 John Wiley \& Sons, Ltd.

Received 14 December 2011; Revised 6 January 2013; Accepted 4 February 2013
KEY WORDS: symmetric tensor; maximum $Z$-eigenvalue; semismoothness; spectral graph theory; characteristic tensor; polynomial optimization

## 1. INTRODUCTION

An $m$ th-order $n$-dimensional tensor $\mathcal{A}$ consists of $n^{m}$ entries in real number:

$$
\begin{equation*}
\mathcal{A}=\left(\mathcal{A}_{i_{1} i_{2} \ldots i_{m}}\right), \quad \mathcal{A}_{i_{1} i_{2} \ldots i_{m}} \in \mathbb{R}, \quad 1 \leqslant i_{1}, i_{2}, \ldots, i_{m} \leqslant n . \tag{1.1}
\end{equation*}
$$

We say a tensor $\mathcal{A}$ is symmetric if the value of $\mathcal{A}_{i_{1} i_{2} \ldots i_{m}}$ is invariant under any permutation of its index $\left\{i_{1}, i_{2}, \ldots, i_{m}\right\}$. Clearly, when $m=2$, a symmetric tensor is nothing but a symmetric matrix. A symmetric tensor uniquely defines an $m$ th-degree homogeneous polynomial function $f$ with real coefficient: for all $x=\left(x_{1}, \ldots, x_{n}\right)^{T} \in \mathbb{R}^{n}$,

$$
f(x)=\mathcal{A} x^{m}:=\sum_{i_{1}, \ldots, i_{m}=1}^{n} \mathcal{A}_{i_{1} i_{2} \cdots i_{m}} x_{i_{1}} \ldots x_{i_{m}}
$$

[^0]Over the last few years, there has been a realization that there is a reasonably complete and consistent theory of eigenvalues and singular values for tensors of higher order, proposed by Lim and Qi independently, that generalizes the theory of matrix eigenvalues and singular values in various manners and extent. Recently, numerical study on tensors also has attracted many researchers because of its wide applications in polynomial optimization [1], hypergraph theory [2, 3], higher-order Markov chain [4], signal processing [5], and image science [6]. In particular, various efficiently numerical schemes have been proposed to find the low rank approximations of a tensor and the eigenvalues/eigenvectors of a tensor with specific structure (cf. [7-12, 43]).

Among the various definitions of an eigenvalue of a symmetric tensor, there are two particular interesting definitions called $Z$-eigenvalues and $H$-eigenvalues (see the definition later on). Recall that a tensor is said to be positive semidefinite if the corresponding homogeneous polynomial function of the tensor uniquely determined always takes nonnegative values. As shown in [13], a tensor is positive semidefinite if and only if its $Z$-eigenvalues (resp. $H$-eigenvalues) are all nonnegative. So, the $Z$-eigenvalues and $H$-eigenvalues play an important role in determining whether a symmetric tensor is positive semidefinite or not. On the other hand, $Z$-eigenvalues and $H$-eigenvalues can be fundamentally different as investigated in [13]. For example, finding an $H$-eigenvalue of a symmetric tensor is equivalent to solving a homogeneous polynomial equation, whereas calculating a $Z$-eigenvalue is equivalent to solving nonhomogeneous polynomial equations. Moreover, a diagonal symmetric tensor A has exactly $n$ many $H$-eigenvalues and may have more than $n$ Z-eigenvalues (for more details, see [13]).

Very recently, in our preceding paper [14], we investigated the analytic properties of the maximum $H$-eigenvalue function of a symmetric tensor. In particular, we showed that, for an $m$ thorder $n$-dimensional symmetric tensor $\mathcal{A}$, the $H$-maximum eigenvalue function is $\frac{1}{(2 m-1)^{n}}$ th-order semismooth at $\mathcal{A}$ when the geometric multiplicity of $\mathcal{A}$ is one. As an application, we proposed a generalized Newton method to solve the space tensor conic linear programming (STCLP) problem that arises in medical imaging area. Local convergence rate of this method was established by using the semismooth property of the maximum $H$-eigenvalue function. In this paper, we continue our study and examine the analytic properties of $Z$-eigenvalues. We first show that the maximum $Z$-eigenvalue function is continuous, convex, and differentiable almost everywhere, extending the fundamental analytic properties of the maximum eigenvalue of a symmetric matrix. Then, we establish that the normalized eigenspace associated with maximum $Z$-eigenvalue function is $\rho$ th-order Hölder stable at $\mathcal{A}$ with $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. As a by-product, without the geometric multiplicity assumption, we also establish that the maximum $Z$-eigenvalue function is always at least $\rho$ th-order semismooth at $\mathcal{A}$. As an application, we introduce the characteristic tensor of a hypergraph and show that the maximum $Z$-eigenvalue function of the associated characteristic tensor provides a natural link for the combinatorial structure and the analytic structure of the underlying hypergraph. A variational formula for the second largest $Z$-eigenvalue for the characteristic tensor of a hypergraph is also provided.

The organization of this paper is as follows. We first fix the notations and collect some basic definitions in Section 2. In Section 3, by using the variational analysis techniques, we show that the maximum $Z$-eigenvalue function is continuous and convex and hence differentiable almost everywhere. In particular, we obtain the formula for calculating the convex conjugate of and $\epsilon$-convex subdifferential for the maximum $Z$-eigenvalue function. In Section 4, for an $m$ th-order $n$-dimensional tensor $\mathcal{A}$, we show that the normalized eigenspace associated with maximum $Z$-eigenvalue function is $\rho$ th-order Hölder stable at $\mathcal{A}$ with $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. We also show that the maximum $Z$-eigenvalue function is always at least $\rho$ th-order semismooth at $\mathcal{A}$. Sufficient condition ensuring the strong semismoothness of the maximum $Z$-eigenvalue function is also provided. In Section 5, we introduce the characteristic tensor of a hypergraph and show that the maximum $Z$-eigenvalue function of the associated characteristic tensor provides a natural link for the combinatorial structure and the analytic structure of the underlying hypergraph. Moreover, we establish a variational formula for the second largest $Z$-eigenvalue for the characteristic tensor of a hypergraph. In Section 6, numerical examples are provided to show how one can compute the largest and second largest $Z$-eigenvalues of a real symmetric tensor using polynomial
optimization technique. Finally, we conclude our paper and present some future research topics in Section 7.

## 2. PRELIMINARIES

In this section, we fix the notations and collect some basic definitions and facts that we will use later on. Let $X, Y$ be finite dimensional inner product spaces. We use $\mathbb{B}_{X}$ (resp. $\mathbb{B}_{Y}$ ) to denote the unit open ball in $X$ (resp. $Y$ ). Denote the space of all linear map from $X$ to $Y$ by $L(X, Y)$. The norm of $X$ is defined by $\|x\|=\sqrt{\langle x, x\rangle_{X}}$ for all $x \in X$, where $\langle\cdot, \cdot\rangle_{X}$ is the inner product in $X$. Consider a locally Lipschitz function $G: X \rightarrow Y$. By the Rademacher's theorem, $G$ is differentiable almost everywhere on $X$. Let $D_{G}$ be the set consisting of all the points where $G$ is differentiable. Then, for any $x \in D_{G}$, the derivative of $G, \nabla G(x)$ exists. Denote

$$
J_{B} G(x)=\left\{V \in L(X, Y): V=\lim _{x_{k} \rightarrow x} \nabla G\left(x_{k}\right), x_{k} \in D_{G}\right\}
$$

Then, its Clarke's generalized Jacobian [41, 42] is defined by $J_{C} G(x)=\operatorname{conv} J_{B} G(x)$. In particular, if $Y=\mathbb{R}$ and $G=g$ where $g: X \rightarrow \mathbb{R}$ is locally Lipschitz, by identifying $X^{*}$ as $X$, then the Clarke's generalized Jacobian reduces to the Clarke's subdifferential defined by

$$
\partial_{C} g(x)=\left\{\xi \in X:\langle\xi, v\rangle_{X} \leqslant g^{\circ}(x ; v) \text { for all } v \in X\right\}
$$

where $\langle\cdot, \cdot\rangle_{X}$ is the inner product in $X$ and $g^{\circ}(x ; v)$ is the Clarke directional derivative of $g$ at the point $x$ in the direction $v$ given by

$$
g^{\circ}(x ; v)=\limsup _{y \rightarrow x, t \downarrow 0} \frac{g(y+t v)-g(y)}{t}
$$

The Clarke subdifferential $\partial_{C} g(x)$ is a nonempty, convex, and compact subset of $X$ for each $x \in X$. Recall that $g$ is convex on $X$ if

$$
g\left(\lambda x_{1}+(1-\lambda) x_{2}\right) \leqslant \lambda g\left(x_{1}\right)+(1-\lambda) g\left(x_{2}\right), \forall \lambda \in[0,1] \text { and } x_{1}, x_{2} \in X
$$

For each $\epsilon \geqslant 0$, define the $\epsilon$-convex subdifferential of $g$ at $x$ by

$$
\partial_{\epsilon} g(x):=\left\{\xi \in X:\langle\xi, z-x\rangle_{X} \leqslant g(z)-g(x)+\epsilon \text { for all } z \in X\right\}
$$

If $\epsilon=0$, we simply call it convex subdifferential of $g$ and denote it by $\partial g(x)$. If $g$ is convex on $X$, then $\partial_{C} g(x)=\partial g(x)$ for all $x \in X$. The Fenchel conjugate function of a convex function $g$ on $X$ is denoted by $g^{*}$ and is defined by

$$
g^{*}(\xi)=\sup _{x \in X}\left\{\langle\xi, x\rangle_{X}-g(x)\right\} \text { for all } \xi \in X
$$

An important property relating the Fenchel conjugate and the $\epsilon$-subdifferential is the so-called generalized Fenchel inequality (cf. [15, Theorem 2.4.2 (ii)])

$$
g(x)+g^{*}(\xi) \leqslant\langle\xi, x\rangle_{X}+\epsilon \Leftrightarrow \xi \in \partial_{\epsilon} g(x)
$$

We are now ready to state the definitions of semismooth functions and $\rho$ th-order semismooth functions.

## Definition 2.1

Let $G: X \rightarrow Y$ be a locally Lipschitz and directionally differentiable function. Then, the function $G$ is said to be semismooth at $x$ if

$$
G(x+\Delta x)-G(x)-V \Delta x=o(\|\Delta x\|), \forall V \in J_{C} G(x+\Delta x)
$$

Moreover, $G$ is said to be $\rho$ th-order semismooth function at $x$ for some $\rho \in(0,1]$ if

$$
G(x+\Delta x)-G(x)-V \Delta x=O\left(\|\Delta x\|^{1+\rho}\right), \forall V \in J_{C} G(x+\Delta x)
$$

In particular, if $\rho=1$, we say $G$ is strongly semismooth at $x$. We also say $G: X \rightarrow Y$ is a semismooth (resp. $\rho$ th-order semismooth, strongly semismooth) function if $G$ is semismooth (resp. $\rho$ th-order semismooth, strongly semismooth) at $x$ for all $x \in X$.

The concept of a semismooth function was originally given by Mifflin [16] when $Y=\mathbb{R}$. Later on, Qi and Sun [17] (see also [48]) extended the definition to vector value functions and showed that semismooth functions play an important role in establishing the local convergence rate of the generalized Newton method for solving nonsmooth equations. From the definitions of the semismooth functions, it is clear that scalar multiplication and sums of semismooth (resp. $\rho$ th-order semismooth) functions are still semismooth (resp. $\rho$ th-order semismooth) functions. An important example of the strongly semismooth function is the eigenvalue function of a symmetric matrix [18]. The next result [18, Theorem 3.7] provides a convenient tool for proving $\rho$ th-order semismoothness.

Lemma 2.1
Suppose that $G: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m}$ is locally Lipschitzian and directionally differentiable in a neighborhood of $x$. Let $p \in(0,1]$. Then, $G$ is $\rho$ th-order semismooth if and only if for any $x+\Delta x \in$ $D_{G}$,

$$
G(x+\Delta x)-G(x)-\nabla G(x+\Delta x) \Delta x=O\left(\|\Delta x\|^{1+\rho}\right)
$$

Next, we recall some basic definitions and facts of tensor and its eigenvalues. Let $n \in \mathbb{N}$ and let $m$ be an even number. Consider

$$
S=\{\mathcal{A}: \mathcal{A} \text { is an } m \text { th-order } n \text {-dimensional symmetric tensor }\} \text {. }
$$

Clearly, $S$ is a vector space under the addition and multiplication defined as follows: for any $t \in \mathbb{R}$, $\mathcal{A}=\left(\mathcal{A}_{i_{1}, \ldots, i_{m}}\right)_{1 \leqslant i_{1}, \ldots, i_{m} \leqslant n}$ and $\mathcal{B}=\left(\mathcal{B}_{i_{1}, \ldots, i_{m}}\right)_{1 \leqslant i_{1}, \ldots, i_{m} \leqslant n}$

$$
\mathcal{A}+\mathcal{B}=\left(\mathcal{A}_{i_{1}, \ldots, i_{m}}+\mathcal{B}_{i_{1}, \ldots, i_{m}}\right)_{1 \leqslant i_{1}, \ldots, i_{m} \leqslant n} \text { and } t \mathcal{A}=\left(t \mathcal{A}_{i_{1}, \ldots, i_{m}}\right)_{1 \leqslant i_{1}, \ldots, i_{m} \leqslant n} .
$$

For each $\mathcal{A}, \mathcal{B} \in S$, we define the inner product by

$$
\langle\mathcal{A}, \mathcal{B}\rangle_{S}=\sum_{i_{1}, \ldots, i_{m}=1}^{n} \mathcal{A}_{i_{1}, \ldots, i_{m}} \mathcal{B}_{i_{1}, \ldots, i_{m}}
$$

The corresponding norm is defined by $\|\mathcal{A}\|_{S}=\left(\langle\mathcal{A}, \mathcal{A}\rangle_{S}\right)^{1 / 2}=\left(\sum_{i_{1}, \ldots, i_{m}=1}^{n} \mathcal{A}_{i_{1}, \ldots, i_{m}}^{2}\right)^{1 / 2}$. The unit ball in $S$ is denoted by $\mathbb{B}_{S}$. For a vector $x \in \mathbb{R}^{n}$, we use $x_{i}$ to denotes its $i$ th component. We use $x^{[m-1]}$ to denote a vector in $\mathbb{R}^{n}$ such that $x_{i}^{[m-1]}=\left(x_{i}\right)^{m-1}$. Moreover, for a vector $x \in \mathbb{R}^{n}$, we use $x^{m}$ to denote the $m$ th-order $n$-dimensional symmetric rank one tensor induced by $x$, that is,

$$
\left(x^{m}\right)_{i_{1} \ldots i_{m}}=x_{i_{1}} \ldots x_{i_{m}}, \forall i_{1}, \ldots, i_{m} \in\{1, \ldots, n\} .
$$

Let $\mathcal{A} \in S$. By the tensor product (cf. [5]), $\mathcal{A} x^{m}$ is a real number defined as

$$
\mathcal{A} x^{m}=\sum_{i_{1}, \ldots, i_{m}=1}^{n} \mathcal{A}_{i_{1}, \ldots, i_{m}} x_{i_{1}} \ldots x_{i_{m}}=\left\langle\mathcal{A}, x^{m}\right\rangle_{S}
$$

and $\mathcal{A} x^{m-1}$ is a vector in $\mathbb{R}^{n}$ whose $i$ th component is

$$
\begin{equation*}
\sum_{i_{2} \cdots i_{m}=1}^{n} A_{i i_{2} \cdots i_{m}} x_{i_{2}} \ldots x_{i_{m}} \tag{2.2}
\end{equation*}
$$

## Definition 2.2

Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional real symmetric tensor. We say $\lambda \in \mathbb{R}$ is a $Z$-eigenvalue of $\mathcal{A}$ and $x \neq 0, x \in \mathbb{R}^{n}$ is a $Z$-eigenvector corresponding to $\lambda$ if $(x, \lambda)$ satisfies

$$
\left\{\begin{array}{c}
\mathcal{A} x^{m-1}=\lambda x \\
x^{T} x=1
\end{array}\right.
$$

Moreover, $\lambda \in \mathbb{R}$ is an $H$-eigenvalue of $\mathcal{A}$, and $x \neq 0, x \in \mathbb{R}^{n}$ is an $H$-eigenvector corresponding to $\lambda$ if $(x, \lambda)$ satisfies

$$
\mathcal{A} x^{m-1}=\lambda x^{[m-1]}
$$

This definition of $Z$-eigenvalues and $H$-eigenvalues was introduced by Qi in [13]. Independently, Lim [11] also gave the definitions via a variational approach and established an interesting PerronFrobenius theorem for tensors with nonnegative entries. From [13] and [19], both $Z$-eigenvalues and $H$-eigenvalues for an even order symmetric tensor always exist. Moreover, from the definitions, we can see that finding an $H$-eigenvalue of a symmetric tensor is equivalent to solving a homogeneous polynomial equation, whereas calculating a $Z$-eigenvalue is equivalent to solving nonhomogeneous polynomial equations. In general, the behaviors of $Z$-eigenvalues and $H$-eigenvalues can be quite different. For example, a diagonal symmetric tensor A has exactly $n$ many $H$-eigenvalues and may have more than $n$ many Z-eigenvalues (for more details, see [13]). Recently, by reducing a symmetric tensor to a pseudo-canonical form, Qi et al. [20] proposed a direct method for finding all the Z-eigenvalues in the case of order three and dimension three. More recently, Kolda and Mayo [21] provided a shifted power method for computing a $Z$-eigenvalue and its associated eigenvector for a symmetric tensor. For numerical methods of finding $H$-eigenvalues for tensors with nonnegative entries, see [12]. Recently, a new method for finding the maximum $Z$-eigenvalue of a weakly nonnegative symmetric tensor using sum-of-squares programming problem is also proposed in [22].

## 3. THE MAXIMUM $Z$-EIGENVALUE FUNCTION

In this section, we examine the continuity and differentiability of the maximum $Z$-eigenvalue function. To do this, we first formally define the maximum $Z$-eigenvalue function. Because any real symmetric tensor with even order always has a $Z$-eigenvalue (cf. [13, 19]), it then makes sense to define the maximum $Z$-eigenvalue function $\lambda_{1}^{Z}: S \rightarrow \mathbb{R}$ as follows:

$$
\lambda_{1}^{Z}(\mathcal{A})=\{\lambda \in \mathbb{R}: \lambda \text { is the largest Z-eigenvalue of } \mathcal{A}\}
$$

We first recall the following simple lemma, which will be useful for our later analysis. Its proof can be found in $[11,13]$. However, for the completeness of the paper, we present the proof here.

## Lemma 3.1

Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional real symmetric tensor where $m$ is even. Then, we have

$$
\lambda_{1}^{Z}(\mathcal{A})=\max _{\|x\|=1} \mathcal{A} x^{m}
$$

## Proof

Consider the following optimization problem ( $P$ )

$$
\begin{array}{rl}
(P) \max _{x \in \mathbb{R}^{n}} & \mathcal{A} x^{m} \\
\text { s.t. } & \|x\|^{m}=1
\end{array}
$$

Let $f(x):=\mathcal{A} x^{m}$ and $g(x):=\left(x^{T} x\right)^{\frac{m}{2}}=\|x\|^{m}$. Because $f$ is continuous and the feasible set $\{x: g(x)=1\}$ is compact, a global maximizer of $(P)$ exists. Denote a maximizer of $(P)$ by $x_{0}$. Clearly, $x_{0} \neq 0$. Note that $g$ is a homogeneous polynomial with degree $m$. The Euler identity
implies that $\nabla g(x)^{T} x=m g(x)$. Thus, for any $x$ with $g(x)=1, \nabla g(x) \neq 0$. So, the standard Karush-Kuhn-Tucker (KKT) theory implies that there exists $\lambda_{0} \in \mathbb{R}$ such that

$$
m \mathcal{A} x_{0}^{m-1}-m \lambda_{0}\left\|x_{0}\right\|^{m-2} x_{0}=\nabla f\left(x_{0}\right)-\lambda_{0} \nabla g\left(x_{0}\right)=0 .
$$

This implies that $\lambda_{0}$ is a real eigenvalue of $\mathcal{A}$ and so $\lambda_{0} \leqslant \lambda_{1}^{Z}(\mathcal{A})$. Note that $v(P)=\mathcal{A} x_{0}^{m}=$ $x_{0}^{T}\left(\mathcal{A} x_{0}^{m-1}\right)=x_{0}^{T}\left(\lambda_{0} x_{0}\right)=\lambda_{0}$, where $v(P)$ is the optimal value of $(P)$. It follows that $v(P) \leqslant \lambda_{1}^{Z}(\mathcal{A})$, that is, $\max _{\|x\|_{m}=1} \mathcal{A} x^{m} \leqslant \lambda_{1}^{Z}(\mathcal{A})$. Finally, noting that, for any eigenvector $u$ corresponds to $\lambda_{1}^{Z}(\mathcal{A})$ with $\|u\|=1$, we have

$$
\mathcal{A} u^{m}=u^{T}\left(\mathcal{A} u^{m-1}\right)=\lambda_{1}^{Z}(\mathcal{A}) u^{T} u=\lambda_{1}^{Z}(\mathcal{A})\|u\|^{2}=\lambda_{1}^{Z}(\mathcal{A})
$$

Thus, $\lambda_{1}^{Z}(\mathcal{A})=\max _{\|x\|=1} \mathcal{A} x^{m}$, and so, the conclusion follows .

## Remark 3.1

Define the normalized eigenspace associated with $\lambda_{1}^{Z}(\mathcal{A})$ by $E_{1}^{Z}(\mathcal{A}):=\left\{u: \mathcal{A} u^{m-1}=\right.$ $\left.\lambda_{1}^{Z}(\mathcal{A}) u,\|u\|=1\right\}$. From the proof of the Lemma 3.1, we see that

$$
E_{1}^{Z}(\mathcal{A})=\left\{u: \mathcal{A} u^{m}=\lambda_{1}^{Z}(\mathcal{A}),\|u\|=1\right\}
$$

Next, we show that the maximum $Z$-eigenvalue function is continuous and convex.
Theorem 3.1
The function $\lambda_{1}^{Z}$ is a continuous and convex function on $S$. Moreover, its conjugate $\left(\lambda_{1}^{Z}\right)^{*}$ can be calculated as

$$
\left(\lambda_{1}^{Z}\right)^{*}(\mathcal{B})=\left\{\begin{array}{cl}
0, & \text { if } \quad \mathcal{B} \in \operatorname{conv}\left\{u^{m}:\|u\|=1\right\} \\
+\infty, & \text { else },
\end{array}\right.
$$

where conv $A$ denotes the convex hull of $A$ and is defined by

$$
\operatorname{conv} A=\left\{\sum_{i=1}^{s} \mu_{i} a_{i}: \mu_{i} \geqslant 0, \sum_{i=1}^{s} \mu_{i}=1, a_{i} \in A, s \in \mathbb{N}\right\}
$$

Proof
Let $T:=\left\{u^{m}:\|u\|=1\right\} \subseteq S$. Because $\lambda_{1}^{Z}(\mathcal{A})=\max _{\|x\|=1} \mathcal{A} x^{m}$, we have $\lambda_{1}^{Z}(\mathcal{A})=$ $\max _{\mathcal{B} \in T}\langle\mathcal{B}, \mathcal{A}\rangle_{S}$. Note that $\mathcal{B} \mapsto\langle\mathcal{B}, \mathcal{A}\rangle_{S}$ is affine and the supremum of a series of affine functions is convex. It follows that $\lambda_{1}^{Z}$ is a finite-valued convex function on $S$ and so is continuous and convex. From the definition, it can be verified that $\max _{\mathcal{B} \in T}\langle\mathcal{B}, \mathcal{A}\rangle_{S}=\max _{\mathcal{B} \in \operatorname{conv} T}\langle\mathcal{B}, \mathcal{A}\rangle_{S}$, and so,

$$
\begin{equation*}
\lambda_{1}^{Z}(\mathcal{A})=\max _{\mathcal{B} \in \operatorname{conv} T}\langle\mathcal{B}, \mathcal{A}\rangle_{S} \tag{3.3}
\end{equation*}
$$

It follows that, for each $\mathcal{B} \in S$,

$$
\begin{aligned}
\left(\lambda_{1}^{Z}\right)^{*}(\mathcal{B})=\sup _{\mathcal{A} \in S}\left\{\langle\mathcal{B}, \mathcal{A}\rangle_{S}-\lambda_{1}^{Z}(\mathcal{A})\right\} & =\sup _{\mathcal{A} \in S}\left\{\langle\mathcal{B}, \mathcal{A}\rangle_{S}-\lambda_{1}^{Z}(\mathcal{A})\right\} \\
& =\sup _{\mathcal{A} \in S}\left\{\langle\mathcal{B}, \mathcal{A}\rangle_{S}-\sup _{\mathcal{C} \in \operatorname{conv} T}\langle\mathcal{C}, \mathcal{A}\rangle\right\} \\
& =\sup _{\mathcal{A} \in S} \min _{\mathcal{C} \in \operatorname{conv} T}\left\{\langle\mathcal{B}, \mathcal{A}\rangle_{S}-\langle\mathcal{C}, \mathcal{A}\rangle_{S}\right\} \\
& =\min _{\mathcal{C} \in \operatorname{conv} T} \sup _{\mathcal{A} \in S}\left\{\langle\mathcal{B}, \mathcal{A}\rangle_{S}-\langle\mathcal{C}, \mathcal{A}\rangle_{S}\right\}
\end{aligned}
$$

where the last equality follows from the standard convex-concave minimax theorem (cf. [15, Theorem 2.10.2]). Note that for each $\mathcal{C} \in \operatorname{conv} T$,

$$
\sup _{\mathcal{A} \in S}\left\{\langle\mathcal{B}, \mathcal{A}\rangle_{S}-\langle\mathcal{C}, \mathcal{A}\rangle_{S}\right\}=\left\{\begin{array}{cl}
0, & \text { if } \quad \mathcal{B}=\mathcal{C} \\
+\infty, & \text { else }
\end{array}\right.
$$

Thus, the conclusion follows.
As $\lambda_{1}^{Z}$ is continuous and convex, its convex $\epsilon$-subdifferential $(\epsilon \geqslant 0)$ always exists where the convex $\epsilon$-subdifferential (cf. [23]) $\partial_{\epsilon} \lambda_{1}^{Z}$ is defined by

$$
\partial_{\epsilon} \lambda_{1}^{Z}(\mathcal{A})=\left\{\mathcal{B} \in S:\left\langle\mathcal{B}, \mathcal{A}^{\prime}-A\right\rangle_{S} \leqslant \lambda_{1}^{Z}\left(\mathcal{A}^{\prime}\right)-\lambda_{1}^{Z}(\mathcal{A})+\epsilon \text { for all } \mathcal{A}^{\prime} \in S\right\}
$$

We are now ready to state the formula for the $\epsilon$-subdifferential of maximum $Z$-eigenvalue function.

Theorem 3.2
Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor where $m$ is even. Then, for all $\epsilon \geqslant 0$, we have

$$
\begin{gathered}
\partial_{\epsilon} \lambda_{1}^{Z}(\mathcal{A})=\left\{B \in S: B=\sum_{i=1}^{s} \mu_{i} u_{i}^{m}, \mu_{i} \geqslant 0, \sum_{i=1}^{s} \mu_{i}=1,\left\|u_{i}\right\|=1, s \in \mathbb{N}\right. \\
\left.\quad \text { and } \lambda_{1}^{Z}(\mathcal{A}) \leqslant \sum_{i=1}^{s} \mu_{i}\left(\mathcal{A} u_{i}^{m}\right)+\epsilon\right\}
\end{gathered}
$$

Proof
From the generalized Fenchel inequality, we have

$$
\mathcal{B} \in \partial_{\epsilon} \lambda_{1}^{Z}(\mathcal{A}) \Leftrightarrow \lambda_{1}^{Z}(\mathcal{A})+\left(\lambda_{1}^{Z}\right)^{*}(\mathcal{B}) \leqslant\langle\mathcal{A}, \mathcal{B}\rangle_{S}+\epsilon
$$

Note that

$$
\left(\lambda_{1}^{Z}\right)^{*}(\mathcal{B})=\left\{\begin{array}{cl}
0, & \text { if } \quad \mathcal{B} \in \operatorname{conv}\left\{u^{m}:\|u\|=1\right\} \\
+\infty, & \text { else. }
\end{array}\right.
$$

So, $\mathcal{B} \in \partial_{\epsilon} \lambda_{1}^{Z}(\mathcal{A})$ if and only if $\mathcal{B} \in \operatorname{conv}\left\{u^{m}:\|u\|=1\right\}$ and

$$
\lambda_{1}^{Z}(\mathcal{A}) \leqslant\langle\mathcal{A}, \mathcal{B}\rangle_{S}+\epsilon
$$

This is further equivalent to the fact that there exist $s \in \mathbb{N}, \mu_{i} \geqslant 0, i=1, \ldots, s$ with $\sum_{i=1}^{s} \mu_{i}=1$ and $\left\|u_{i}\right\|=1$ such that

$$
\mathcal{B}=\sum_{i=1}^{s} \mu_{i} u_{i}^{m} \text { and } \lambda_{1}^{Z}(\mathcal{A}) \leqslant\left\langle\mathcal{A}, \sum_{i=1}^{s} \mu_{i} u_{i}^{m}\right\rangle_{S}+\epsilon=\sum_{i=1}^{s} \mu_{i}\left(\mathcal{A} u_{i}^{m}\right)+\epsilon
$$

Thus, the conclusion follows.
When $\epsilon=0$, the convex subdifferential formula can be simplified as follows.

## Corollary 3.1

Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor where $m$ is even. Then, we have

$$
\partial \lambda_{1}^{Z}(\mathcal{A})=\operatorname{conv}\left\{u^{m}: u \in E_{1}^{Z}(\mathcal{A})\right\}
$$

Proof
Let $\epsilon=0$. Then, the preceding theorem shows that

$$
\begin{aligned}
\partial \lambda_{1}^{Z}(\mathcal{A})= & \left\{B \in S: B=\sum_{i=1}^{s} \mu_{i} u_{i}^{m}, \mu_{i} \geqslant 0, \sum_{i=1}^{s} \mu_{i}=1,\left\|u_{i}\right\|=1, s \in \mathbb{N}\right. \\
& \text { and } \left.\lambda_{1}^{Z}(\mathcal{A}) \leqslant \sum_{i=1}^{s} \mu_{i}\left(\mathcal{A} u_{i}^{m}\right)\right\}
\end{aligned}
$$

Note that $\lambda_{1}^{Z}(\mathcal{A})=\max _{\|x\|=1} \mathcal{A} x^{m}$. So, $\lambda_{1}^{Z}(\mathcal{A}) \leqslant \sum_{i=1}^{s} \mu_{i}\left(\mathcal{A} u_{i}^{m}\right)$ with $\mu_{i} \geqslant 0, \sum_{i=1}^{s} \mu_{i}=1$ and $\left\|u_{i}\right\|=1$ is equivalent to $\lambda_{1}^{Z}(\mathcal{A})=\mathcal{A} u_{i}^{m}$, for all $i=1, \ldots, s$. It follows that

$$
\begin{aligned}
\partial \lambda_{1}^{Z}(\mathcal{A})= & \left\{B \in S: B=\sum_{i=1}^{s} \mu_{i} u_{i}^{m}, \mu_{i} \geqslant 0, \sum_{i=1}^{s} \mu_{i}=1,\left\|u_{i}\right\|=1, s \in \mathbb{N}\right. \\
& \text { and } \left.\lambda_{1}^{Z}(\mathcal{A})=\mathcal{A} u_{i}^{m}\right\}
\end{aligned}
$$

Therefore, the conclusion follows Remark 3.1.

## Remark 3.2

If $m=2$, our subdifferential formula for $\lambda_{1}^{Z}$ reduces to

$$
\partial \lambda_{1}^{Z}(\mathcal{A})=\operatorname{conv}\left\{u u^{T}:\left(\lambda_{1}^{Z}(\mathcal{A}), u\right) \text { is an eigenpair of } \mathcal{A} \text { and }\|u\|=1\right\}
$$

which is the classical subdifferential formula of the maximum eigenvalue function in the matrix case (cf. [24]).

Definition 3.1
Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor where $m$ is even. Recall that the normalized eigenspace associated with $\lambda_{1}^{Z}(\mathcal{A})$ is given by $E_{1}^{Z}(\mathcal{A}):=\left\{u: \mathcal{A} u^{m-1}=\lambda_{1}^{Z}(\mathcal{A}) u,\|u\|=1\right\}$. Then, the eigenspace associated with $\lambda_{1}^{Z}(\mathcal{A})$ is $\operatorname{span} E_{1}^{Z}(\mathcal{A})$, where $\operatorname{span} E_{1}^{Z}(\mathcal{A})$ is the subspace generated by $E_{1}^{Z}(\mathcal{A})$, that is, $\operatorname{span} E_{1}^{Z}(\mathcal{A}):=\bigcup_{t \in \mathbb{R}}\left\{t \operatorname{conv} E_{1}^{Z}(\mathcal{A})\right\}$. We now define the geometric multiplicity of $\lambda_{1}^{Z}(\mathcal{A})$ as the dimension of the subspace $\operatorname{span} E_{1}^{Z}(\mathcal{A})$.

## Corollary 3.2

Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor where $m$ is even. Then, the maximum $Z$-eigenvalue function $\lambda_{1}^{Z}$ is locally Lipschitz and is (Fréchet) differentiable almost everywhere. Moreover, $\lambda_{1}^{Z}$ is differentiable at $\mathcal{A} \in S$ if and only if the geometric multiplicity of $\lambda_{1}^{Z}(\mathcal{A})$ is one.

## Proof

From the preceding theorem, the maximum $Z$-eigenvalue function $\lambda_{1}^{Z}$ is continuous and convex. So, $\lambda_{1}^{Z}$ is locally Lipschitz. Then, the Radamecher theorem implies that it is (Fréchet) differentiable almost everywhere. To see the last assertion, as $m$ is even, we see that $u^{m}=(-u)^{m}$. So, the geometric multiplicity of $\lambda_{1}^{Z}(\mathcal{A})$ is one is equivalent to the fact that the set

$$
\partial \lambda_{1}^{Z}(\mathcal{A})=\left\{u^{m}:\left(\lambda_{1}^{Z}(\mathcal{A}), u\right) \text { is an eigenpair of } \mathcal{A} \text { and }\|u\|=1\right\}
$$

is a singleton. Note that a continuous convex function on a finite dimensional space is Fréchet differentiable if and only if its subdifferential is a singleton. Thus, the conclusion follows.

### 3.1. Perturbation bound

Consider $\mathcal{A}(v)=\mathcal{A}+\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}$, where $v=\left(v_{1}, \ldots, v_{r}\right) \in \mathbb{R}^{r}$ and $\mathcal{A}, \mathcal{B}_{j} \in S, j=1, \ldots, r$. Define the map $h: \mathbb{R}^{r} \rightarrow \mathbb{R}$ by

$$
h(v)=\lambda_{1}^{Z}(\mathcal{A}(v))
$$

Then, we see that $h$ is a continuous and convex function on $\mathbb{R}^{r}$. In the next discussion, we present the following sensitivity result of the maximum $Z$-eigenvalue function. In the matrix case, this result collapses to the classical sensitivity result derived in [25].

## Proposition 3.1

Let $\mathcal{A}, \mathcal{B}_{j} \in S, j=1, \ldots, r$. Consider the map $h: \mathbb{R}^{r} \rightarrow \mathbb{R}$ defined by

$$
h(v)=\lambda_{1}^{Z}\left(\mathcal{A}+\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}\right)
$$

where $v=\left(v_{1}, \ldots, v_{r}\right) \in \mathbb{R}^{r}$. Then, we have

$$
\sup _{u \in E_{1}^{Z}(\mathcal{A})} \sum_{j=1}^{r} v_{j}\left\langle\mathcal{B}_{j}, u^{m}\right\rangle_{S} \leqslant h(v)-h(0) \leqslant \lambda_{1}^{Z}\left(\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}\right)
$$

Proof
First of all, from (3.3), we have

$$
\begin{aligned}
h(v)-h(0) & =\max _{\mathcal{C} \in \operatorname{conv} T}\left\langle\mathcal{C}, \mathcal{A}+\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}\right\rangle_{S}-\max _{\mathcal{C} \in \operatorname{conv} T}\langle\mathcal{C}, \mathcal{A}\rangle_{S} \\
& \leqslant \max _{\mathcal{C} \in \operatorname{conv} T}\left\langle\mathcal{C}, \sum_{j=1}^{r} v_{j} \mathcal{B}_{j}\right\rangle_{S}=\lambda_{1}^{Z}\left(\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}\right) .
\end{aligned}
$$

On the other hand, from the convexity of $h$, we have

$$
h(v)-h(0) \geqslant \sup _{a \in \partial h(0)} a^{T} v
$$

Note from the chain rule of the convex subdifferential (cf. [15]) that

$$
a \in \partial h(0)=\left\{\left(\begin{array}{c}
\left\langle\mathcal{B}_{1}, \mathcal{D}\right\rangle_{S} \\
\vdots \\
\left\langle\mathcal{B}_{r}, \mathcal{D}\right\rangle_{S}
\end{array}\right): \mathcal{D} \in \partial \lambda_{1}^{Z}(\mathcal{A})\right\} .
$$

So,

$$
\begin{aligned}
& \sup _{a \in \partial h(0)} a^{T} v=\sup _{\mathcal{D} \in \partial \lambda_{1}^{Z}(\mathcal{A})} \sum_{j=1}^{r}\left\langle\mathcal{B}_{j}, \mathcal{D}\right\rangle_{S} v_{j}=\sup _{\substack{\mu_{i} \geqslant 0, \sum_{i=1}^{S} \mu_{i}=1, s \in \mathbb{N},\left\|u_{i}\right\|=1, \lambda_{1}^{Z}(\mathcal{A})=\mathcal{A} u_{i}^{m}}} \sum_{j=1}^{r}\left\langle\mathcal{B}_{j}, \sum_{i=1}^{s} \mu_{i} u_{i}^{m}\right\rangle_{S} v_{j} \\
& =\sup _{\substack{u_{i} \geqslant 0, \sum_{i=1}^{s} \mu_{i}=1, s \in \mathbb{N},\left\|u_{i}\right\|=1, \lambda_{1}^{Z}(\mathcal{A})=\mathcal{A} u_{i}^{m}}} \sum_{i=1}^{s} \mu_{i}\left\langle\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}, u_{i}^{m}\right\rangle_{S} \\
& =\sup _{\substack{\left.\mu_{i} \geqslant 0, \sum_{i=1}^{s} \\
\left\|u_{i}\right\|=1, \lambda_{1}^{Z} \\
\mu_{i}=1, s \in \mathbb{A}\right)=\mathcal{A} u_{i}^{m}}}\left\langle\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}, \sum_{i=1}^{s} \mu_{i} u_{i}^{m}\right\rangle_{S} \\
& \left\|u_{i}\right\|=1, \lambda_{1}^{Z}(\mathcal{A})=\mathcal{A} u_{i}^{m} \\
& =\sup _{\|u\|=1, \lambda_{1}^{Z}(\mathcal{A})=\mathcal{A} u^{m}}\left\langle\sum_{j=1}^{r} v_{j} \mathcal{B}_{j}, u^{m}\right\rangle_{S} \\
& =\sup _{u \in E_{1}^{Z}(\mathcal{A})} \sum_{j=1}^{r} v_{j}\left\langle\mathcal{B}_{j}, u^{m}\right\rangle_{S},
\end{aligned}
$$

where the last equality follows from Remark 3.1.

## 4. STABILITY ANALYSIS OF THE NORMALIZED EIGENSPACE

In this section, we study the stability of the normalized eigenspace $E_{1}^{Z}(\mathcal{A})$, that is, how the normalized eigenspace $E_{1}^{Z}(\mathcal{A})$ changes when the corresponding symmetric tensor $\mathcal{A}$ perturbs. To achieve the Hölder stability, we need the following two results. The first result gives an effective estimate for the growth rate (Łojasiewicz exponent) of a polynomial with real coefficients (for sharper exponent under specific conditions, see [26, Theorem 2.3], [27, Lemma 4.3] and [45]). Then, the second result
is a local error bound result that estimates how far a point to the lower level set $S_{f}=\{x: f(x) \leqslant 0\}$ is, in terms of its function value (for related error bound result, see [28,29]).

Lemma 4.1 (cf. [30, Theorem 4.2])
Let $f$ be a polynomial with real coefficients on $\mathbb{R}^{n}$ with degree $m \geqslant 2$. Suppose that $f(0)=0$ and $\nabla f(0)=0$. Then, there exist $\epsilon, r, c>0$ such that

$$
\|\nabla f(x)\| \geqslant c|f(x)|^{\tau} \text { for all }\|x\| \leqslant \epsilon \text { with } f(x) \leqslant r
$$

where $\tau \leqslant 1-\left(m(3 m-3)^{n-1}\right)^{-1}$.
Lemma 4.2 (cf. [31, Corollary 2.1])
Let $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ be a continuously differentiable function, and let $\bar{x} \in \mathbb{R}^{n}$ such that $\bar{x} \in$ bdry $S_{f}$, where $S_{f}=\{x: f(x) \leqslant 0\}$ and bdry $C$ denotes the boundary of the set $C$. Suppose that there exist $\epsilon, c>0$ such that $\|\nabla f(x)\| f(x)^{\theta-1} \geqslant c$ for all $x$ with $\|x-\bar{x}\| \leqslant \epsilon$ and $x \notin S_{f}$. Then,

$$
d\left(x, S_{f}\right) \leqslant \frac{1}{c}(\max \{f(x), 0\})^{\theta} \text { for all } x \text { with }\|x-\bar{x}\| \leqslant \frac{\epsilon}{2} .
$$

Recall that, for an $(n \times n)$ matrix $M, M \prec 0$ (resp. $M \preceq 0$ ) means that $M$ is negative definite (resp. negative semidefinite). Moreover, we use $\mathbf{I}_{\mathbf{n}}$ to denote the ( $n \times n$ ) identity matrix. We now provide the stability result of the normalized eigenspace. We achieve this by noting that the normalized eigenspace $E_{1}^{Z}(\mathcal{A})$ is just the optimal solution of the constraint polynomial optimization problem

$$
\left(P_{\mathcal{A}}\right) \max \left\{\mathcal{A} x^{m}:\|x\|=1\right\}
$$

and the stability of the optimal solution of the parameterized optimization problem $\left(P_{\mathcal{A}}\right)$ can be approached by examining the growth property of a related real polynomial, which can be regarded as a generalized Lagrangian function of the constraint optimization problem $\left(P_{\mathcal{A}}\right)$.

Theorem 4.1
Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor where $m$ is even and $m \geqslant 4$.
(i) Then, the normalized eigenspace is Hölder stable at $\mathcal{A}$ with the order $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$, that is, there exist $\epsilon>0$ and $\alpha>0$ such that for any tensor $\mathcal{B}$ with $\|\mathcal{B}-\mathcal{A}\|_{S} \leqslant \epsilon$,

$$
\begin{equation*}
E_{1}^{Z}(\mathcal{B}) \subseteq E_{1}^{Z}(\mathcal{A})+\alpha\left(\|\mathcal{B}-\mathcal{A}\|_{S}\right)^{\frac{1}{d}} \mathbb{B}_{\mathbb{R}^{n}} \tag{4.4}
\end{equation*}
$$

where $d \in \mathbb{N}$ with $d \leqslant m(3 m-3)^{n-1}-1$, and $\mathbb{B}_{\mathbb{R}^{n}}$ is the unit ball in $\mathbb{R}^{n}$.
(ii) If we further assume that the following second-order condition holds: $\forall u \in E_{1}^{Z}(\mathcal{A})$

$$
\begin{equation*}
(m-1) \mathcal{A} u^{m-2}-\lambda_{1}^{Z}(\mathcal{A})\left((m-2) u u^{T}+\mathbf{I}_{\mathbf{n}}\right) \prec 0 \text { on } C_{u}=\left\{h \in \mathbb{R}^{n}: h^{T} u=0\right\}, \tag{4.5}
\end{equation*}
$$

then the integer $d$ in (4.4) can be set as 1 , that is, there exist $\epsilon>0$ and $\alpha>0$ such that for any tensor $\mathcal{B}$ with $\|\mathcal{B}-\mathcal{A}\|_{S} \leqslant \epsilon$,

$$
E_{1}^{Z}(\mathcal{B}) \subseteq E_{1}^{Z}(\mathcal{A})+\alpha\|\mathcal{B}-\mathcal{A}\|_{S} \mathbb{B}_{\mathbb{R}^{n}}
$$

Proof of ( $i$ )
Fix any $u \in E_{1}^{Z}(\mathcal{A})$. Let $\gamma>0$, and let

$$
f(x):=\lambda_{1}^{Z}(\mathcal{A})\|x+u\|^{m}-\mathcal{A}(x+u)^{m}+\left(\|x+u\|^{2}-1\right)^{2} .
$$

It can be seen that $f$ is a real polynomial on $\mathbb{R}^{n}$ with degree $m$. Moreover, it can be verified that $f$ is nonnegative, $f(0)=0$ and $\nabla f(0)=0$ (as 0 is the minimizer of $f$ ). So, Lemma 4.1 implies that there exist $\gamma_{u}, r_{u}, c_{u}>0$ such that

$$
\|\nabla f(x)\| \geqslant c_{u}|f(x)|^{\tau}=f(x)^{\tau} \text { for all }\|x\| \leqslant \gamma_{u} \text { with } f(x) \leqslant r_{u},
$$

where $\tau \leqslant 1-\left(m(3 m-3)^{n-1}\right)^{-1}$. Choose $\delta_{u}<\gamma_{u} / 2$ such that $f(x) \leqslant r_{u}$ for all $\|x\| \leqslant \delta_{u}$ (this can be carried out as $f(0)=0$ and $f$ is continuous). Then, we see that

$$
\|\nabla f(x)\||f(x)|^{(1-\tau)-1}=\|\nabla f(x)\||f(x)|^{-\tau} \geqslant c_{u} \text { for all }\|x\| \leqslant 2 \delta_{u}
$$

Letting $h(x):=f(x-u)=\lambda_{1}^{Z}(\mathcal{A})\|x\|^{m}-\mathcal{A}(x)^{m}+\left(\|x\|^{2}-1\right)^{2}$, we obtain that

$$
\|\nabla h(x)\||h(x)|^{(1-\tau)-1} \geqslant c_{u} \text { for all }\|x-u\| \leqslant 2 \delta_{u}
$$

As $E_{1}^{Z}(\mathcal{A}) \subseteq\{x:\|x\|=1\}, E_{1}^{Z}(\mathcal{A})$ has no interior point and hence $u \in \operatorname{bdry} E_{1}^{Z}(\mathcal{A})$. Note that $h$ is nonnegative and $\{x: h(x) \leqslant 0\}=\{x: h(x)=0\}=E_{1}^{Z}(\mathcal{A})$. Applying Lemma 4.2 with $f=h$ and $\bar{x}=u$, for each $u \in E_{1}^{Z}(\mathcal{A})$, we can find $c_{u}>0$ such that

$$
\begin{equation*}
d\left(x, E_{1}^{Z}(\mathcal{A})\right) \leqslant \frac{1}{c_{u}} h(x)^{\theta} \text { for all } x \text { with }\|x-u\| \leqslant \delta_{u}, \tag{4.6}
\end{equation*}
$$

where $\theta=1-\tau \geqslant\left(m(3 m-3)^{n-1}\right)^{-1}$. Now, using standard compactness argument, we show that there exist $c>0$ and $\delta>0$ such that

$$
\begin{equation*}
d\left(x, E_{1}^{Z}(\mathcal{A})\right) \leqslant \frac{1}{c} h(x)^{\theta} \text { for all } x \text { with } d\left(x, E_{1}^{Z}(\mathcal{A})\right) \leqslant \delta \tag{4.7}
\end{equation*}
$$

As $E_{1}^{Z}(\mathcal{A})$ is compact and

$$
E_{1}^{Z}(\mathcal{A}) \subseteq \bigcup_{u \in E_{1}^{Z}(\mathcal{A})}\left\{x:\|x-u\| \leqslant \frac{\delta_{u}}{2}\right\}
$$

there exist $l \in \mathbb{N}$ and $\left\{u_{1}, \ldots, u_{l}\right\} \subseteq E_{1}^{Z}(\mathcal{A})$ such that

$$
\begin{equation*}
E_{1}^{Z}(\mathcal{A}) \subseteq \bigcup_{i=1}^{l}\left\{x:\left\|x-u_{i}\right\| \leqslant \frac{\delta_{u_{i}}}{2}\right\} \tag{4.8}
\end{equation*}
$$

Take $\delta=\frac{1}{2} \min \left\{\delta_{u_{1}}, \ldots, \delta_{u_{l}}\right\}>0$ and $c=\min \left\{c_{u_{1}}, \ldots, c_{u_{l}}\right\}>0$. Then, for each $x$ with $d\left(x, E_{1}^{Z}(\mathcal{A})\right) \leqslant \delta$, we can find $a \in E_{1}^{Z}(\mathcal{A})$ such that $\|x-a\| \leqslant \delta$. By (4.8), there exists $i_{0} \in\{1, \ldots, l\}$ such that $\left\|a-u_{i_{0}}\right\| \leqslant \frac{\delta_{u_{i_{0}}}}{2}$. So,

$$
\left\|x-u_{i_{0}}\right\| \leqslant\|x-a\|+\left\|a-u_{i_{0}}\right\| \leqslant \delta+\frac{\delta_{u_{i_{0}}}}{2} \leqslant \delta_{u_{i_{0}}}
$$

Then, (4.6) implies that

$$
d\left(x, E_{1}^{Z}(\mathcal{A})\right) \leqslant \frac{1}{c_{u_{i_{0}}}} h(x)^{\theta} \leqslant \frac{1}{c} h(x)^{\theta}
$$

Thus, (4.7) holds. Letting $\beta:=c^{\frac{1}{\theta}}$, this shows that for any $x$ with $d\left(x, E_{1}^{Z}(\mathcal{A})\right) \leqslant \delta$,

$$
\begin{equation*}
\lambda_{1}^{Z}(\mathcal{A})\|x\|^{m}-\mathcal{A}(x)^{m}+\left(\|x\|^{2}-1\right)^{2}=h(x) \geqslant c^{\frac{1}{\theta}} d\left(x, E_{1}^{Z}(\mathcal{A})\right)^{\frac{1}{\theta}}=\beta d\left(x, E_{1}^{Z}(\mathcal{A})\right)^{\frac{1}{\theta}} \tag{4.9}
\end{equation*}
$$

For the aforementioned $\delta$, there exists $\epsilon>0$ such that for any $\mathcal{B} \in S$ with $\|\mathcal{B}-\mathcal{A}\|_{S} \leqslant \epsilon$,

$$
\begin{equation*}
E_{1}^{Z}(\mathcal{B}) \subseteq E_{1}^{Z}(\mathcal{A})+\delta \mathbb{B}_{\mathbb{R}^{n}} \tag{4.10}
\end{equation*}
$$

(Otherwise, there exists a sequence of symmetric tensors $B_{n}$ with $\left\|\mathcal{B}_{n}-\mathcal{A}\right\|_{S} \rightarrow 0$ and $y_{n} \in S_{\mathcal{B}_{n}}$ such that $d\left(y_{n}, E_{1}^{Z}(\mathcal{A})\right) \geqslant \delta$. As $\left\|y_{n}\right\|=1$, by passing to subsequence, we may assume that $y_{n} \rightarrow y$ for some $y$ with $\|y\|=1$. Then, we have $d\left(y, E_{1}^{Z}(\mathcal{A})\right) \geqslant \delta>0$. Now, as $\lambda_{1}^{Z}$ is continuous, we have $\mathcal{B}_{n}\left(y_{n}\right)^{m}=\lambda_{1}^{Z}\left(\mathcal{B}_{n}\right) \rightarrow \lambda_{1}^{Z}(\mathcal{A})$. So, passing to limit, we see that $\mathcal{A} y^{m}=\lambda_{1}^{Z}(\mathcal{A})$.

Note that $\|y\|=1$. So, $y \in E_{1}^{Z}(\mathcal{A})$, which makes contradiction.) Now, fix any arbitrary $\mathcal{B}$ with $\|\mathcal{B}-\mathcal{A}\|_{S} \leqslant \epsilon$. From (4.10), we have $E_{1}^{Z}(\mathcal{B}) \subseteq E_{1}^{Z}(\mathcal{A})+\delta \mathbb{B}_{\mathbb{R}^{n}}$. Let $r>0$ be a constant such that

$$
\begin{equation*}
\left\|x^{m}-y^{m}\right\|_{S} \leqslant r\|x-y\| \text { for all } x, y \in K:=\{x:\|x\|=1\} \tag{4.11}
\end{equation*}
$$

Let $d=\frac{1}{\theta}-1 \leqslant m(3 m-3)^{n-1}-1$ and $\alpha=\left(\beta^{-1} r\right)^{d}>0$. For any arbitrary $v \in E_{1}^{Z}(\mathcal{B})$, take $u \in E_{1}^{Z}(\mathcal{A})$ be such that $\|v-u\|=d\left(v, E_{1}^{Z}(\mathcal{A})\right)$. To finish the proof, it suffices to show that

$$
\begin{equation*}
\|v-u\| \leqslant \alpha\|\mathcal{B}-\mathcal{A}\|_{S}^{\frac{1}{d}} \tag{4.12}
\end{equation*}
$$

To see this, we first note that $\|v-u\| \leqslant \delta$ as $E_{1}^{Z}(\mathcal{B}) \subseteq E_{1}^{Z}(\mathcal{A})+\delta \mathbb{B}_{\mathbb{R}^{n}}$. So, (4.9) and $\|v\|=1$ (as $v \in E_{1}^{Z}(\mathcal{B})$ ) imply that

$$
\lambda_{1}^{Z}(\mathcal{A})-\mathcal{A} v^{m}=\lambda_{1}^{Z}(\mathcal{A})\|v\|^{m}-\mathcal{A} v^{m}+\left(\|v\|^{2}-1\right)^{2} \geqslant \beta\|v-u\|^{\frac{1}{\theta}}
$$

As $u$ is an optimal solution of $P_{\mathcal{A}}$, we have $\|u\|=1$ and $\mathcal{A} u^{m}=\lambda_{1}^{Z}(\mathcal{A})$, and so,

$$
\begin{equation*}
\|v-u\|^{\frac{1}{\theta}} \leqslant \beta^{-1}\left(\lambda_{1}^{Z}(\mathcal{A})-\mathcal{A} v^{m}\right)=\beta^{-1}\left(\mathcal{A} u^{m}-\mathcal{A} v^{m}\right) \tag{4.13}
\end{equation*}
$$

Then, as $v$ is optimal for $\left(P_{\mathcal{B}}\right)$, so $\|v\|=1$ and $\mathcal{B} u^{m} \leqslant \mathcal{B} v^{m}$. It follows from (4.11) that

$$
\begin{aligned}
\mathcal{A} u^{m}-\mathcal{A} v^{m} & =\left(\mathcal{B} u^{m}-\mathcal{B} v^{m}\right)+\left((\mathcal{A}-\mathcal{B}) u^{m}-(\mathcal{A}-\mathcal{B}) v^{m}\right) \\
& \leqslant(\mathcal{A}-\mathcal{B}) u^{m}-(\mathcal{A}-\mathcal{B}) v^{m} \\
& \leqslant\left\|u^{m}-v^{m}\right\| S\|\mathcal{A}-\mathcal{B}\|_{S} \\
& \leqslant r\|u-v\|\|\mathcal{B}-\mathcal{A}\|_{S}
\end{aligned}
$$

This together with (4.13) implies that

$$
\|v-u\|^{\frac{1}{\theta}} \leqslant \beta^{-1}\left(\mathcal{A} u^{m}-\mathcal{A} v^{m}\right) \leqslant \beta^{-1} r\|u-v\|\|\mathcal{B}-\mathcal{A}\|_{S}
$$

So, we have $\|v-u\|^{d}=\|v-u\|^{\frac{1}{\theta}-1} \leqslant \beta^{-1} r\|\mathcal{B}-\mathcal{A}\|_{S}$. Note that $\alpha=\left(\beta^{-1} r\right)^{\frac{1}{d}}$. Then,

$$
d\left(v, E_{1}^{Z}(\mathcal{A})\right)=\|v-u\| \leqslant \alpha\|\mathcal{B}-\mathcal{A}\|_{S}^{\frac{1}{d}}
$$

Therefore, (4.12) holds, and so, statement (i) follows.

## Proof of (ii)

Suppose that the second-order condition (4.5) holds. Fix an arbitrary $u \in E_{1}^{Z}(\mathcal{A})$ and consider the minimization problem

$$
\begin{array}{rll}
\left(P_{0}\right) & \min _{x \in \mathbb{R}^{n}} & f(x):=-\mathcal{A} x^{m} \\
\text { s.t. } & \|x\|^{m}=1
\end{array}
$$

Clearly, $u$ satisfies the KKT condition of $\left(P_{0}\right)$ with Lagrange multiplier $\lambda_{1}^{Z}(\mathcal{A})$. Note that the usual second-order sufficient condition for this problem reduces to (4.5). So, the following second-order growth condition holds at $u$ (for example, see [32, Corollary 1]): there exist $\beta>0$ and $\delta>0$ such that

$$
-\mathcal{A} x^{m}+\lambda_{1}^{Z}(\mathcal{A})=f(x)-f(u) \geqslant \beta\|x-u\|^{2} \text { for all } x \in \mathbb{R}^{n} \text { with }\|x-u\| \leqslant \delta
$$

Now, using the same method of the proof as in part (1), we see that the conclusion holds.

### 4.1. Semismoothness of the maximum $Z$-eigenvalue function

In this subsection, as an application of the preceding stability result of the normalized eigenspace, we examine the semismoothness of the maximum $Z$-eigenvalue function. Now, consider a function $f: S \rightarrow \mathbb{R}$, where $S$ is the symmetric tensor space on $\mathbb{R}^{n}$. Note that the symmetric tensor space $S$ can be identified as a finite dimensional space with an appropriate dimension. The definition of semismoothness of $f$ can be translated as follows:

Definition 4.1
Let $f: S \rightarrow \mathbb{R}$ be a locally Lipschitz and directionally differentiable function. Then, the function $f: S \rightarrow \mathbb{R}$ is said to be semismooth at $\mathcal{A} \in S$ if

$$
f(\mathcal{A}+\Delta \mathcal{A})-f(\mathcal{A})-\langle V(\Delta \mathcal{A}), \Delta \mathcal{A}\rangle_{S}=o\left(\|\Delta \mathcal{A}\|_{S}\right), \forall V(\Delta \mathcal{A}) \in \partial_{C} f(\mathcal{A}+\Delta \mathcal{A})
$$

Moreover, $f: S \rightarrow \mathbb{R}$ is said to be $\rho$ th-order semismooth function for some $\rho \in(0,1]$ at $\mathcal{A} \in S$ if

$$
f(\mathcal{A}+\Delta \mathcal{A})-f(\mathcal{A})-\langle V(\Delta \mathcal{A}), \Delta \mathcal{A}\rangle_{S}=O\left(\|\Delta \mathcal{A}\|_{S}^{1+\rho}\right), \forall V(\Delta \mathcal{A}) \in \partial_{C} f(\mathcal{A}+\Delta \mathcal{A})
$$

In particular, if $\rho=1$, we say $f$ is a strongly semismooth function. We also say $f: S \rightarrow \mathbb{R}$ is a semismooth (resp. $\rho$ th-order semismooth) function on $S$ if $f$ is semismooth (resp. $\rho$ th-order semismooth) at $\mathcal{A}$ for all $\mathcal{A} \in S$.

To achieve this, we first observe from Lemma 3.1 that the maximum eigenvalue of tensors $\lambda_{1}^{Z}(\mathcal{A})$ is indeed the optimal value of a polynomial optimization problem $\left(P_{\mathcal{A}}\right) \max \left\{\mathcal{A} x^{m}:\|x\|=1\right\}$. Thus, the semismooth properties of the maximum $Z$-eigenvalue function can be approached by examining the stability of the parameterized optimization problem $\left(P_{\mathcal{A}}\right)$, which we have already studied earlier. We now study the $\rho$ th-order semismooth properties of the maximum $Z$-eigenvalue function and establish some explicit estimation of $\rho$.

## Theorem 4.2

Let $\mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor ( $m$ is even). Then, we have
(i) The maximum $Z$-eigenvalue function $\lambda_{1}^{Z}$ is at least $\rho$ th-order semismooth at $\mathcal{A}$ with $\rho=$ $\frac{1}{m(3 m-3)^{n-1}-1}$.
(ii) Moreover, if we further assume that the following second-order condition holds: for all $u \in E_{1}^{Z}(\mathcal{A})$

$$
\begin{equation*}
(m-1) \mathcal{A} u^{m-2}-\lambda_{1}^{Z}(\mathcal{A})\left((m-2) u u^{T}+\mathbf{I}_{\mathbf{n}}\right) \prec 0 \text { on } C_{u}=\left\{h \in \mathbb{R}^{n}: h^{T} u=0\right\} \tag{4.14}
\end{equation*}
$$

then the maximum $Z$-eigenvalue function $\lambda_{1}^{Z}$ is strongly semismooth at $\mathcal{A}$.

## Proof of (i)

Suppose that $m=2$. Then, $\mathcal{A}$ is an $(n \times n)$ symmetric matrix. From [18], we see that the maximum eigenvalue function is strongly semismooth. So, without loss of generality, we may assume that $m \geqslant 4$. Let $\mathcal{A}$ be an arbitrary $m$ th-order $n$-dimensional symmetric tensor, and let $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. Let $\Delta \mathcal{A}$ be an $m$ th-order $n$-dimensional symmetric tensor such that $\|\Delta \mathcal{A}\|_{S}>0$ and $\lambda_{1}$ is differentiable at $\mathcal{A}+\Delta \mathcal{A}$. So, Corollary 3.2 implies that $\nabla \lambda_{1}(\mathcal{A}+\Delta \mathcal{A})=\left(w_{\Delta \mathcal{A}}\right)^{m}$, where $w_{\Delta \mathcal{A}} \in E_{1}^{Z}(\mathcal{A}+\Delta \mathcal{A})$. Note that $\lambda_{1}$ is continuous and convex (and so, is directionally differentiable) and the Clarke subdifferential and the convex subdifferential of $\lambda_{1}$ coincide. To see the conclusion, we only need to show that

$$
\begin{equation*}
\lambda_{1}^{Z}(\mathcal{A}+\Delta \mathcal{A})-\lambda_{1}^{Z}(\mathcal{A})-\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S}=O\left(\|\Delta \mathcal{A}\|_{S}^{1+\rho}\right) \tag{4.15}
\end{equation*}
$$

To see (4.15), let $r>0$ be a constant such that

$$
\begin{equation*}
\left\|x^{m}-y^{m}\right\|_{S} \leqslant r\|x-y\| \text { for all } x, y \in K:=\{x:\|x\|=1\} \tag{4.16}
\end{equation*}
$$

Let $v \in E_{1}^{Z}(\mathcal{A})$ be such that $\left\|w_{\Delta \mathcal{A}}-v\right\|=d\left(w_{\Delta \mathcal{A}}, E_{1}^{Z}(\mathcal{A})\right)$. Then, the preceding proposition implies that there exists $c>0$ such that $\left\|w_{\Delta \mathcal{A}}-v\right\| \leqslant c\|\Delta \mathcal{A}\|_{S}^{\rho}$. Clearly, $v^{m} \in \partial \lambda_{1}^{Z}(\mathcal{A})$. It follows from (4.16) that

$$
\begin{aligned}
\lambda_{1}^{Z}(\mathcal{A}+\Delta \mathcal{A})-\lambda_{1}^{Z}(\mathcal{A})-\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S} & \geqslant\left\langle v^{m}, \Delta \mathcal{A}\right\rangle_{S}-\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S} \\
& \geqslant-\left\|v^{m}-\left(w_{\Delta \mathcal{A}}\right)^{m}\right\|_{S}\|\Delta \mathcal{A}\|_{S} \\
& \geqslant-r\left\|v-w_{\Delta \mathcal{A}}\right\| \Delta \mathcal{A} \|_{S} \\
& \geqslant-r c\|\Delta \mathcal{A}\|_{S}^{1+\rho}
\end{aligned}
$$

On the other hand, as $\nabla \lambda_{1}(\mathcal{A}+\Delta \mathcal{A})=\left(w_{\Delta \mathcal{A}}\right)^{m}$ and $\lambda_{1}^{Z}$ is convex, we see that

$$
\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m},-\Delta \mathcal{A}\right\rangle_{S}=\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \mathcal{A}-(\mathcal{A}+\Delta \mathcal{A})\right\rangle_{S} \leqslant \lambda_{1}^{Z}(\mathcal{A})-\lambda_{1}^{Z}(\mathcal{A}+\Delta \mathcal{A})
$$

This implies that

$$
\lambda_{1}^{Z}(\mathcal{A}+\Delta \mathcal{A})-\lambda_{1}^{Z}(\mathcal{A})-\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S} \leqslant\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S}-\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S}=0
$$

Therefore, we have

$$
\lambda_{1}^{Z}(\mathcal{A}+\Delta \mathcal{A})-\lambda_{1}^{Z}(\mathcal{A})-\left\langle\left(w_{\Delta \mathcal{A}}\right)^{m}, \Delta \mathcal{A}\right\rangle_{S}=O\left(\|\Delta \mathcal{A}\|_{S}^{1+\rho}\right)
$$

and so, the conclusion follows.

## Proof of (ii)

Using the same method of proof as in part (i) and using Proposition 4.1(ii) (instead of Proposition 4.1(i)), we see that the conclusion follows.

## Remark 4.1

In Theorem 4.2, for an $m$ th-order $n$-dimensional tensor $\mathcal{A}$, we showed that the maximum $Z$-eigenvalue function is always at least $\rho$ th-order semismooth at $\mathcal{A}$ with $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. In our preceding paper [14], $\rho$ th-order semismoothness of the maximum $H$-eigenvalue function at $\mathcal{A}$ with $\rho=\frac{1}{(2 m-1)^{n}}$ was shown under an additional assumption that the geometric multiplicity at $\mathcal{A}$ is one. At this moment, it is not clear whether our method of proof here can be used to relax the geometric multiplicity assumption in our previous $\rho$ th-order semismoothness result for maximum $H$-eigenvalue function with some appropriate exponent $\rho$ (as we made use of the fact that an eigenvector associated with the maximum $Z$-eigenvalue is of norm one in Proposition 4.1). Moreover, one can observe that the degree of the semismooth property is different for the maximum $H$-eigenvalue function and the maximum $Z$-eigenvalue function. At this moment, it is not clear for us which type of maximum eigenvalue function has a better analytic properties.

### 4.2. Nonsmooth Newton method for space tensor problems

Let $S(4,3)$ be the space consisting of all the fourth-order three-dimensional symmetric tensors. It is shown [6] that $S(4,3)$ is of dimension 15 , and so, there exists a one-to-one mapping $L: \mathbb{R}^{15} \rightarrow$ $S(4,3)$ (indeed the mapping $L$ can be explicitly constructed; see [6,33] for details). Let $n$ be a natural number, $\mathcal{A}_{i} \in S(4,3), i=0,1, \ldots, n$ and $b_{i} \in \mathbb{R}, i=1, \ldots, n$. Consider the STCLP problem, which was proposed and studied in [33]:

$$
\begin{aligned}
(S T C L P) \min _{\mathcal{X} \in S(4,3)} & \left\langle\mathcal{A}_{0}, \mathcal{X}\right\rangle_{S} \\
\text { s.t. } & \left\langle\mathcal{A}_{i}, \mathcal{X}\right\rangle_{S} \leqslant b_{i}, i=1, \ldots, n \\
& \mathcal{X} \in-\mathcal{C}(4,3)
\end{aligned}
$$

where $\mathcal{C}(4,3)$ is the cone of all negative semidefinite fourth-order three-dimensional symmetric tensor, that is, $\mathcal{C}(4,3):=\left\{\mathcal{A} \in S(4,3): \mathcal{A} x^{4} \leqslant 0, \forall x \in \mathbb{R}^{3}\right\}$. The problem (STCLP) arises from the medical imaging area where a high-order tensor is used to describe the non-Gaussian diffusion
feature [6]. Recently, Li et al. [14] proposed a nonsmooth Newton method based on the maximum $H$-eigenvalue for solving the space tensor problem and established that the corresponding nonsmooth Newton method converges superlinearly to a solution with order $1+\rho$ for some unknown constant $\rho>0$. Here, we propose a new nonsmooth Newton method based on the maximum $Z$-eigenvalue. An advantage of the new method is that we are now able to obtain the superlinear convergence with an explicit estimate of the order $1+\rho$, using the semismooth property of the maximum $Z$-eigenvalue established in Theorem 4.2.

Note that $\mathcal{C}(4,3)=\left\{\mathcal{A} \in S(4,3): \lambda_{1}^{Z}(\mathcal{A}) \leqslant 0\right\}$ and $\lambda_{1}^{Z}$ is convex. We see that problem (STCLP) is a convex programming problem. By assuming the Slater constraint qualification, solving (STCLP) is equivalent to solving its KKT system. As shown in [34] (see also [33]), $\mathcal{C}(4,3)$ is a self-dual cone and $(\mathcal{C}(4,3))^{*}=\mathcal{C}(4,3)=\mathcal{U}(4,3)$, where $(\mathcal{C}(4,3))^{*}$ is the usual dual cone of $\mathcal{C}(4,3)$, and $\mathcal{U}(4,3)$ is the rank one tensor space in $S(4,3)$ defined by $\mathcal{U}(4,3)=\{\mathcal{A} \in S(4,3): \mathcal{A}=$ $\left.\sum_{j=1}^{r}\left(a_{j}\right)^{4}, a_{j} \in \mathbb{R}^{3}, r \in \mathbb{N}\right\}$, where $a^{4}$ is the fourth-order three-dimensional symmetric rank one tensor defined by $\left(a^{4}\right)_{i_{1} i_{2} i_{3} i_{4}}=a_{i_{1}} a_{i_{2}} a_{i_{3}} a_{i_{4}}$ for each $i_{1}, i_{2}, i_{3}, i_{4} \in\{1,2,3\}$. So, its KKT system is to find $y_{1}, \ldots, y_{n} \in \mathbb{R}$ and $\mathcal{X} \in S(4,3)$ such that

$$
(K K T)\left\{\begin{array}{ll}
\mathcal{X} \in-\mathcal{C}(4,3) & \text { (Primal Feasible) }  \tag{4.17}\\
-\mathcal{A}_{0}-\sum_{i=1}^{n} y_{i} \mathcal{A}_{i} \in \mathcal{C}(4,3) & \text { (Dual Feasible) } \\
0 \leqslant y_{i} \perp\left(\left\langle\mathcal{A}_{i}, \mathcal{X}\right\rangle_{S}-b_{i}\right) \leqslant 0 & \text { (Complementary Slackness I) } \\
\left\langle\mathcal{A}_{0}+\sum_{i=1}^{n} y_{i} \mathcal{A}_{i}, \mathcal{X}\right\rangle_{S}=0 & \text { (Complementary Slackness II) }
\end{array} .\right.
$$

The following proposition establishes a useful observation: Solving the KKT problem is equivalent to solving the nonsmooth equation $F(x)=0$, where $F: \mathbb{R}^{n+15} \rightarrow \mathbb{R}^{n+2}$ is defined by

$$
F(x)=\left(\begin{array}{c}
\max \left\{\lambda_{1}^{Z}\left(-\mathcal{A}_{0}-\sum_{i=1}^{n} y_{i} \mathcal{A}_{i}\right), \lambda_{1}^{Z}(-L z)\right\}  \tag{4.18}\\
\left\langle\mathcal{A}_{0}+\sum_{i=1}^{n} y_{i} \mathcal{A}_{i}, L z\right\rangle_{S} \\
\max \left\{-y_{1},\left\langle\mathcal{A}_{1}, L z\right\rangle_{S}-b_{1}\right\} \\
\vdots \\
\max \left\{-y_{n},\left\langle\mathcal{A}_{n}, L z\right\rangle_{S}-b_{n}\right\}
\end{array}\right), x=(y, z) \in \mathbb{R}^{n} \times \mathbb{R}^{15},
$$

where $L: \mathbb{R}^{15} \rightarrow S(4,3)$ is the one-to-one linear mapping between $\mathbb{R}^{15}$ and $S(4,3)$. Its proof is similar to [14, Proposition 5.1], so we omit the proof here.

## Proposition 4.1

Let $y \in \mathbb{R}^{n}, \mathcal{X} \in S(4,3)$, and $x:=\left(y, L^{-1}(\mathcal{X})\right) \in \mathbb{R}^{n} \times \mathbb{R}^{15}$. Then, $(y, \mathcal{X}) \in \mathbb{R}^{n} \times S(4,3)$ solves the KKT system (4.17) if and only if $F(x)=0$.

From the proceeding proposition, to obtain a solution of the KKT system, it suffices to solve the nonsmooth underdetermined equation $F(x)=0$, where $F: \mathbb{R}^{n+15} \rightarrow \mathbb{R}^{n+2}$. Now, we state a nonsmooth Newton method for solving the space tensor conic linear problem (STCLP):

## Algorithm 1

Step 0. Choose $\left(y^{(0)}, \mathcal{X}^{(0)}\right) \in \mathbb{R}^{n} \times S(4,3)$. Compute $z^{(0)}=L^{-1}\left(\mathcal{X}^{(0)}\right)$, and let $x^{(0)}=$ $\left(y^{(0)}, z^{(0)}\right)$. If $F\left(x^{(0)}\right) \neq 0$, then set $k:=0$. Otherwise, output $\left(y^{(0)}, \mathcal{X}^{(0)}\right)$.
Step 1. Compute a $V_{k} \in \mathbb{R}^{(n+2) \times(n+15)}$ such that $V_{k} \in J_{C} F\left(x^{(k)}\right)$. ${ }^{\ddagger}$
Step 2. Let $x^{(k+1)}=x^{(k)}+\Delta x^{(k)}$, where $\Delta x^{(k)}=-\left(V_{k}^{T} V_{k}\right)^{-1} V_{k}^{T} F\left(x^{(k)}\right)$.

[^1]Step 3. If $F\left(x^{(k+1)}\right) \neq 0$, then replace $k$ by $k+1$ and go back to step 1 . Otherwise, let $x^{(k+1)}=\left(y^{(k+1)}, z^{(k+1)}\right) \in \mathbb{R}^{n} \times \mathbb{R}^{15}$ and output $\left(y^{(k+1)}, L^{-1}\left(z^{(k+1)}\right)\right) \in \mathbb{R}^{n} \times S(4,3)$.

Next, we present the superlinear convergence result of Algorithm 1 with an explicit estimate of the convergence rate for a regular solution (see the definition in the Appendix), by using the known convergence result for general nonsmooth Newton method (Lemma 7.1) and the semismooth property of maximum $Z$-eigenvalue (Theorem 4.2).

## Theorem 4.3

Let $\left(y^{*}, \mathcal{X}^{*}\right) \in \mathbb{R}^{n} \times S(4,3)$ be a solution of the KKT system of (STCLP). Let $x^{*}=$ $\left(y^{*}, L^{-1}\left(\mathcal{X}^{*}\right)\right) \in \mathbb{R}^{n} \times \mathbb{R}^{15}$. Let $x^{*}$ be a regular solution, and let $P$ be the regular space associated with $x^{*}$. Then, there exists a neighborhood $N$ of $x^{*}$ such that Algorithm 1 is well defined for any initial point $x^{(0)} \in N_{0} \cap\left(x^{*}+P\right)$ and Algorithm 1 either terminates infinitely many iterations or generates a sequence $\left\{x^{(k)}\right\}$ such that $x^{(k)}$ converges superlinearly to $x^{*}$ with order at least $1+\rho$, where $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$, that is,

$$
\begin{equation*}
\lim _{k \rightarrow \infty} \frac{\left\|x^{(k+1)}-x^{*}\right\|}{\left\|x^{(k)}-x^{*}\right\|^{1+\rho}}<+\infty \tag{4.19}
\end{equation*}
$$

## Proof

Clearly, $F$ is locally Lipschitz. Moreover, we note that $(a, b) \mapsto \max \{a, b\}$ is strongly semismooth, any continuous differentiable function with locally Lipschitz gradient is strongly semismooth, and composition of ( $\rho$ th-order) semismooth function is still ( $\rho$ th-order) semismooth. Then, Theorem 4.2(1) implies that $\mathcal{X} \mapsto \lambda_{1}^{Z}(\mathcal{X})$ is at least $\rho$ th-order semismooth with $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. It follows that $F$ is a vector valued function where each of its coordinate is at least $\rho$ th-order semismooth. Thus, $F$ is also at least $\rho$ th-order semismooth, and so, the conclusion follows from Lemma 7.1.

## 5. APPLICATION TO SPECTRAL HYPERGRAPH THEORY

Now, consider an (undirected) hypergraph, which is a pair $G=(V, E)$, where $V=\{1, \ldots, n\}$ is a finite set of vertices and $E \subseteq 2^{V}$ is a set of subsets of $V$ (each of which is called a hyperedge). A hypergraph $G$ is said to be $m$-uniform for an integer $m \geqslant 2$ if $|e|=m$ for all $e \in E$, where $|\cdot|$ denotes the cardinality. That is, for an $m$-uniform hypergraph, each hyperedge has the same cardinality $m$. If $m=2$, then 2-uniform graphs are typically called graphs. A finite path from vertex $i$ to vertex $j$ is a finite sequence of vertices with start from $i$ and end with $j$ such that each of its vertices and the next vertex belong to some hyperedge. Two vertices are called connected if there is a finite path between them. A connected component $X$ of $G$ is a subset of $V$ such that any two vertices in $X$ are connected and no other vertex in $V \backslash X$ is connected to any vertex in $X$.

Consider an $m$-uniform hypergraph $G=(V, E)$. Let $E=\left\{E_{1}, \ldots, E_{p}\right\}$, where each $E_{l}$, $l=1, \ldots, p$, is a hyperedge. We define a homogeneous polynomial $f$ associated with the hypergraph $G$ defined by

$$
f_{G}\left(x_{1}, \ldots, x_{n}\right)=\sum_{i=1}^{n} \sum_{j=1}^{n} \delta_{i j}\left(x_{i}-x_{j}\right)^{m}
$$

where $\delta_{i j}$ is defined by

$$
\delta_{i j}=\left\{\begin{array}{l}
1, \quad \text { if } \quad i<j \text { and }\{i, j\} \subseteq E_{l}, \text { for some } l=1, \ldots, p, \\
0, \quad \text { else. }
\end{array}\right.
$$

Note that any homogeneous polynomial uniquely determines a symmetric tensor. We can define the Laplacian tensor $\mathcal{L}$ of the hypergraph $G$ by

$$
\mathcal{L} x^{m}=f_{G}(x) \text { for all } x \in \mathbb{R}^{n}
$$

In the special case when $m=2$, our definition of Laplacian tensor reduces to the Laplacian matrix as the Laplacian matrix $L$ of a graph $G=(V, E)$ satisfies the following property: for all $x \in \mathbb{R}^{n}$,

$$
x^{T} L x=\sum_{i<j,(i, j) \in E}\left(x_{i}-x_{j}\right)^{2}
$$

## Remark 5.1

Recently, Hu and Qi [3] studied the algebraic connectivity of a 4-uniform hypergraph by introducing a different definition of the Laplacian tensor as follows: For a 4-uniform hypergraph $G=(V, E)$, its Laplacian tensor $T$ corresponds to the following quartic form:

$$
T x^{4}:=\sum_{E_{p} \in E} L\left(E_{p}\right) x^{4}
$$

where each $E_{p}$ is an hyperedge of $G$ and

$$
\begin{aligned}
L\left(E_{p}\right) x^{4}= & \frac{1}{84}\left[\left(x_{i}+x_{j}+x_{k}-3 x_{l}\right)^{4}+\left(x_{i}+x_{j}+x_{l}-3 x_{k}\right)^{4}\right. \\
& \left.+\left(x_{i}+x_{k}+x_{l}-3 x_{j}\right)^{4}+\left(x_{j}+x_{k}+x_{l}-3 x_{i}\right)^{4}\right]
\end{aligned}
$$

It can be seen that our definition of a Laplacian tensor is a direct extension of the matrix cases and works for $m$-uniform hypergraph with $m>4$.

We now define the characteristic tensor $\mathcal{C}$ of an $m$-uniform hypergraph $G$ by $\mathcal{C}=-\mathcal{L}$. The following theorem summarizes some basic features of the characteristic tensor.

## Theorem 5.1

Let $\mathcal{C}$ be the characteristic tensor of an $m$-uniform hypergraph $G$ where $m$ is even. Then, the following statements hold:
(1) The characteristic tensor $\mathcal{C}$ is negative semidefinite in the sense that $\mathcal{C} x^{m} \leqslant 0$ for all $x \in \mathbb{R}^{m}$.
(2) The maximum $Z$-eigenvalue of the characteristic tensor $\mathcal{C}$ is 0 , and $a=\frac{1}{\sqrt{n}} \mathbf{1}_{n}$ is a corresponding eigenvector where $\mathbf{1}_{n} \in \mathbb{R}^{n}$ is the vector whose components are all equal to one.
(3) The dimension of the eigenspace of the maximum $Z$-eigenvalue equals the number of the connected component of the $m$-uniform hypergraph $G$.

## Proof of (1)

Clearly, $f_{G}(x) \geqslant 0$ for all $x \in \mathbb{R}^{n}$. So, we see that

$$
\mathcal{C} x^{m}=-\mathcal{L} x^{m}=-f_{G}(x) \leqslant 0 \text { for all } x \in \mathbb{R}^{n}
$$

Proof of (2)
Consider $a=\frac{1}{\sqrt{n}} \mathbf{1}_{n}$. Then, $a=\left(a_{1}, \ldots, a_{n}\right)$ with $a_{i}=\frac{1}{\sqrt{n}}$ for each $i=1, \ldots, n$ and $\|a\|=1$. Note that

$$
\mathcal{C} a^{m}=-f_{G}(a)=-\sum_{i=1}^{n} \sum_{j=1}^{n} \delta_{i j}\left(a_{i}-a_{j}\right)^{m}=0
$$

 Thus, the conclusion follows from Lemma 3.1.

Proof of (3)
Let the connected component of $G$ be $\left\{V^{1}, \ldots, V^{q}\right\}$, where $q \in \mathbb{N}$. For each $l=1, \ldots, q$, define $a^{l}=\left(a_{1}^{l}, \ldots, a_{n}^{l}\right)$, where for each $i=1, \ldots, n$,

$$
a_{i}^{l}=\left\{\begin{array}{l}
1, \quad \text { if } \quad i \in V_{l}, \\
0, \\
\text { else }
\end{array}\right.
$$

Then, we see that

$$
\mathcal{C}\left(a^{l}\right)^{m}=-f_{G}\left(a^{l}\right)=-\sum_{i=1}^{n} \sum_{j=1}^{n} \delta_{i j}\left(a_{i}^{l}-a_{j}^{l}\right)^{m}=0
$$

where the last equality holds as $a_{i}^{l}-a_{j}^{l}=0$ if $i, j$ in the same connected component and $\delta_{i j}=0$ if $i, j$ in different connected components. So, each $a^{l}$ is an eigenvector associated with the maximum $Z$-eigenvalue of $\mathcal{L}, l=1, \ldots, q$. Note that $\left\{a^{1}, \ldots, a^{q}\right\}$ is linearly independent. So, the dimension of the eigenspace is at least $q$. Now, consider any eigenvector $x=\left(x_{1}, \ldots, x_{n}\right)$ associated with the maximum $Z$-eigenvalue. Then,

$$
\mathcal{C} x^{m}=-f_{G}(x)=-\sum_{i=1}^{n} \sum_{j=1}^{n} \delta_{i j}\left(x_{i}-x_{j}\right)^{m}=0
$$

As $m$ is even, this implies that $\delta_{i j}\left(x_{i}-x_{j}\right)=0$ for all $i, j=1, \ldots, n$. Take each connected component $V^{l}$, it follows that $x_{i}=x_{j}$ if $i, j \in V^{l}$. Note that $\bigcup_{l=1}^{q} V^{l}=V$. It follows that

$$
x=\alpha_{1} a^{1}+\cdots+\alpha_{q} a^{q}
$$

for some $\left(\alpha_{1}, \ldots, \alpha_{q}\right) \neq 0$. This shows that the dimension of the eigenspace equals $q$.
The following corollary provides the link between the combinatorial structure and the analytic structure of the hypergraph.

## Corollary 5.1

Let $G$ be an $m$-uniform hypergraph $G$ where $m$ is even, and let $\mathcal{C}$ be its characteristic tensor. Then, the following statements are equivalent:
(1) $G$ is connected.
(2) The geometric multiplicity of $\lambda_{1}^{Z}(\mathcal{C})$ is one.
(3) The maximum $Z$-eigenvalue function $\lambda_{1}^{Z}$ is differentiable at $\mathcal{C}$.

## Proof

$[(1) \Leftrightarrow(2)]$ This equivalence follows from the preceding theorem by letting the number of the connected component be one.
$[(2) \Leftrightarrow(3)]$ This equivalence follows by Corollary 3.2.

Next, we denote the second largest $Z$-eigenvalue of the characteristic tensor by $\lambda_{2}^{Z}(\mathcal{C})$, that is,

$$
\lambda_{2}^{Z}(\mathcal{C}):=\max \left\{\lambda \in \mathbb{R}: \lambda \text { is a Z-eigenvalue of } \mathcal{C} \text { and } \lambda \neq \lambda_{1}^{Z}(\mathcal{C})\right\} .
$$

The following proposition provides a variational characterization of the second largest $Z$-eigenvalue of the characteristic tensor.

## Proposition 5.1

Let $\mathcal{C}$ be the characteristic tensor of an $m$-uniform hypergraph $G$, where $m$ is even. Then, we have

$$
\lambda_{2}^{Z}(\mathcal{C})=\max _{x \in \mathbb{R}^{n}}\left\{\mathcal{C} x^{m}:\|x\|=1 \text { and } x \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}\right\} .
$$

## Proof

Consider the maximization problem

$$
\left(P_{0}\right) \max _{x \in \mathbb{R}^{n}}\left\{\mathcal{C} x^{m}:\|x\|^{m}=1 \text { and } x \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}\right\}
$$

To finish the proof, it suffices to show that $v\left(P_{0}\right)=\lambda_{2}^{Z}(\mathcal{C})$, where $v\left(P_{0}\right)$ is the optimal value of $\left(P_{0}\right)$.

Let $a$ be a maximizer of the problem $\left(P_{0}\right)$. Then, by the KKT condition, we have $\|a\|=1$, $a \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}$, and there exist $\lambda \in \mathbb{R}$ and $u \in \operatorname{span} E_{1}^{Z}(\mathcal{C})$ such that

$$
-m \mathcal{C} a^{m-1}+m \lambda a+u=0
$$

This implies that

$$
-m\left(\mathcal{C} a^{m-1}\right)^{T} u+\|u\|^{2}=\left(-m \mathcal{C} a^{m-1}+m \lambda a+u\right)^{T} u=0
$$

From the construction of the characteristic tensor, one can write $\mathcal{C}=\sum_{j=1}^{r} \gamma_{j} w_{j}^{m}$, where $\gamma_{j} \leqslant 0$, $w_{j} \in \mathbb{R}^{n}, j=1, \ldots, r, r \in \mathbb{N}$, and $y^{m}$ denotes the rank one tensor defined by $\left(y^{m}\right)_{i_{1}, \ldots, i_{m}}=$ $y_{i_{1}} \ldots y_{i_{m}}$. So, we have $\mathcal{C} a^{m-1}=\sum_{j=1}^{r} \gamma_{j}\left(\left\langle w_{j}, a\right\rangle\right)^{m-1} w_{j}$, and hence

$$
-m \sum_{j=1}^{r} \gamma_{j}\left(w_{j}^{T} a\right)^{m-1} w_{j}^{T} u+\|u\|^{2}=0
$$

As $u \in \operatorname{span} E_{1}^{Z}(\mathcal{C})$ and $E_{1}^{Z}(\mathcal{C})$ is symmetric (that is, if $v \in E_{1}^{Z}(\mathcal{C})$, then $-v \in E_{1}^{Z}(\mathcal{C})$ ), there exist $s \in \mathbb{N}, \alpha_{i} \geqslant 0, i=1, \ldots, s$, such that $u=\sum_{i=1}^{s} \alpha_{i} v_{i}$ with $v_{i} \in E_{1}^{Z}(\mathcal{C})$. Note that, for each $i=1, \ldots, s$,

$$
0=\lambda_{1}^{Z}(\mathcal{C})=\mathcal{C} v_{i}^{m}=\sum_{j=1}^{r} \gamma_{j}\left(w_{j}^{T} v_{i}\right)^{m}
$$

This together with $\gamma_{j} \leqslant 0$ and $m$ even implies that $\gamma_{j} w_{j}^{T} v_{i}=0$ for all $j=1, \ldots, r$ and $i=1, \ldots, s$. So, $\gamma_{j} w_{j}^{T} u=\gamma_{j} w_{j}^{T}\left(\sum_{i=1}^{s} \alpha_{i} v_{i}\right)=0$ for all $j=1, \ldots, r$, and hence

$$
0=-m \sum_{j=1}^{r} \gamma_{j}\left(w_{j}^{T} a\right)^{m-1} w_{j}^{T} u+\|u\|^{2}=\|u\|^{2}
$$

Thus, $u=0$ and $\mathcal{C} a^{m-1}=\lambda a$. So, $\lambda$ is a $Z$-eigenvalue of $\mathcal{C}$ with an eigenvector $a \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}$ with $\|a\|=1$. Moreover, we also have $\lambda=\lambda a^{T} a=\mathcal{C} a^{m}=v\left(P_{0}\right)$. We now show that $\lambda \neq \lambda_{1}(\mathcal{A})$. To see this, we proceed by the method of contradiction and suppose that $\lambda=\lambda_{1}^{Z}(\mathcal{C})$. Then, $\mathcal{C} a^{m}=\lambda_{1}^{Z}(\mathcal{C})$ and $\|a\|=1$. Then, we see that $a$ is an eigenvector of the maximum $Z$-eigenvalue. This contradicts the fact that $a \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}$ with $\|a\|=1$. So, by the definition of $\lambda_{2}^{Z}(\mathcal{C})$, we have

$$
v\left(P_{0}\right)=\lambda \leqslant \lambda_{2}^{Z}(\mathcal{C})
$$

To see the reverse inequality, let $a$ be an eigenvector of $\lambda_{2}^{Z}(\mathcal{C})$. Then, we can write $a=\rho x+u$ with $\rho \geqslant 0, x \in \operatorname{span} E_{1}^{Z}(\mathcal{C})$ with $\|x\|=1$ and $u \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}$. It follows that $a^{T} x=(\rho x+u)^{T} x=$ $\rho\|x\|^{2}=\rho$. Note that there exists $r \in \mathbb{N}$ such that $\mathcal{C}=\sum_{j=1}^{r} \gamma_{j} w_{j}^{m}$. It follows that

$$
\lambda_{2}^{Z}(\mathcal{C}) a=\mathcal{C} a^{m-1}=\sum_{j=1}^{r} \gamma_{j}\left(w_{j}^{T} a\right)^{m-1} w_{j}
$$

As $x \in \operatorname{span} E_{1}^{Z}(\mathcal{C})$, similar as before, we can show that $\gamma_{j} w_{j}^{T} x=0, j=1, \ldots, r$. It then follows that

$$
\rho \lambda_{2}^{Z}(\mathcal{C})=\lambda_{2}^{Z}(\mathcal{C}) a^{T} x=\sum_{j=1}^{r} \gamma_{j}\left(w_{j}^{T} a\right)^{m-1} w_{j}^{T} x=0
$$

As $\lambda_{2}^{Z}(\mathcal{C})<\lambda_{1}^{Z}(\mathcal{C})=0$. So, $\rho=0$. This implies that $a=u \in\left(\operatorname{span} E_{1}^{Z}(\mathcal{C})\right)^{\perp}$. This together with $\|a\|=1$ gives us that $a$ is feasible for $\left(P_{0}\right)$. Thus, we see that

$$
\lambda_{2}^{Z}(\mathcal{C}) \leqslant v\left(P_{0}\right)
$$

This completes the proof.
Consider a hypergraph $G=(V, E)$, where $V=\{1, \ldots, n\}$ is a finite set of vertices and $E \subseteq 2^{V}$ is a set consisting of all the hyperedges. The edge cut or coboundary, $E_{X}$, of the set $X \subseteq V$ is defined as the set of all hyperedges $e \in E$ such that there are two vertices $u, v \in e$ with $u \in X$ and $v \notin X$. A bisection of $G$ is a two-partition $\{X, Y\}$ of the vertex set $V=\{1, \ldots, n\}$ in which $|X|=|Y|$ if $n$ is even or $|X|=|Y|-1$ if $n$ is odd. The bisection problem is to find a bisection for which $E_{X}$ is as small as possible. The bipartition width $b w(G)$ of the hypergraph $G$ is defined as the optimal value of the bisection problem, that is,

$$
b w(G):=\min \left\{\left|E_{X}\right|: X \subseteq V ;|X|=\left[\frac{n}{2}\right]\right\},
$$

where $[a]$ denotes the integer part of the number $a$. Calculating the exact value of the bipartition width is, in general, a hard problem even for the graph case. Below, we shall see that one can use the second largest eigenvalue to provide a lower bound for the bipartition width of a connected hypergraph.

## Proposition 5.2

Let $G=(V, E)$ be an $m$-uniform connected hypergraph where $m$ is an even number and $V=\{1, \ldots, n\}$ is a finite set of vertices. Let $\mathcal{C}$ be the characteristic tensor of $G$ and let $\lambda_{2}^{Z}(\mathcal{C})$ be the second largest eigenvalue of $\mathcal{C}$. Let $X$ be a subset of the vertex set $V$. Then, we have

$$
\left|E_{X}\right| \geqslant \frac{-4 \lambda_{2}^{Z}(\mathcal{C})}{m^{2}}\left(\frac{|X||n-X|}{n}\right)^{\frac{m}{2}}
$$

Moreover, we have

$$
b w(G) \geqslant\left\{\begin{array}{cl}
\frac{-4 \lambda_{2}^{Z}(\mathcal{C})}{m^{2}}\left(\frac{n^{2}-1}{4 n}\right)^{\frac{m}{2}}, & \text { if } n \text { is odd } \\
\frac{-4 \lambda_{2}^{Z}(\mathcal{C})}{m^{2}}\left(\frac{n}{4}\right)^{\frac{m}{2}}, & \text { if } n \text { is even }
\end{array}\right.
$$

Proof
The inequality is immediate if $X=\emptyset$ and $X=V$. So, let us consider the case when $X$ is a proper nonempty subset of $V$. Let $w=\sum_{i \in X} e_{i}$, where $e_{i}=(0, \ldots, 0,1,0, \ldots, 0) \in \mathbb{R}^{n}$ is the vector whose $i$ th component is one and all the other components are all zero. Let $w=\beta \mathbf{1}_{n}+u$, where $\beta=\frac{|X|}{n}$ and $u=\sum_{i \in X} e_{i}-\frac{|X|}{n} \mathbf{1}_{n}$. Then, $u^{T} \mathbf{1}_{n}=0$. Note that $\mathbf{1}_{n} \in E_{1}^{Z}(\mathcal{C})$ and $\operatorname{span} E_{1}^{z}(\mathcal{C})$ is a one-dimensional subspace as $G$ is connected. So, $u \in\left(\operatorname{span} E_{1}^{z}(\mathcal{C})\right)^{\perp}$. It is easy to see that $\|u\|=\sqrt{\frac{|X|(n-|X|)}{n}}$. Let $\bar{u}=\frac{u}{\|u\|}$. Then, $\bar{u} \in\left(\operatorname{span} E_{1}^{z}(\mathcal{C})\right)^{\perp}$ and $\|\bar{u}\|=1$. So, from the preceding proposition, we see that

$$
\begin{equation*}
\lambda_{2}^{Z}(\mathcal{C}) \geqslant \mathcal{C} \bar{u}^{m}=\frac{\mathcal{C} u^{m}}{\|u\|^{m}}=\mathcal{C} u^{m}\left(\frac{n}{|X|(n-|X|)}\right)^{\frac{m}{2}} \tag{5.20}
\end{equation*}
$$

Now, from the construction of the characteristic tensor, one can write $\mathcal{C}=\sum_{k=1}^{r} \gamma_{k} v_{k}^{m}$ with $v_{k}^{T} \mathbf{1}_{n}=0$ for some $r \in \mathbb{N}$ and $v_{k} \in \mathbb{R}^{n}, k=1, \ldots, r$. So, we see that
$\mathcal{C} u^{m}=\mathcal{C}\left(w-\beta \mathbf{1}_{n}\right)^{m}=\sum_{k=1}^{r} \gamma_{k}\left(v_{k}^{T}\left(w-\beta \mathbf{1}_{n}\right)\right)^{m}=\sum_{k=1}^{r} \gamma_{k}\left(v_{k}^{T} w\right)^{m}=\mathcal{C} w^{m}=-\sum_{i, j} \delta_{i j}\left(w_{i}-w_{j}\right)^{m}$.
Note that $w=\sum_{i \in X} e_{i}$, and hence,

$$
\sum_{i, j} \delta_{i j}\left(w_{i}-w_{j}\right)^{m}=\sum_{i \in X, j \notin X, i<j,\{i, j\} \subseteq E} 1 \leqslant \max _{e \in E}\{|e \cap X|(m-|e \cap X|)\}\left|E_{X}\right| \leqslant \frac{m^{2}}{4}\left|E_{X}\right|
$$

where the last inequality follows as the discrete function $f:\{1, \ldots, m\} \rightarrow \mathbb{R}$ defined by $k \mapsto k(m-k)$ attains its maximum at $k=m / 2$. Thus, it follows from (5.20) and (5.21) that

$$
\begin{aligned}
\lambda_{2}^{Z}(\mathcal{C}) \geqslant \mathcal{C} u^{m}\left(\frac{n}{|X|(n-|X|)}\right)^{\frac{m}{2}} & =-\sum_{i, j} \delta_{i j}\left(w_{i}-w_{j}\right)^{m}\left(\frac{n}{|X|(n-|X|)}\right)^{\frac{m}{2}} \\
& \geqslant-\frac{m^{2}}{4}\left|E_{X}\right|\left(\frac{n}{|X|(n-|X|)}\right)^{\frac{m}{2}}
\end{aligned}
$$

and hence, the first assertion follows. The second assertion follows by taking $|X|=\frac{n}{2}$ if $n$ is even and $|X|=\frac{n-1}{2}$ if $n$ is odd in the first assertion.

Remark 5.2
When $m=2$, our lower bounds for the edge cut and bipartition width collapses the classical result for the connected graph cases (cf. [35,36]). As an illustration, when $m=2$, the lower bound for edge cuts reads

$$
\begin{aligned}
\left|E_{X}\right| & \geqslant-\lambda_{2}^{Z}(\mathcal{C})\left(\frac{|X||n-X|}{n}\right) \\
& =-\max \left\{x^{T} \mathcal{C} x: x^{T} \mathbf{1}_{n}=0,\|x\|=1\right\}\left(\frac{|X||n-X|}{n}\right) \\
& =\min \left\{x^{T} L x: x^{T} \mathbf{1}_{n}=0,\|x\|=1\right\}\left(\frac{|X||n-X|}{n}\right)
\end{aligned}
$$

where $L$ is the Laplacian matrix of the graph $G$. Note that $\min \left\{x^{T} L x: x^{T} \mathbf{1}_{n}=0,\|x\|=1\right\}=\mu_{1}$, where $\mu_{1}$ is the second smallest eigenvalue of the Laplacian matrix $L$. It follows that $\left|E_{X}\right| \geqslant$ $\mu_{1} \frac{|X||n-X|}{n}$, which is the classical result of the lower bound for the edge cut for a connected graph. The case for the bipartition width is also similar.

## 6. COMPUTATION OF THE LARGEST/SECOND-LARGEST Z-EIGENVALUE

In this section, we explain how one can compute the largest/second-largest $Z$-eigenvalue of symmetric tensors, using polynomial optimization techniques and our variational formula developed in the previous sections.

### 6.1. Computation of the largest Z-eigenvalue

In this subsection, we first discuss how to compute the largest $Z$-eigenvalue for a general symmetric tensor. From Lemma 3.1, the largest $Z$-eigenvalue of an $m$ th-order $n$-dimensional symmetric tensor can be computed as

$$
\lambda_{1}^{Z}(\mathcal{A})=\max _{\|x\|=1} \mathcal{A} x^{m}=-\min \{f(x): h(x)=0\}
$$

Let $f(x)=-\mathcal{A} x^{m}$ and $h(x)=\|x\|^{2}-1$. Then, the largest $Z$-eigenvalue cam be found by solving the following polynomial optimization problem.

$$
\left(P^{0}\right) \quad \lambda_{1}^{Z}(\mathcal{A})=-\min \{f(x): h(x)=0\}
$$

In general, solving $\left(P^{0}\right)$ is an NP-hard (non-deterministic polynomial-time hard) problem when $m \geqslant 4$. So, finding the largest $Z$-eigenvalue problem is, in general, also an NP-hard problem. Recently, Kolda and Mayo [21] proposed a shifted power method for finding a $Z$-eigenvalue of a general symmetric tensor. However, the $Z$-eigenvalue found by the method developed in [21] need not be the largest one. Moreover, another power-type method was proposed in [12] for calculating the largest $H$-eigenvalue (which is a different notion of $Z$-eigenvalue) of a nonnegative tensor. Unfortunately, this method does not work for finding the $Z$-eigenvalue and heavily relies on the nonnegative assumption. To the best of our knowledge, the only exact method for finding the largest $Z$-eigenvalue was established in [20], which only works for the fourth-order three-dimensional symmetric tensor. In what follows, we introduce a new method for finding the largest $Z$-eigenvalue for a general symmetric tensor, which utilizes the polynomial optimization technique. With the help of large-scale semidefinite programming problem (SDP) solvers, it can be used to find the largest $Z$-eigenvalue for a medium-size symmetric tensor.

To do this, we recall some basic facts as below. For a real polynomial (polynomial with real coefficients) $f$, we use $\operatorname{deg} f$ to denote the degree of $f$. We say that a real polynomial $f$ is sum of squares (SOS) if there exist real polynomials $f_{j}, j=1, \ldots, r$, such that $f=\sum_{j=1}^{r} f_{j}^{2}$. An important property of the SOS polynomials is that checking whether a polynomial is SOS or not is equivalent to solving an SDP. For details, see [37].

For each $k \in \mathbb{N}$, the $k$ th Lasserre's relaxation for solving $\left(P^{0}\right)$ is

$$
\left.\begin{array}{rl}
\left(R P_{k}^{0}\right) & \max
\end{array}\right)
$$

As shown in [37], for each fixed $k \in \mathbb{N},\left(R P_{k}^{0}\right)$ can be equivalently rewritten as an SDP. We note that an SDP can be efficiently solved and has found numerous application in various areas. Moreover, it is easy to see that, for each $k \in \mathbb{N}$, max $\left(R P_{k}^{0}\right) \leqslant \max \left(R P_{k+1}^{0}\right) \leqslant \min \left(P^{0}\right)$. We now show that the optimal value of the relaxation problem $\left(R P_{k}^{0}\right)$ asymptotically converges to $\min \left(P^{0}\right)$ using the celebrated positivity characterization from real algebraic geometry (Appendix).

## Proposition 6.1

It holds that $\lim _{k \rightarrow \infty} \max \left(R P_{k}^{0}\right)=\min \left(P^{0}\right)$.
Proof
Let $r:=\min \left(P^{0}\right)$ and fix any $\epsilon>0$. Then, we have $f(x)-r+\epsilon>0$ for all $x \in K$, where $K=\{x: h(x)=0\}=\{x: \pm h(x) \geqslant 0\}$. Define the quadratic module as follows:

$$
\mathbf{M}(h,-h)=\left\{\sigma_{0}+\left(\sigma_{1}-\sigma_{2}\right) h \mid \sigma_{i} \text { is SOS, } i=0,1,2,3,4\right\}
$$

It is clear that $-h \in \mathbf{M}(h,-h)$ and $\{x:-h(x) \geqslant 0\}$ is compact. Thus, Putinar positivstellensatz (Lemma 7.2) implies that $f-r+\epsilon \in M(h,-h)$. This shows that there exists $k \in \mathbb{N}$, $r-\epsilon \leqslant \max \left(R P_{k}\right)$. Thus, the conclusion follows.

Numerical experience indicates that (and recently was justified theoretically in a generical sense in [38]) the relaxation is often exact for small relaxation order $k$ (usually $k \leqslant 4$ ). Moreover, one can certify the global optimality by verifying some suitable technical condition called flat truncation condition [37]. On the other hand, the size of the equivalent SDP problem of the relaxation problem increases dramatically when the dimension/order of the tensor increases. For example, as illustrated in Table I, for a fourth-order 60-dimensional tensor, the equivalent SDP problem for the fourth relaxation problem has 1830 variables and 595, 664 constraints. Fortunately, a robust SDP
software (SDPNAL [39]) has been established very recently, which enables us to solve large-scale SDP (dimension up to 2000 and number of constraint of the SDP up to 1 million). This, in turn, helps us to find the largest $Z$-eigenvalue for medium size problem.

## Example 6.1

A fourth-order $n$-dimensional symmetric tensor defined by

$$
\mathcal{A}_{i j k l}=\frac{1}{24}(i+j-k-l), \quad 1 \leqslant i<j<k<l \leqslant n, \text { and } \mathcal{A}_{i j k l}=\mathcal{A}_{\sigma(i j k l)}
$$

where $\sigma(i j k l)$ denotes a permutation of its index $\{i, j, k, l\}$. The corresponding polynomial optimization problem $\left(P^{0}\right)$, in this case, becomes

$$
\begin{array}{ll}
\min & f(x):=\sum_{1 \leqslant i<j<k<l \leqslant n}(-i-j+k+l) x_{i} x_{j} x_{k} x_{l} \\
\text { s.t. } & h(x)=\|x\|^{2}-1=0
\end{array}
$$

## Example 6.2

An fourth-order $n$-dimensional symmetric tensor defined by

$$
\mathcal{A}_{i j k l}=\frac{1}{24}(-i-j-k-l), \quad 1 \leqslant i<j<k<l \leqslant n, \text { and } \mathcal{A}_{i j k l}=\mathcal{A}_{\sigma(i j k l)}
$$

where $\sigma(i j k l)$ denotes a permutation of its index $\{i, j, k, l\}$. The corresponding polynomial optimization problem $\left(P^{0}\right)$, in this case, becomes

$$
\begin{array}{ll}
\min & f(x):=\sum_{1 \leqslant i<j<k<l \leqslant n}(i+j+k+l) x_{i} x_{j} x_{k} x_{l} \\
\text { s.t. } & h(x)=\|x\|^{2}-1=0
\end{array}
$$

## Example 6.3

Let $n$ be an even number. Let $V=\{1, \ldots, n\}$. We generate a random graph $G=(V, E)$ with $|V|=n$ as follows. Select a random subset $M \subseteq V$ with $|M|=n / 2$. The edges $e_{i, j},(i, j) \nsubseteq M$, are generated with probability $1 / 2$. A fourth-order $n$-dimensional symmetric tensor is defined by

$$
\begin{gathered}
\mathcal{A}_{i i i i}=-1, i=1, \ldots, n \\
\mathcal{A}_{i i j j}=-\frac{1}{3}, \quad(i, j) \in E, \text { and } \mathcal{A}_{i j k l}=0 \text { otherwise. }
\end{gathered}
$$

The corresponding polynomial optimization problem $\left(P^{0}\right)$, in this case, becomes

$$
\begin{aligned}
\min & f(x):=\sum_{1 \leqslant i \leqslant n} x_{i}^{4}+2 \sum_{(i, j) \in E} x_{i}^{2} x_{j}^{2} \\
\text { s.t. } & h(x)=\|x\|^{2}-1=0
\end{aligned}
$$

In fact, the optimal value of this optimization problem $\left(P^{0}\right)$ indeed returns the stability number of the random graph $G$ we generated.

Table I summarizes the numerical results of Examples 6.1-6.3, where we compute the largest eigenvalue by first converting the fourth-order relaxation of the equivalent polynomial optimization to an SDP problem and solving this SDP problem using SDPNAL. All numerical experiments are performed on a desktop, with 3.47 GHz quad-core Intel E5620 Xeon 64-bit CPUs and 4 GB RAM, equipped with MATLAB 7.13 (R2011b). In particular, the data in Table I are explained as follows.

- $m$ : the order of the symmetric tensor;
- $n$ : the dimension of the symmetric tensor;

Table I. Numerical Results for Example 6.1-6.3

| Problem | m | n | NV | NC | $\lambda_{1}^{Z}(\mathcal{A})$ | Global optimality (Yes/No) | Time $(\mathrm{s})$ |
| :--- | :---: | :---: | ---: | :---: | :---: | :---: | :---: |
| Example 6.1 | 4 | 20 | 210 | 8854 | 21.4745 | Yes | 2.20 |
| Example 6.1 | 4 | 30 | 465 | 40,919 | 48.3792 | Yes | 44.59 |
| Example 6.1 | 4 | 40 | 820 | 123,409 | 87.1374 | Yes | 246.59 |
| Example 6.1 | 4 | 50 | 1275 | 292,824 | 140.405 | Yes | 1159.5 |
| Example 6.2 | 4 | 20 | 210 | 8854 | 46.0150 | Yes | 14.41 |
| Example 6.2 | 4 | 30 | 465 | 40,919 | 88.8139 | Yes | 49.72 |
| Example 6.2 | 4 | 40 | 820 | 123,409 | 136.4154 | Yes | 290.20 |
| Example 6.2 | 4 | 50 | 1275 | 292,824 | 187.6926 | Yes | 1494.64 |
| Example 6.3 | 4 | 40 | 820 | 123,409 | 15.9999 | Yes | 159.75 |
| Example 6.3 | 4 | 50 | 1275 | 292,824 | 18.9999 | Yes | 806.18 |
| Example 6.3 | 4 | 60 | 1830 | 595,664 | 24.0001 | Yes | 3387.17 |

- $N V$ : the number of variables of the equivalent SDP problem;
- $N C$ : the number of constraints in the equivalent SDP problem;
- $\lambda_{1}^{Z}(\mathcal{A})$ : the calculated largest $Z$-eigenvalue;
- Global optimality (Yes/No): whether the global optimality is certified or not; and
- Time: the CPU-time measured in seconds.

We observe that, for all the aforementioned numerical examples, the largest $Z$-eigenvalues can be found successfully for medium size tensors (dimension ranges from 20 to 60). On the other hand, in general, to find the largest eigenvalue of a large-scale tensor, one has to exploit the structure (e.g., sparsity structure) of the underlying tensor. It would be interesting to see how one could exploit the structure of a tensor to reduce the corresponding computation cost in the aforementioned algorithm.

### 6.2. Computation of the second largest eigenvalue for connected hypergraph

In this subsection, we discuss the computation of the second largest eigenvalue for the characteristic tensor of a hypergraph. For the characteristic tensor $\mathcal{C}$ of a connected hypergraph $G$, we have seen that its largest eigenvalue is zero and its second largest eigenvalue provides a lower bound for the bipartition width. Thus, it is important to compute or estimate the second largest eigenvalue of the characteristic tensor. From our variational formula for second largest eigenvalue (Proposition 5.1) and noting that $G$ is connected, Theorem 5.1 implies that

$$
\lambda_{2}^{Z}(\mathcal{C})=\max \left\{\mathcal{C} x^{m}: \mathbf{1}_{n}^{T} x=0,\|x\|=1\right\}
$$

Let $f(x)=-\mathcal{C} x^{m}, h_{1}(x)=\mathbf{1}_{n}^{T} x$, and $h_{2}(x)=\|x\|^{2}-1$. Then, the second largest eigenvalue for Laplacian tensor for a connected hypergraph $G$ equals the negative optimal value of the following polynomial optimization problem

$$
(P) \quad \min \left\{f(x): h_{1}(x)=0, h_{2}(x)=0\right\}
$$

For each $k \in \mathbb{N}$, the $k$ th Lasserre's relaxation for solving $(P)$ is

$$
\begin{array}{cc}
\left(R P_{k}\right) \quad \max & \gamma \\
\text { s.t. } & f(x)-\gamma=\sigma_{0}(x)+\phi_{l}(x) h_{1}(x)+\phi_{2}(x) h_{2}(x) \\
& \sigma_{0} \text { is } \operatorname{SOS}, \operatorname{deg} \sigma_{0} \leqslant k, \\
& \phi_{l} \text { are real polynomials, } \operatorname{deg}\left(\phi_{l} h_{l}\right) \leqslant k, l=1,2
\end{array}
$$

Similarly, we can show that the optimal value of $(P)$ can be approximated by a sequence of SDP.

## Proposition 6.2

It holds that $\lim _{k \rightarrow \infty} \max \left(R P_{k}\right)=\min (P)$.

## Proof

Let $r:=\min (P)$ and fix any $\epsilon>0$. Then, we have $f(x)-r+\epsilon>0$ for all $x \in K$, where $K=\left\{x: h_{1}(x)=0, h_{2}(x)=0\right\}=\left\{x: \pm h_{1}(x) \geqslant 0, \pm h_{2}(x) \geqslant 0\right\}$. Define the quadratic module as follows:

$$
M\left(h_{1},-h_{1}, h_{2},-h_{2}\right)=\left\{\sigma_{0}+\left(\sigma_{1}-\sigma_{2}\right) h_{1}+\left(\sigma_{3}-\sigma_{4}\right) h_{2} \mid \sigma_{i} \text { is SOS, } i=0,1,2,3,4\right\} .
$$

It is clear that $-h_{2} \in M\left(h_{1},-h_{1}, h_{2},-h_{2}\right)$ and $\left\{x:-h_{2}(x) \geqslant 0\right\}$ is compact. Thus, Putinar positivstellensatz (Lemma 7.2) implies that $f-r+\epsilon \in M\left(h_{1},-h_{1}, h_{2},-h_{2}\right)$. This shows that there exists $k \in \mathbb{N}, r-\epsilon \leqslant \max \left(R P_{k}\right)$. Thus, the conclusion follows.

In what follows, as an illustration, we use two simple numerical examples to explain how one can compute/approximate the second largest eigenvalue for the characteristic tensor of a hypergraph (and so, the bipartition width) using the common global polynomial optimization software Gloptipoly3 [40] (for other popular software see [46, 47]). An important feature of the software Gloptipoly3 is that it can certify the global optimality using tools from real algebraic geometry [40, 44].

## Example 6.4

Consider the following connected uniform four-hypergraph $G=(V, E)$, where $V=\{1,2,3,4,5,6\}$ and $E=\{(1,2,3,4),(1,2,5,6),(3,4,5,6)\}$. It can be verified that

$$
\begin{aligned}
\lambda_{2}^{Z}(\mathcal{C}) & =\max \left\{\mathcal{C} x^{m}: \mathbf{1}_{n}^{T} x=0,\|x\|=1\right\} \\
& =-\min \left\{f(x): \sum_{i=1}^{6} x_{i}=0, \sum_{i=1}^{6} x_{i}^{2}-1=0\right\},
\end{aligned}
$$

where

$$
\begin{aligned}
f(x)= & \left(x_{1}-x_{2}\right)^{4}+\left(x_{1}-x_{3}\right)^{4}+\left(x_{1}-x_{4}\right)^{4}+\left(x_{1}-x_{5}\right)^{4}+\left(x_{1}-x_{6}\right)^{4} \\
& +\left(x_{2}-x_{3}\right)^{4}+\left(x_{2}-x_{4}\right)^{4}+\left(x_{2}-x_{5}\right)^{4}+\left(x_{2}-x_{6}\right)^{4} \\
& +\left(x_{3}-x_{4}\right)^{4}+\left(x_{3}-x_{5}\right)^{4}+\left(x_{3}-x_{6}\right)^{4}+\left(x_{4}-x_{5}\right)^{4} \\
& +\left(x_{4}-x_{6}\right)^{4}+\left(x_{5}-x_{6}\right)^{4} .
\end{aligned}
$$

As we explained before, the problem $(P) \min \left\{f(x): \sum_{i=1}^{6} x_{i}=0, \sum_{i=1}^{6} x_{i}^{2}-1=0\right\}$ can be solved by using Gloptipoly3. Indeed, by running the following simple code using Gloptipoly3, the software indicates that the fourth-order relaxation $\left(R P_{4}\right)$ gives the global minimum 4 of $(P)$ and returns a global minimizer $[-0.4082,0.4082,-0.4082,0.4082,0.4082,-0.4082]^{T}$.
Thus, we see that $\lambda_{2}^{Z}(\mathcal{C})=-4$. Letting $m=4$ and $n=6$, our estimate of bipartition width gives us that $b w(G) \geqslant \frac{-4 \lambda_{2}^{Z}(\mathcal{C})}{m^{2}}\left(\frac{n}{4}\right)^{\frac{m}{2}}=\frac{9}{4}$. Note that $b w(G)$ must be an integer. So, $b w(G)=3$. On the other hand, it is easy to verify directly from the definition of bipartition width that $b w(G)=3$ in this case.

## Example 6.5

Consider the following connected uniform four-hypergraph $G=(V, E)$, where $V=$ $\{1,2,3,4,5,6,7,8,9,10\}$ and

$$
E=\{(1,2,3,4),(3,4,5,6),(5,6,7,8),(6,7,8,9),(7,8,9,10),(1,2,9,10)\} .
$$

It can be verified that

$$
\begin{aligned}
\lambda_{2}^{Z}(\mathcal{C}) & =\max \left\{\mathcal{C} x^{m}: \mathbf{1}_{n}^{T} x=0,\|x\|=1\right\} \\
& =-\min \left\{f(x): \sum_{i=1}^{10} x_{i}=0, \sum_{i=1}^{10} x_{i}^{2}-1=0\right\},
\end{aligned}
$$

where

$$
\begin{aligned}
f(x)= & \sum_{j \in\{2,3,4,9,10\}}(x(1)-x(j))^{4}+\sum_{j \in\{3,4,9,10\}}(x(2)-x(j))^{4}+\sum_{j \in\{4,5,6\}}(x(3)-x(j))^{4} \\
& +\sum_{j \in\{5,6\}}(x(4)-x(j))^{4}+\sum_{j \in\{\{6,7,8\}}(x(5)-x(j))^{4}+\sum_{j \in\{7,8,9\}}(x(6)-x(j))^{4} \\
& \sum_{j \in\{8,9,10\}}(x(7)-x(j))^{4}+\sum_{j \in\{9,10\}}(x(8)-x(j))^{4}+(x(9)-x(10))^{4} ;
\end{aligned}
$$

As we explained before, the problem $(P) \min \left\{f(x): \sum_{i=1}^{10} x_{i}=0, \sum_{i=1}^{10} x_{i}^{2}-1=0\right\}$ can be solved by using Gloptipoly3. Indeed, by running a similar simple code as in the preceding example via Gloptipoly3, the software indicates that the third-order relaxation $\left(R P_{3}\right)$ gives the global minimum 0.4 of $(P)$ and returns a global minimizer

$$
[0.4509,0.4509,0.2129,0.2129,-0.1939,-0.1879,-0.4636,-0.4636,-0.0107,-0.0079]^{T}
$$

So, in this case, $\lambda_{2}(\mathcal{C})=-0.4920$.

## 7. CONCLUSION AND REMARKS

In this paper, using variational analysis techniques, we examined some fundamental analytic properties of $Z$-eigenvalues of a symmetric tensor with even order. As applications, we introduced the characteristic tensor of a hypergraph and showed that the maximum $Z$-eigenvalue function of the associated characteristic tensor provides a natural link for the combinatorial structure and the analytic structure of the underlying hypergraph. We also established a variational formula for the second largest $Z$-eigenvalue for the characteristic tensor of a hypergraph.

Below, we present a few open questions and remarks:

- For an $m$ th-order $n$-dimensional tensor $\mathcal{A}$, we showed that the maximum $Z$-eigenvalue function is always at least $\rho$ th-order semismooth at $\mathcal{A}$ with $\rho=\frac{1}{m(3 m-3)^{n-1}-1}$. In our preceding paper [14], $\rho$ th-order semismoothness of the maximum $H$-eigenvalue function at $\mathcal{A}$ with $\rho=\frac{1}{(2 m-1)^{n}}$ is shown under an additional assumption that the geometric multiplicity at $\mathcal{A}$ is one. It would be interesting to see whether our method of proof here can be used to relax the geometric multiplicity assumption in our previous $\rho$ th-order semismoothness result for maximum $H$-eigenvalue function with some appropriate exponent $\rho$.
- We made use of the concept of $Z$-eigenvalue to study some very basic properties of hypergraphs. It would be useful to exploit more in this direction to establish further results (for example, bounds on the maximal cliques and chromatic number). Recently, some progresses has been made in [2] using the concept of $H$-eigenvalues via an algebraic approach. It would be also interesting to see how one could link these two approaches together.
- We established the variational characterization of the second largest $Z$-eigenvalue for the characteristic tensor of a uniform graph. It would be useful to see whether this variational characterization continues to hold for more general tensors or not. Moreover, we have explained how one can compute the largest and second largest $Z$-eigenvalues for using polynomial optimization techniques together with our variational characterization. Detail study of the convergence rate and effective error estimate of this method would be interesting topics to examine.

These will be our future research topics and will be examined in a forthcoming study.

## APPENDIX A

## Nonsmooth Newton method for underdetermined equations

Consider a general nonsmooth equation $G(x)=0$, where $G: \mathbb{R}^{m} \rightarrow \mathbb{R}^{l}$ with $m \geqslant l$. The general algorithm of the nonsmooth Newton method for a underdetermined equation is stated as follows:

## Algorithm 0

Step 0. Choose $x^{(0)} \in \mathbb{R}^{m}$. If $G\left(x^{(0)}\right) \neq 0$, then set $k:=0$ and go to step 1 . Otherwise, output $x^{(0)}$.
Step 1. Compute a $V_{k} \in \mathbb{R}^{l \times m}$ such that $V_{k} \in J_{C} G\left(x^{(k)}\right)$.
Step 2. Let $x^{(k+1)}=x^{(k)}+\Delta x^{(k)}$, where

$$
\Delta x^{(k)}=-V_{k}^{T}\left(V_{k} V_{k}^{T}\right)^{-1} G\left(x^{(k)}\right)
$$

Step 3. If $G\left(x^{(k+1)}\right) \neq 0$, then replace $k$ by $k+1$ and go back to step 1 . Otherwise, output $x^{(k+1)} \in \mathbb{R}^{m}$.

We now recall the definition of regular solution and the local convergence result of a nonsmooth Newton method for underdetermined equations, which was presented in [14].

## Definition 7.1

For a nonlinear equation $G(x)=0$, where $G: \mathbb{R}^{m} \rightarrow \mathbb{R}^{l}$ with $m \geqslant l$. We say $x^{*}$ is a regular solution of this equation if $G\left(x^{*}\right)=0$ and there exists a neighborhood $N$ of $x^{*}$ such that
(a) $\operatorname{rank}(V)=l$ for all $V \in J_{C} G(x)$ and $x \in N \cap\{x: G(x) \neq 0\}$; and
(b) $R\left(V^{T}\left(V V^{T}\right)^{-1}\right) \equiv P$ for all $V \in J_{C} G(x)$ and $x \in N \cap\{x: G(x) \neq 0\}$, where $R(A)$ denotes the range of an $(m \times l)$ matrix $A$, which is defined by $R(A)=\left\{A x: x \in \mathbb{R}^{l}\right\} \subseteq \mathbb{R}^{m}$, and $P$ is some vector space in $\mathbb{R}^{m}$.
Moreover, the vector space $P$ in (b) is called the regular space associated with $x^{*}$.
Lemma 7.1 (cf. [14, Theorem 5.1])
Consider an underdetermined equation $G(x)=0$, where $G: \mathbb{R}^{m} \rightarrow \mathbb{R}^{l}$ with $m \geqslant l$. Let $x^{*}$ be a regular solution and $P$ be the regular space associated with $x^{*}$. If we further assume that $G$ is $\rho$ th-order semismooth for some $\rho \in(0,1]$. Then, Algorithm 0 either terminates infinitely many iterations or generates a sequence $\left\{x^{(k)}\right\}$ such that $x^{(k)}$ converges to $x^{*}$ with order $(1+\rho)$, that is,

$$
\lim _{k \rightarrow \infty} \frac{\left\|x^{(k+1)}-x^{*}\right\|}{\left\|x^{(k)}-x^{*}\right\|^{1+\rho}}<+\infty
$$

## Positivity characterization from real algebraic geometry

A quadratic module generated by polynomials $g_{1}, \ldots, g_{m} \in \mathbb{R}[x]$ is defined as $\mathbf{M}\left(g_{1}, \ldots, g_{m}\right):=$ $\left\{\sigma_{0}+\sigma_{1} g_{1}+\cdots+\sigma_{m} g_{m} \mid \sigma_{i}\right.$ is sum of squares polynomial, $\left.i=0,1, \ldots, m\right\}$. We say $\mathbf{M}\left(g_{1}, \ldots, g_{m}\right)$ is archimedean if there exists $p \in \mathbf{M}\left(g_{1}, \ldots, g_{m}\right)$ such that $\{x: p(x) \geqslant 0\}$ is compact.

We now recall the following important certificate for positivity of a polynomial over a semialgebraic set under the assumption that the associated quadratic module is archimedean (and hence, the semialgebraic set must be compact).

## Lemma 7.2 (Putinar positivstellensatz)

Let $f, g_{j}, j=1, \ldots, m$, be real polynomials with $K:=\left\{x: g_{j}(x) \geqslant 0, j=1, \ldots, m\right\} \neq$ $\emptyset$. Suppose that $f(x)>0$ for all $x \in K$ and $M\left(g_{1}, \ldots, g_{m}\right)$ is archimedean. Then, $f \in$ $M\left(g_{1}, \ldots, g_{m}\right)$.

## ACKNOWLEDGEMENTS

The authors are grateful to the referees and the editor for their constructive comments and helpful suggestions that have contributed to the final preparation of the paper. This work was partially supported by the Hong Kong Research Grant Council (grant nos PolyU 501909, 502510, 502111, and 501212), the National Natural Science Foundation of China (grant nos 11001060 and 61262026), the JGZX programm of Jiangxi Province (20112BCB23027), the science and technology programm of Jiangxi Education Committee (LDJH12088), and a research grant from Australian Research Council.

## REFERENCES

1. Ni Q, Qi L, Wang F. An eigenvalue method for the positive definiteness identification problem. IEEE Transactions on Automatic Control 2008; 53:1096-1107.
2. Cooper J, Dutle A. Spectra of uniform hypergraphs,. Linear Algebra and its Applications 2012; 436:3268-3292.
3. Hu S, Qi L. Algebraic connectivity of an even uniform hypergraph. Journal of Combinatorial Optimization 2012; 24:564-579.
4. Li X, Ng M, Ye Y. Finding stationary probability vector of a transition probability tensor arising from a higher-order Markov chain, 2011. preprint.
5. Qi L, Teo KL. Multivariate polynomial minimization and its application in signal processing. Journal of Global Optimization 2003; 46:419-433.
6. Qi L, Yu G, Wu EX. Higher order positive semi-definite diffusion tensor imaging. SIAM Journal on Imaging Sciences 2010; 3:416-433.
7. De Lathauwer L, De Moor B. From matrix to tensor: multilinear algebra and signal processing. In Mathematics in Signal Processing IV, Selected Papers Presented at 4th IMA Int. Conf. on Mathematics in Signal Processing, McWhirter J (ed.). Oxford University Press: Oxford, United Kingdom, 1998; 1-15.
8. De Lathauwer L, De Moor B, Vandewalle J. On the best rank-1 and rank- $\left(R^{1}, R^{2}, \ldots, R^{N}\right)$ approximation of higher-order tensor. SIAM Journal on Matrix Analysis and Applications 2000; 21:1324-1342.
9. Kofidis E, Regalia Ph. On the best rank-1 approximation of higher-order symmetric tensors. SIAM Journal on Matrix Analysis and Applications 2002; 23:863-884.
10. Kolda TG, Bader BW. Tensor decompositions and applications. SIAM Review 2009; 51(3):455-500.
11. Lim LH. Singular values and eigenvalues of tensors: a variational approach. Proceeding of 1st IEEE International Workshop on Computational Advances of Multi-Tensor Adaptive Processing, Puerto Vallarta, 2005; 129-132.
12. Ng M, Qi L, Zhou G. Finding the largest eigenvalue of a nonnegative tensor. SIAM Journal on Matrix Analysis and Applications 2009; 31:1090-1099.
13. Qi L. Eigenvalues of a real symmetric tensor. Journal of Symbolic Computation 2005; 40:1302-1324.
14. Li G, Qi L, Yu G. Semismoothness of the maximum $H$-eigenvalue function of a symmetric tensor and its application. Linear Algebra and its Applications 2013; 438(2):813-833.
15. Zalinescu C. Convex Analysis in General Vector Spaces. World Scientific: River Edge, New Jersey, 2002.
16. Mifflin R. Semismooth and semiconvex functions in constrained optimization. SIAM Journal on Control and Optimization 1977; 15(6):959-972.
17. Qi L, Sun J. A nonsmooth version of Newton's method. Mathematical Programming 1993; 78:353-368.
18. Sun D, Sun J. Semismooth matrix-valued functions. Mathematics of Operation Research 2002; 27:150-169.
19. Chang KC, Pearson K, Zhang T. On eigenvalue problems of real symmetric tensors. Journal of Mathematical Analysis and Applications 2009; 350:416-422.
20. Qi L, Wang F, Wang Y. Z-eigenvalue methods for a global polynomial optimization problem. Mathematical Programming 2009; 118:301-316.
21. Kolda TG, Mayo JR. Shifted power method for computing tensor eigenpairs. SIAM Journal on Matrix Analysis and Applications 2011; 32:1095-1124.
22. Hu S, Li G, Qi L, Song Y. Finding the maximum eigenvalue of essentially nonnegative symmetric tensors via sum of squares programming. Journal of Optimization Theory and Applications 2013. DOI:10.1007/s10957-013-0293-9.
23. Hiriart-Urruty JB, Lemaréchal C. Convex Analysis and Minimization Algorithms. I. Fundamentals, Grundlehren der Mathematischen Wissenschaften [Fundamental Principles of Mathematical Sciences], Vol. 305. Springer-Verlag: Berlin, 1993.
24. Lewis AS. Nonsmooth analysis of eigenvalues. Mathematical Programming 1999; 84:1-24.
25. Ye D. Sensitivity analysis of the greatest eigenvalue of a symmetric matrix via the $\epsilon$-subdifferential of the associated convex quadratic form. Journal of Optimization Theory and Applications 1993; 76(2):287-304.
26. Gwoździewicz J. The Łojasiewicz exponent of an analytic function at an isolated zero. Commentarii Mathematici Helvetici 1999; 74:364-375.
27. Li G. On the asymptotically well behaved functions and global error bound for convex polynomials. SIAM Journal on Optimization 2010; 20(4):1923-1943.
28. Li G, Ng KF. Error bounds of generalized D-gap functions for nonsmooth and nonmonotone variational inequality problems. SIAM Journal on Optimization 2009; 20(2):667-690.
29. Li G, Tang C, Wei Z. Error bound results for generalized D-gap functions of nonsmooth variational inequality problems. Journal of Computational and Applied Mathematics 2010; 233(11):2795-2806.
30. D'Acunto D, Kurdyka K. Explicit bounds for the Łojasiewicz exponent in the gradient inequality for polynomials. Annales Polonici Mathematici 2005; 87:51-61.
31. Ngai HV, Thèra M. Error bounds for systems of lower semicontinuous functions in Asplund spaces. Mathematical Programming 2009; 116(1-2):397-427.
32. Fiacco AV. Second order sufficient conditions for weak and strict constrained minima. SIAM Journal on Applied Mathematics 1968; 16:105-108.
33. Qi L, Ye Y. Space tensor conic programming. preprint.
34. Reznick B. Some concrete aspects of Hilbert's 17th Problem. In Real Algebraic Geometry and Ordered Structures (Baton Rouge, LA , 1996), Vol. 253, Contemporary Mathematics. American Mathematical Society: Providence, RI, 2000; 251-272.
35. Mohar B. Some applications of Laplace eigenvalues of graphs. In Graph Symmetry: Algebraic Methods and Applications, Vol. 497, Hahn G, Sabidussi G (eds), NATO ASI Ser. C. Kluwer: Dordrecht, 1997; 225-275.
36. Mohar B. Laplace eigenvalues of graphs: a survey. Discrete Mathematics 1992; 109:171-183.
37. Lasserre JB. Global optimization with polynomials and the problem of moments. SIAM Journal on Optimization 2001; 11:796-817.
38. Nie JW. Optimality conditions and finite convergence of Lasserre's hierarchy, 2012. arXiv:1206.0319 [math.OC].
39. Zhao X, Sun D, Toh KC. A Newton-CG augmented lagrangian method for semidefinite programming. SIAM Journal on Optimization 2010; 20:1737-1765.
40. Henrion D, Lasserre JB, Loefberg J. GloptiPoly 3: moments, optimization and semidefinite programming. Optimization Methods and Software 2009; 24:761-779.
41. Clarke FH. Optimization and Nonsmooth Analysis. Wiley: New York, 1983.
42. Bonnans JF, Shapiro A. Perturbation Analysis of Optimization Problems. Springer: New York, 2000.
43. Chang KC, Pearson K, Zhang T. Perron-Frobenius theorem for nonnegative tensors. Communications in Mathematical Sciences 2008; 6:507-520.
44. Jeyakumar V, Li G. Necessary global optimality conditions for nonlinear programming with polynomial constraint. Mathematical Programming 2011; 126(2):393-399.
45. Li G. Global error bounds for piecewise convex polynomials. Mathematical Programming 2013; 1-2:37-34. DOI: 10.1007/s10107-011-0481-z.
46. Löfberg J. Pre- and post-processing sum-of-squares programs in practice. IEEE Transactions on Automatic Control 2009; 54:1007-1011.
47. Löfberg J. YALMIP: a toolbox for modeling and optimization in MATLAB. In Proceedings of the CACSD Conference, Taipei, 2004; 284-289.
48. Pang J, Qi L. Nonsmooth equations: motivation and algorithms. SIAM Journal on Optimization 1993; 3(3):443-465.

[^0]:    *Correspondence to: Guoyin Li, Department of Applied Mathematics, University of New South Wales, Sydney 2052, Australia.
    ${ }^{\dagger}$ E-mail: g.li@unsw.edu.au

[^1]:    ${ }^{*}$ For a fourth-order three-dimensional tensor, one can efficiently find an eigenvector associated with its maximum eigenvalue (e.g., see [6]), and so, one can also efficiently compute a member of the (Clarke) generalized Hessian $J_{C} F\left(x^{(k)}\right) \subseteq \mathbb{R}^{(n+2) \times(n+15)}$.

